



Shifted Partial Quadratics, their Pushforwards, and Signature Invariants for Tangles

Abstract. Following a general discussion of the computation of zombians of unfinished columbaria (with examples), I will tell you about my recent joint work w/ Jessica Liu on what we feel is the “textbook” extension of knot signatures to tangles, which for unknown reasons, is not in any of the textbooks that we know.



Jessica Liu



Columbaria in an East Sydney Cemetery



Jacobian, Hamiltonian, Zombian

Kashaev's Conjecture [Ka]

$$\sigma_{Kas} = 2\sigma_{TL}.$$

Liu's Theorem [Li].

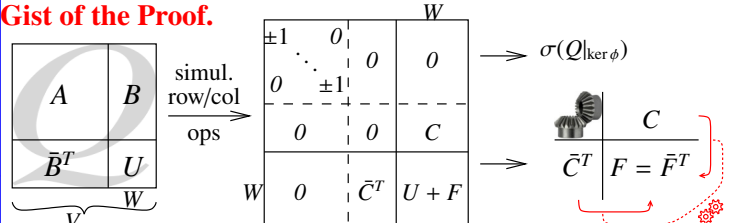
A **Partial Quadratic (PQ)** on V is a quadratic Q defined only on a subspace $\mathcal{D}_Q \subset V$. We add PQs with $\mathcal{D}_{Q_1+Q_2} := \mathcal{D}_{Q_1} \cap \mathcal{D}_{Q_2}$. Given a linear $\psi: V \rightarrow W$ and a PQ Q on W , there is an obvious **pullback** ψ^*Q , a PQ on V .

Theorem 1. Given a linear $\phi: V \rightarrow W$ and a PQ Q on V , there is a unique **pushforward** PQ ϕ_*Q on W such that for every PQU on W , $\sigma_V(Q + \phi^*U) = \sigma_{\ker\phi}(Q|_{\ker\phi}) + \sigma_W(U + \phi_*Q)$.

(If you must, $\mathcal{D}(\phi_*Q) = \phi(\text{ann}_Q(\mathcal{D}(Q) \cap \ker\phi))$ and $(\phi_*Q)(w) = Q(v)$, where v is s.t. $\phi(v) = w$ and $Q(v, \text{rad } Q|_{\ker\phi}) = 0$).

Prior Art on signatures for tangles / braids. Gambaudo and Ghys [GG], Cimasoni and Conway [CC], Conway [Co], Merz [Me]. All define signatures of tangles / braids by first closing them to links and then work hard to derive composition properties.

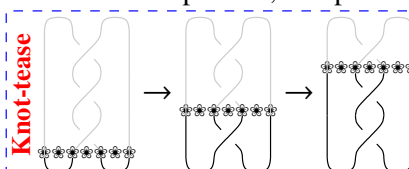
Gist of the Proof.



Why Tangles? • Faster!

- Conceptually clearer proofs of invariance (and of skein relations).
- Often fun and consequential:
 - The Jones Polynomial \rightsquigarrow The Temperley-Lieb Algebra.
 - Khovanov Homology \rightsquigarrow “Unfinished complexes”, complexes in a category.
 - The Kontsevich Integral \rightsquigarrow Associators.
 - HFK \rightsquigarrow OMG, type D , type A , $\mathcal{A}_\infty, \dots$

$$2^{n/2} + 2^{n/2} + 2^{\sqrt{n}} \ll 2^n$$



Computing Zombians of Unfinished Columbaria.

- Must be no slower than for finished ones.
- Future zombies must be able to complete the computation.
- Future zombies must not even know the size of the task that today's zombies were facing.
- We must be able to extend to ZPUCs, Zombie Processed Unfinished Columbaria!

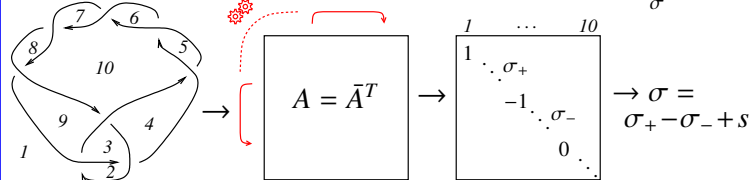


Columbarium near Assen

Example / Exercise. Compute the determinant of a $1,000 \times 1,000$ matrix in which 50 entries are not yet given.

