

Pensieve header: The “Speedy” engine, from pensieve://Projects/SL2Portfolio2/.

Program

## The “Speedy” Engine

Program

### Internal Utilities

Program

Canonical Form:

Program

```
In[ ]:= CCF[ $\mathcal{E}$ _] := PP_CCF@ExpandDenominator@ExpandNumerator@PP_Together@Together[PP_Exp[
Expand[ $\mathcal{E}$ ] /. e^x_ e^y_ -> e^{x+y} /. e^x_ -> e^{CCF[x]}]];
CF[ $\mathcal{E}$ _List] := CF /@  $\mathcal{E}$ ;
CF[ $sd\_SeriesData$ ] := MapAt[CF,  $sd$ , 3];
CF[ $\mathcal{E}$ _] := PP_CF@Module[
{vs = Cases[ $\mathcal{E}$ , (y | b | t | a | x |  $\eta$  |  $\beta$  |  $\tau$  |  $\alpha$  |  $\xi$ )_,  $\infty$ ] U {y, b, t, a, x,  $\eta$ ,  $\beta$ ,  $\tau$ ,  $\alpha$ ,  $\xi$ },
Total[CoefficientRules[Expand[ $\mathcal{E}$ ], vs] /. (ps_ -> c_) -> CCF[c] (Times@@vs^{ps})]
];
CF[ $\mathcal{E}\_E$ ] := CF /@  $\mathcal{E}$ ; CF[E_sp___[ $\mathcal{E}S\_\_\_\_$ ]] := CF /@ E_sp[ $\mathcal{E}S$ ];
```

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The Kronecker  $\delta$ :

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```
In[ ]:= K $\delta$  /: K $\delta$ _{i,j} := If[i === j, 1, 0];
```

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Equality, multiplication, and degree-adjustment of perturbed Gaussians;  $E[L, Q, P]$  stands for  $e^{L+Q} P$ :

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```
In[ ]:= E /: E[L1_, Q1_, P1_]  $\equiv$  E[L2_, Q2_, P2_] :=
CF[L1 == L2]  $\wedge$  CF[Q1 == Q2]  $\wedge$  CF[Normal[P1 - P2] == 0];
E /: E[L1_, Q1_, P1_] E[L2_, Q2_, P2_] := E[L1 + L2, Q1 + Q2, P1 * P2];
E[L_, Q_, P_]_{ $k$ } := E[L, Q, Series[Normal@P, { $\epsilon$ , 0,  $k$ }]];
```

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### Zip and Bind

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Variables and their duals:

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```
In[ ]:= {t*, b*, y*, a*, x*, z*} = { $\tau$ ,  $\beta$ ,  $\eta$ ,  $\alpha$ ,  $\xi$ ,  $\zeta$ };
{ $\tau^*$ ,  $\beta^*$ ,  $\eta^*$ ,  $\alpha^*$ ,  $\xi^*$ ,  $\zeta^*$ } = {t, b, y, a, x, z}; (u_{-i})^* := (u^*)_i;
```

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```
In[*]:= U21 = {B_{i-}^{p-} -> e^{-p h \gamma b_i}, B_{-}^{p-} -> e^{-p h \gamma b}, T_{i-}^{p-} -> e^{p h t_i}, T_{-}^{p-} -> e^{p h t}, \mathcal{A}_{i-}^{p-} -> e^{p \gamma \alpha_i}, \mathcal{A}_{-}^{p-} -> e^{p \gamma \alpha}};
12U = {e^{c-} b_{i+d-} -> B_{i-}^{c/(h \gamma)} e^d, e^{c-} b+d- -> B^{-c/(h \gamma)} e^d,
e^{c-} t_{i+d-} -> T_{i-}^{c/h} e^d, e^{c-} t+d- -> T^{c/h} e^d,
e^{c-} \alpha_{i+d-} -> \mathcal{A}_{i-}^{c/\gamma} e^d, e^{c-} \alpha+d- -> \mathcal{A}^{c/\gamma} e^d,
e^{\beta-} -> e^{Expand@\beta}}
```

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```
In[*]:= D_b[f_] := \partial_b f - h \gamma B \partial_B f; D_{b_i}[f_] := \partial_{b_i} f - h \gamma B_i \partial_{B_i} f;
D_t[f_] := \partial_t f + h T \partial_T f; D_{t_i}[f_] := \partial_{t_i} f + h T_i \partial_{T_i} f;
D_\alpha[f_] := \partial_\alpha f + \gamma \mathcal{A} \partial_{\mathcal{A}} f; D_{\alpha_i}[f_] := \partial_{\alpha_i} f + \gamma \mathcal{A}_i \partial_{\mathcal{A}_i} f;
D_v[f_] := \partial_v f; D_{\{v, \theta\}}[f_] := f; D_{\{\}}[f_] := f; D_{\{v, n\_Integer\}}[f_] := D_v[D_{\{v, n-1\}}[f]];
D_{\{L\_List, Ls\_ \_\_\}}[f_] := D_{\{Ls\}}[D_L[f]];
```

Program

Finite Zips:

Program