Homework / Research Projects. • What with ZPUCs? • Use this to get an Alexander tangle invariant.

Reminders. {knots} \rightleftharpoons {matrices / quadratic forms} $\xrightarrow{\text{signature } \sigma} \mathbb{Z}$:



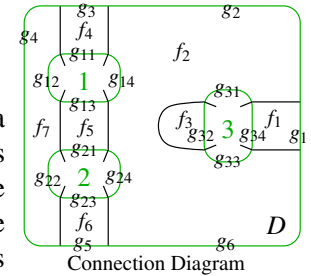
With $|\omega| = 1, t = 1 - \omega, r = t + \bar{t}, v = \text{Re}(\omega),$ and $u = \text{Re}(\omega^{1/2})$:

$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} -r & -t & 2t & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ 2\bar{t} & t & -r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A += \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$\bar{X}_{-i,j,k,-l}$	$A += \begin{pmatrix} r & -t & -2\bar{t} & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ -2t & t & r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A -= \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$

Theorem 2. ψ^* and ϕ_* are functorial. Also, if $\alpha // \beta = \gamma // \delta$, α is surjective, β is injective, and $\text{im } \gamma \supset \ker \delta$, then $\gamma^* // \alpha_* = \delta_* // \beta^*$. Finally, ψ^* is additive but ϕ_* isn't.

Definition. $S \left(\begin{matrix} g_2 \\ g_3 \\ \dots \end{matrix} \right) := \left\{ \begin{matrix} \text{SPQ } S \\ \text{on } \langle g_i \rangle \end{matrix} \right\}$.

Theorem 3. $\{S(\text{cyclic sets})\}$ is a planar algebra, with compositions $S(D)((S_i)) := \phi_*^D(\psi_D^*(\bigoplus_i S_i))$, where $\psi_D: \langle f_i \rangle \rightarrow \langle g_{\alpha i} \rangle$ maps every face of D to the sum of the input gaps adjacent to it and $\phi^D: \langle f_i \rangle \rightarrow \langle g_i \rangle$ maps every face to the sum of the output gaps adjacent to it. So for our D , ψ_D is $f_1 \mapsto g_{34}, f_2 \mapsto g_{31} + g_{14} + g_{24} + g_{33}, f_3 \mapsto g_{32}, f_4 \mapsto g_{11}, f_5 \mapsto g_{13} + g_{21}, f_6 \mapsto g_{23}, f_7 \mapsto g_{12} + g_{22}$ and ϕ^D is $f_1 \mapsto g_1, f_2 \mapsto g_2 + g_6, f_3 \mapsto 0, f_4 \mapsto g_3, f_5 \mapsto 0, f_6 \mapsto g_5, f_7 \mapsto g_4$.



Theorem 4. TL and Kas, defined on X and \bar{X} as before, extend to planar algebra morphisms {tangles} \rightarrow {S}.



Implementation (sources: <http://drorbn.net/icerm23/ap>). I like it most when the implementation matches the math perfectly. We failed here.

Once[<< KnotTheory`];

Loading KnotTheory` version

of February 2, 2020, 10:53:45.2097.

Read more at <http://katlas.org/wiki/KnotTheory>.

Utilities. The step function, algebraic numbers, canonical forms.

$\theta[x_]$ /; NumericQ[x] := UnitStep[x]

```
 $\omega 2[v\_][p\_]$  := Module[{q = Expand[p], n, c},
  If[q == 0, 0,
    c = Coefficient[q,  $\omega$ , n = Exponent[q,  $\omega$ ]];
     $c v^n + \omega 2[v][q - c(\omega + \omega^{-1})^n]$ ];
```

```
sign[ $\mathcal{E}$ _] := Module[{n, d, v, p, rs, e, k},
  {n, d} = NumeratorDenominator[ $\mathcal{E}$ ];
  {n, d} /=  $\omega^{\text{Exponent}[n, \omega]/2 + \text{Exponent}[n, \omega, \text{Min}]/2}$ ;
  p = Factor[ $\omega 2[v]@n * \omega 2[v]@d / . v \rightarrow 4 u^2 - 2$ ];
  rs = Solve[p == 0, u, Reals];
  If[rs == {}, Sign[p /. u -> 0],
    rs = Union@{u /. rs};
    Sign[(-1)e=Exponent[p, u] Coefficient[p, u, e]] + Sum[
      k = 0;
      While[{d = RootReduce[ $\partial_{\{u, ++k\}} p / . u \rightarrow r$ ]} == 0];
      If[EvenQ[k], 0, 2 Sign[d]] *  $\theta[u - r]$ ,
      {r, rs}]]
  ]
]
```

SetAttributes[B, Orderless];

```
CF[b_B] := RotateLeft[#, First@Ordering[#] - 1] & /@
DeleteCases[b, {}]
```

```
CF[ $\mathcal{E}$ _] := Module[{ $\gamma$ s = Union@Cases[ $\mathcal{E}$ ,  $\gamma$ _ |  $\bar{\gamma}$ _,  $\infty$ ]},
  Total[CoefficientRules[ $\mathcal{E}$ ,  $\gamma$ s] /.
    (ps_ -> c_) := Factor[c]  $\times$  Times@@ $\gamma$ sps]]
```

CF[{}] = {};

CF[C_List] :=

```
Module[{ $\gamma$ s = Union@Cases[C,  $\gamma$ _,  $\infty$ ],  $\gamma$ },
  CF /@ DeleteCases[0] [
    RowReduce[Table[ $\partial_{\gamma} r$ , {r, C}, { $\gamma$ ,  $\gamma$ s}]] .  $\gamma$ s ]
```

(\mathcal{E} _)* := $\mathcal{E} / . \{\bar{\gamma} \rightarrow \gamma, \gamma \rightarrow \bar{\gamma}, \omega \rightarrow \omega^{-1}, c_Complex \rightarrow c^*\}$;

r_Rule* := {r, r*}

RulesOf[γ_i + rest_] := ($\gamma_i \rightarrow -rest$)*;

```
CF[PQ[C_, q_]] := Module[{nC = CF[C]},
  PQ[nC, CF[q /. Union@@RulesOf /@nC]] ]
```

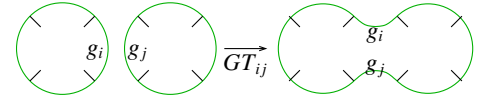
CF[Σ_b [σ _, pq_]] := $\Sigma_{CF[b]}$ [σ , CF[pq]]

Pretty-Printing.

```
Format[ $\Sigma_{b,B}[\sigma_, PQ[C_, q_]]$ ] := Module[{ $\gamma$ s},
   $\gamma$ s =  $\gamma$ # & /@ Join@@b;
  Column[{TraditionalForm@ $\sigma$ ,
    TableForm[Join[
      Prepend[""] /@ Table[TraditionalForm[ $\partial_c r$ ],
        {r, C}, {c,  $\gamma$ s}],
      {Prepend[""] [
        Join@@
          (b /. {L_, m___, r_} =>
            {DisplayForm@RowBox[{"(", L}],
              m, DisplayForm@RowBox[{r, ")"}]}) / .
            i_Integer =>  $\gamma_i$  ]},
      MapThread[Prepend,
        {Table[TraditionalForm[ $\partial_{r,c} q$ ], {r,  $\gamma$ s*},
          {c,  $\gamma$ s}],  $\gamma$ s*}]]
    ], TableAlignments -> Center]
  ], Center] ];
```

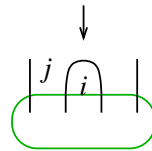
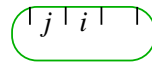
The Face-Centric Core.

```
 $\Sigma_{b1}[\sigma_1, PQ[C1_, q1_]] \oplus \Sigma_{b2}[\sigma_2, PQ[C2_, q2_]] \wedge :=$ 
  CF@ $\Sigma_{\text{Join}[b1, b2]}$ [ $\sigma_1 + \sigma_2, PQ[C1 \cup C2, q1 + q2]$ ];
```



GT for Gap Touch:

```
GTi,j@ $\Sigma_B[\{li___, i, ri___\}, \{lj___, j, rj___\}, bs___]$ [ $\sigma$ _,
  PQ[C_, q_]] :=
  CF@ $\Sigma_B[\{ri, li, j, rj, lj, i\}, bs]$ [ $\sigma$ _, PQ[C  $\cup$  { $\gamma_i - \gamma_j$ }, q]]
```



cor·don (kôr'dn)

n.