```
In[*]:= collect[sd_SeriesData, \xi_] := MapAt[collect[#, \xi] &, sd, 3];
collect[\xi_, \xi_] := PPCollect@Collect[\xi, \xi];
Zip_{\{\}}[P_] := P;
Zip_{\xi\xi}[Ps_List] := Zip_{\xi\xi} /@Ps;
Zip_{\{\xi, \xis\_ \_\_\}}[P_] := PPZip[
(collect[P // Zip_{\{\xi\xi\}}, \xi] /. f_{-} . \xi^{d-} -> (D_{\{\xi^*, d\}}[f])) /. \xi^* -> \theta /.
((\xi^* /. {b -> B, t -> T, \alpha -> \mathcal{A}}) -> 1)]
```

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QZip implements the “Q-level zips” on  $E(L, Q, P) = e^{L+Q} P(\epsilon)$ . Such zips regard the  $L$  variables as scalars.

$$\left\langle P(z_i, \zeta^j) e^{c + \eta^i z_i + y_j \zeta^j + q^i z_i \zeta^j} \right\rangle = |\tilde{q}| \left\langle P(z_i, \zeta^j) e^{c + \eta^i z_i} \Big|_{z_i \rightarrow \tilde{q}_i^k(z_k + y_k)} \right\rangle$$

$$= |\tilde{q}| e^{c + \eta^i \tilde{q}_i^k y_k} \left\langle P(\tilde{q}_i^k(z_k + y_k), \zeta^j + \eta^i \tilde{q}_i^j) \right\rangle.$$

Program

```
QZip_{\xi\xi\_List}@E[L_, Q_, P_] := PPQZip@Module[{ \xi, z, zs, c, ys, \eta s, qt, zrule, \xi rule, out},
zs = Table[\xi^*, {\xi, \xi s}];
c = CF[Q /. Alternatives @@ (\xi s \cup zs) -> \theta];
ys = CF@Table[\partial_{\xi} (Q /. Alternatives @@ zs -> \theta), {\xi, \xi s}];
\eta s = CF@Table[\partial_z (Q /. Alternatives @@ \xi s -> \theta), {z, zs}];
qt = CF@Inverse@Table[K\delta_{z, \xi^*} - \partial_{z, \xi} Q, {\xi, \xi s}, {z, zs}];
zrule = Thread[zs -> CF[qt . (zs + ys)]];
\xi rule = Thread[\xi s -> \xi s + \eta s . qt];
CF /@ E[L, c + \eta s . qt . ys, Det[qt] Zip_{\xi\xi}[P /. (zrule \cup \xi rule)]]];
```

Program

Upper to lower and lower to Upper:

Program

```
In[*]:=
U21 = {B_{i-}^{p-} \to e^{-p \hbar \gamma b_i}, B_{-}^{p-} \to e^{-p \hbar \gamma b}, T_{i-}^{p-} \to e^{p \hbar t_i}, T_{-}^{p-} \to e^{p \hbar t}, \mathcal{A}_{i-}^{p-} \to e^{p \gamma \alpha_i}, \mathcal{A}_{-}^{p-} \to e^{p \gamma \alpha}};
12U = {e^{c_{-} \cdot b_i + d_{-}} \to B_{i-}^{c/(h \gamma)} e^d, e^{c_{-} \cdot b + d_{-}} \to B_{-}^{c/(h \gamma)} e^d,
e^{c_{-} \cdot t_i + d_{-}} \to T_{i-}^{c/h} e^d, e^{c_{-} \cdot t + d_{-}} \to T_{-}^{c/h} e^d,
e^{c_{-} \cdot \alpha_i + d_{-}} \to \mathcal{A}_{i-}^{c/\gamma} e^d, e^{c_{-} \cdot \alpha + d_{-}} \to \mathcal{A}_{-}^{c/\gamma} e^d,
e^{\beta_{-}} \to e^{\text{Expand}[\beta]}};
```

Program

LZip implements the “L-level zips” on  $\mathbb{E}(L, Q, P) = P e^{L+Q}$ . Such zips regard all of  $P e^Q$  as a single “P”. Here the z’s are  $b$  and  $\alpha$  and the  $\zeta$ ’s are  $\beta$  and  $a$ .

Program