1. A line of people, military posts, or ships stationed around an area to enclose or guard it: a *police cordon*.
2. A rope, line, tape, or similar border stretched around an area, usually by the police, indicating that access is restricted.



$$s \begin{pmatrix} 0 & \phi C_{\text{rest}} \\ \bar{\phi}^T & \lambda \theta \\ \bar{C}_{\text{rest}}^T & \bar{\theta}^T A_{\text{rest}} \end{pmatrix} \rightarrow \begin{cases} \exists p \phi_p \neq 0 & \text{use } \phi_p \text{ to kill its row and} \\ \phi = 0, \lambda \neq 0 & \text{column, drop a } \begin{pmatrix} 01 \\ 10 \end{pmatrix} \text{ summand} \\ \phi = 0, \lambda = 0 & \text{use } \lambda \text{ to kill } \theta, \text{ let } s += \text{sign}(\lambda) \\ & \text{append } \theta \text{ to } C_{\text{rest}}. \end{cases}$$

```
Cordoni@ $\Sigma_B[\{li___, i, ri___\}, bs___]$ [ $\sigma$ _, PQ[C_, q_]] :=
```

```
Module[{ $\phi = \partial_{\gamma_i} C$ ,  $\lambda = \partial_{\bar{\gamma}_i, \gamma_i} q$ , n $\sigma = \sigma$ , nC, nq, p},
  {p} = FirstPosition[ (# != 0) & /@  $\phi$ , True, {0}];
  {nC, nq} = Which[
    p > 0, {C, q} /. ( $\gamma_i \rightarrow -C[[p]] / \phi[[p]]$ )* /. ( $\gamma_i \rightarrow \theta$ )*,
     $\lambda \neq 0$ , (n $\sigma += \text{sign}[\lambda]$ ;
      {C, q /. ( $\gamma_i \rightarrow -(\partial_{\bar{\gamma}_i} q) / \lambda$ )* /. ( $\gamma_i \rightarrow \theta$ )*}),
     $\lambda == 0$ , {C  $\cup$  { $\partial_{\bar{\gamma}_i} q$ }, q /. ( $\gamma_i \rightarrow \theta$ )*}];
  CF@ $\Sigma_B[\text{Most}@\{ri, li\}, bs]$ [n $\sigma$ ,
    PQ[nC, nq] /. ( $\gamma_{\text{Last}@\{ri, li\}} \rightarrow \gamma_{\text{First}@\{ri, li\}}$ )* ] ]
```

Strand Operations. c for contract, mc for magnetic contract:

$$C_{i,j}@t : \Sigma_B[\{li_ , i, ri_ \}, \{ _, j, _ \}, _] [_] := t // GT_{j, First\{ri, li\}} // Cordon_j$$

$$C_{i,j}@t : \Sigma_B[\{ _, i, j, _ \}, _] [_] := Cordon_j @ t$$

$$C_{i,j}@t : \Sigma_B[\{j, _, i, _ \}, _] [_] := Cordon_j @ t$$

$$C_{i,j}@t : \Sigma_B[\{ _, j, i, _ \}, _] [_] := Cordon_i @ t$$

$$C_{i,j}@t : \Sigma_B[\{i, _, j, _ \}, _] [_] := Cordon_i @ t$$

$$mc[\mathcal{E}_] := \mathcal{E} // .$$

$$t : \Sigma_B[\{ _, i, _ \}, \{ _, j, _ \}, _] [_] | \Sigma_B[\{ _, i, j, _ \}, _] [_] | \Sigma_B[\{j, _, i, _ \}, _] [_] / ; i + j == \theta \Rightarrow C_{i,j}@t$$

The Crossings (and empty strands).

$$Kas@P_{i,j} := CF@ \Sigma_B[\{i,j\}] [\theta, PQ[\{\}, \theta]] ;$$

$$TL@P_{i,j} := CF@ \Sigma_B[\{i,j\}] [\theta, PQ[\{\}, \theta]]$$

$$Kas[x : X[i, j, k, l]] :=$$

$$Kas@If[PositiveQ[x], X_{-i,j,k,-l}, \bar{X}_{-j,k,l,-i}] ;$$

$$Kas[(x : X | \bar{X})_{fs_}] := Module[\{v = 2u^2 - 1, p, \gamma s, m\},$$

$$\gamma s = \gamma_{\#} \& /@ \{fs\}; p = (x == X);$$

$$m = If[p, \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}, -\begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}];$$

$$CF@ \Sigma_B[\{fs\}] [If[p, -1, 1], PQ[\{\}, \gamma s^* . m . \gamma s]]]$$

$$TL[x : X[i, j, k, l]] :=$$

$$TL@If[PositiveQ[x], X_{-i,j,k,-l}, \bar{X}_{-j,k,l,-i}] ;$$

$$TL[(x : X | \bar{X})_{fs_}] := Module[\{t = 1 - \omega, r, \gamma s, m\},$$

$$r = t + t^*; \gamma s = \gamma_{\#} \& /@ \{fs\};$$

$$m = If[x == X,$$

$$\begin{pmatrix} -r & -t & 2t & t^* \\ -t^* & \theta & t^* & \theta \\ 2t^* & t & -r & -t^* \\ t & \theta & -t & \theta \end{pmatrix}, \begin{pmatrix} r & -t & -2t^* & t^* \\ -t^* & \theta & t^* & \theta \\ -2t & t & r & -t^* \\ t & \theta & -t & \theta \end{pmatrix}];$$

$$CF@ \Sigma_B[\{fs\}] [\theta, PQ[\{\}, \gamma s^* . m . \gamma s]]]$$

Evaluation on Tangles and Knots.

$$Kas[K_] := Fold[mc[\#1 \oplus \#2] \&, \Sigma_B[] [\theta, PQ[\{\}, \theta]], List@@ (Kas /@ PD@K)] ;$$

$$KasSig[K_] := Expand[Kas[K][[1]] / 2]$$

$$TL[K_] :=$$

$$Fold[mc[\#1 \oplus \#2] \&, \Sigma_B[] [\theta, PQ[\{\}, \theta]], List@@ (TL /@ PD@K)] / .$$

$$\theta[c_ + u] / ; Abs[c] \geq 1 \Rightarrow \theta[c] ;$$

$$TLSig[K_] := TL[K][[1]]$$

Reidemeister 3.

$$R3L = PD[X_{-2,5,4,-1}, X_{-3,7,6,-5}]$$

$$X_{-6,9,8,-4} ;$$

$$R3R = PD[X_{-3,5,4,-2}, X_{-4,6,8,-1}]$$

$$X_{-5,7,9,-6} ;$$

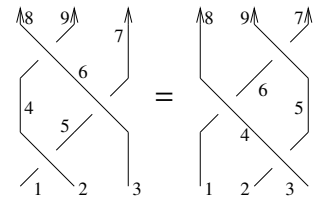
$$\{TL@R3L == TL@R3R, Kas@R3L == Kas@R3R\}$$

$$\{True, True\}$$

Kas@R3L

$$2\theta(u - \frac{1}{2}) - 2\theta(u + \frac{1}{2}) - 2$$

	γ_3	γ_7	γ_9	γ_8	γ_{-1}	γ_{-2}
$\bar{\gamma}_3$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$
$\bar{\gamma}_7$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2(2u^2-1)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$
$\bar{\gamma}_9$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$
$\bar{\gamma}_8$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$
$\bar{\gamma}_{-1}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2(2u^2-1)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$
$\bar{\gamma}_{-2}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$



Reidemeister 2.