```
LZip_{\zeta S\_List}@E[L_, Q_, P_] :=
PP_{LZip}@Module[{z, zs, Zs, c, ys, \eta s, lt, zrulerule, Zrulerule, \zeta rulerule, Q1, EEQ, EQ},
zs = Table[\zeta^*, {\zeta, \zeta S}];
Zs = zs /. {b \to B, t \to T, \alpha \to \mathcal{A}};
c = L /. Alternatives@@(\zeta S \cup zs) \to 0;
ys = Table[\partial_{\zeta} (L /. Alternatives@@zs \to 0), {\zeta, \zeta S}];
\eta s = Table[\partial_z (L /. Alternatives@@\zeta S \to 0), {z, zs}];
lt = Inverse@Table[K\delta_{z, \zeta^*} - \partial_{z, \zeta} L, {\zeta, \zeta S}, {z, zs}];
zrulerule = Thread[zs \to lt.(zs + ys)];
Zrulerule = Join[zrulerule,
zrulerule /. r\_Rule \to ((U = r[[1]] /. {b \to B, t \to T, \alpha \to \mathcal{A}}) \to (U /. U21 /. r // 12U))];
\zeta rulerule = Thread[\zeta S \to \zeta S + \eta s.lt];
Q1 = Q /. (Zrulerule \cup \zeta rulerule);
EEQ[ps___] := EEQ[ps] = PP^{EEQ}@
(CF[e^{-Q1} D_{Thread[{zs, {ps}]}][e^{Q1}]] /. {Alternatives@@zs \to 0, Alternatives@@Zs \to 1});
CF@E[c + \eta s.lt.y s, Q1 /. {Alternatives@@zs \to 0, Alternatives@@Zs \to 1},
Det[lt] (Zip_{\zeta S}[(EQ@@zs) (P /. (Zrulerule \cup \zeta rulerule))] /.
Derivative[ps___][EQ][___] \to EEQ[ps] /. _EQ \to 1) ]];
```

Program

```
In[*]:=
B_{i} [L_, R_] := L R;
B_{is} [L_{E}, R_{E}] := PP_B@Module[{n},
Times[
L /. Table[(v : b | B | t | T | a | x | y)_i \to v_{n@i}, {i, {is}}],
R /. Table[(v : \beta | \tau | \alpha | \mathcal{A} | \zeta | \eta)_i \to v_{n@i}, {i, {is}}]
] // LZJoin@Table[{\beta_{n@i}, \tau_{n@i}, a_{n@i}}, {i, {is}}] // QZipJoin@Table[{\zeta_{n@i}, y_{n@i}}, {i, {is}}];
B_{is} [L_, R_] := B_{is} [L, R];
```

Program

## E morphisms with domain and range.

Program

```
In[ ]:=
BisList[Ed1→r1[L1_, Q1_, P1_], Ed2→r2[L2_, Q2_, P2_]] :=
  E(d1∪Complement[d2, is])→(r2∪Complement[r1, is]) @@ Bis[E[L1, Q1, P1], E[L2, Q2, P2]];
Ed1→r1[L1_, Q1_, P1_] // Ed2→r2[L2_, Q2_, P2_] :=
  Br1∩d2[Ed1→r1[L1, Q1, P1], Ed2→r2[L2, Q2, P2]];
Ed1→r1[L1_, Q1_, P1_] ≡ Ed2→r2[L2_, Q2_, P2_] ^:=
  (d1 == d2) ∧ (r1 == r2) ∧ (E[L1, Q1, P1] ≡ E[L2, Q2, P2]);
Ed1→r1[L1_, Q1_, P1_] Ed2→r2[L2_, Q2_, P2_] ^:=
  E(d1∪d2)→(r1∪r2) @@ (E[L1, Q1, P1] E[L2, Q2, P2]);
Ed→r[L_, Q_, P_] $k_ := Ed→r @@ E[L, Q, P] $k;
E[_E___][i_] := {E}[[i]];
```

Program

## “Define” Code

Program

Define[lhs = rhs, ...] defines the lhs to be rhs, except that rhs is computed only once for each value of \$k. Fancy Mathematica not for the faint of heart. Most readers should ignore.

Program

```
In[ ]:=
SetAttributes[Define, HoldAll];
Define[def_, defs__] := (Define[def]; Define[defs]);
Define[op_is = E_] := Module[{SD, ii, jj, kk, isp, nis, nisp, sis}, Block[{i, j, k},
  ReleaseHold[Hold[
    SD[opnisp, $k_Integer, PPBoot@Block[{i, j, k}, opisp, $k = E; opnis, $k]];
    SD[opisp, op{is}, $k]; SD[opsis, op{sis}];
  ] /. {SD → SetDelayed,
    isp → {is} /. {i → i_, j → j_, k → k_},
    nis → {is} /. {i → ii, j → jj, k → kk},
    nisp → {is} /. {i → ii_, j → jj_, k → kk_}
  } ] ]]
```