$$TL@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}]$$

$$\begin{matrix} & \theta & & & \\ & 1 & \theta & -1 & \theta \\ (\gamma_{-2} & & \gamma_6 & \gamma_5 & \gamma_{-1}) \\ \bar{\gamma}_{-2} & \theta & \theta & \theta & \theta \\ \bar{\gamma}_6 & \theta & \theta & \theta & \theta \\ \bar{\gamma}_5 & \theta & \theta & \theta & \theta \\ \bar{\gamma}_{-1} & \theta & \theta & \theta & \theta \end{matrix}$$

$$\{TL@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}] == GT_{5,-2}@TL@PD[P_{-1,5}, P_{-2,6}], Kas@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}] == GT_{5,-2}@Kas@PD[P_{-1,5}, P_{-2,6}]\}$$

$$\{True, True\}$$

Reidemeister 1.

$$\{TL@PD[X_{-3,3,2,-1}] == TL@P_{-1,2},$$

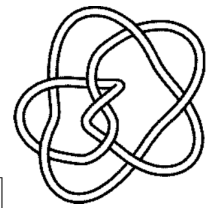
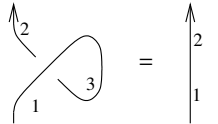
$$Kas@PD[X_{-3,3,2,-1}] == Kas@P_{-1,2}\}$$

$$\{True, True\}$$

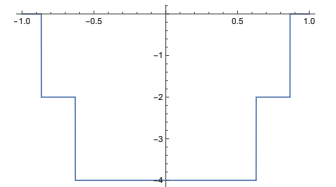
A Knot.

$$f = TLSig[Knot[8, 5]]$$

$$2\theta\left[-\frac{\sqrt{3}}{2} + u\right] - 2\theta\left[\frac{\sqrt{3}}{2} + u\right] - 2\theta\left[u - \left(\text{clockwise} - 0.630\dots\right)\right] + 2\theta\left[u - \left(\text{counterclockwise} 0.630\dots\right)\right]$$



$$Plot[f, \{u, -1, 1\}]$$



Proof of Theorem 1.

Uniqueness: If A and B are 2 pushforwards, then $\sigma_W(U + A) = \sigma_W(U + B)$ for all PQs U on W .

Thus $\mathcal{D}_A = \mathcal{D}_B$, because otherwise if $w \in \mathcal{D}_A \setminus \mathcal{D}_B$, by taking $U(w) = 1$ on $\mathcal{D}_U = \text{span}\{w\}$, we get $\sigma_W(U + A) = 1 \neq 0 = \sigma_W(U + B)$. Furthermore, A and B must agree where they are both defined, because by taking $U(w) = \frac{-A(w) - B(w)}{2}$ on $\mathcal{D}_U = \text{span}\{w\}$ we get $(U + A)(w) = \frac{A(w) - B(w)}{2} = -(U + B)(w)$, so we must have $A(w) = B(w)$ to satisfy $\sigma_W(U + A) = \sigma_W(U + B)$.

Existence: Define ϕ_*Q by $\mathcal{D}_{\phi_*Q} = \phi(\text{ann}_Q(\ker \phi))$ and $\phi_*Q(w) = Q(v)$ where $v \in \text{ann}_Q(\ker \phi)$. Note that ϕ_*Q is well-defined.

First consider when $U = 0$ on all of W . Let K be a maximal non-degenerate subspace of $\ker \phi$. Then $Q = Q|_K \oplus Q|_{\text{ann}_Q(K)}$, and we can write $\text{ann}_Q(K) = R \oplus A \oplus B$ where $R = \text{rad}_Q(\ker \phi)$ and A, B are chosen so that $A \subseteq \text{ann}_Q(R)$ and $B \subseteq \text{ann}_Q(K) \setminus \text{ann}_Q(R)$. Since $Q : R \rightarrow B^*$ is surjective, for any $v \in \mathcal{D}_Q$ there is some $r_v \in R$ such that $Q(r_v, B) = Q(v, B)$. If we choose the r_v so that $r_{v_1} + r_{v_2} = r_{v_1+v_2}$, then we can replace A by $A' = \{a - r_a : a \in A\}$ and B by $B' = \{b - \frac{1}{2}r_b : b \in B\}$ to get $Q = Q|_K \oplus Q|_{R \oplus B'} \oplus Q|_{A'}$. Then notice that

- $\sigma_V(Q|_K) = \sigma_{\ker \phi}(Q|_{\ker \phi})$
- $\sigma_V(Q|_{R \oplus B'}) = 0$
- $\sigma_V(Q|_{A'}) = \sigma_W(\phi_*Q)$

so we get $\sigma_V(Q) = \sigma_{\ker \phi}(Q|_{\ker \phi}) + \sigma_W(\phi_*Q)$.

Now for an arbitrary U , note that $(Q + \phi^*U)|_{\ker \phi} = Q|_{\ker \phi}$ and $\phi_*(Q + \phi^*U) = \phi_*Q + U$ so we can replace Q in the $U = 0$ case by $Q + \phi^*U$ to get the general case.

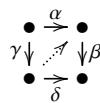
Proof of Theorem 2.

It's clear that pullback is functorial and that pushforward by the identity is the identity. To show $(\phi\psi)_* = \phi_*\psi_*$, use theorem 1 repeatedly to get

$$\begin{aligned} & \sigma((\phi\psi)_*Q + U) \\ &= \sigma(Q + (\phi\psi)^*U) \\ &= \sigma(Q + \psi^*\phi^*U) - \sigma(Q|_{\ker \phi\psi}) \\ &= \sigma(\psi_*Q + \phi^*U) + \sigma(Q|_{\ker \psi}) - \sigma(Q|_{\ker \phi\psi}) \\ &= \sigma(\phi_*\psi_*Q + U) + \sigma(Q|_{\ker \psi}) + \sigma(\psi_*Q|_{\ker \phi}) - \sigma(Q|_{\ker \phi\psi}) \\ &= \sigma(\phi_*\psi_*Q + U) \end{aligned}$$

for any U , where the last step uses theorem 1 on $Q|_{\phi\psi}$ with the map $\psi : \ker \phi\psi \rightarrow \ker \phi$.

To show $\alpha_*\gamma^* = \beta^*\delta_*$, first note that $\beta^*\beta_*$ is the identity on any PQ since β is injective, so



$$\alpha_*\gamma^*Q = \beta^*(\beta\alpha)_*\gamma^*Q = \beta^*(\delta\gamma)_*\gamma^*Q = \beta^*\delta_*\gamma_*\gamma^*Q$$

As $\beta^*\delta_*\gamma_*\gamma^*Q$ and $\beta^*\delta_*Q$ have the same values where they are both defined, it remains to show that they have the same domain. Since α is surjective and γ is surjective onto $\ker(\delta)$, we see that

$$\beta^{-1}\delta(A) = \beta^{-1}\delta(A \cap \text{im } \gamma)$$

for any subspace A . By taking $A = \text{ann}_Q(\ker \delta)$, the two sides of the equality become the domains of $\beta^*\delta_*Q$ and $\beta^*\delta_*\gamma_*\gamma^*Q$.

Proof of Thm 1: Fix W and consider triples $(V, Q, \phi: V \rightarrow W)$, where $Q = (S, D, \varphi)$ is an SPQ on V . Declare $(V_1, Q_1, \phi_1) \sim (V_2, Q_2, \phi_2)$ if for every quadratic U on W , $\sigma_{V_1}(Q_1 + \phi_1^*U) = \sigma_{V_2}(Q_2 + \phi_2^*U)$.

We need to show, given our (V, Q, ϕ) , 1. \exists SPQ Q' on W s.t.

$$(V, Q, \phi) \sim (W, Q', \mathbb{I})$$

2. If $(W, Q', \mathbb{I}) \sim (W, Q'', \mathbb{I})$ then $Q' = Q''$.

Property 2 easy. Property 1 follows from

Claim 1 $(V, Q, \phi) \sim (D(Q), Q, \phi|_{D(Q)})$, so wlog, Q is "full"; $D(Q) = V$.

Claim 2 If Q is full, $\forall v \in \ker \phi$, and $Q(v) \neq 0$, then $(V, Q, \phi) \sim$

$$(V/\langle v \rangle, \text{sign}(Q(v)) + Q + \frac{Q(-,v) \otimes Q(v,-)}{1|Q(v)|^2}, \phi|_{V/\langle v \rangle})$$

so wlog, $Q|_{\ker \phi} = 0$.

Claim 3 If Q is full and $Q|_{\ker \phi} = 0$, then $(V, Q, \phi) \sim (W, \phi|_{\ker Q(v,-)}, \mathbb{I})$