

- If $\mu_i = (\lambda_i; \omega_i) \in M(T_i; H_i)$ for $i = 1, 2$ (and, of course, $T_1 \cap T_2 = \emptyset = H_1 \cap H_2$), set

$$\mu_1 * \mu_2 := (\lambda_1 * \lambda_2; \iota_1(\omega_1) + \iota_2(\omega_2)),$$

where ι_i are the obvious inclusions $\iota_i: CW(T_i) \rightarrow CW(T_1 \cup T_2)$.

- The only truly new definition is that of tha^{ux} :

$$(\lambda; \omega) \parallel tha^{ux} := (\lambda; \omega + J_u(\lambda_x)) \parallel RC_u^{\lambda_x}.$$

Thus the “new” tha^{ux} is just the “old” tha^{ux} , with an added term of $J_u(\lambda_x)$.

Theorem 4.2. *M, with the operations defined above, is a meta-group-action (MGA). Furthermore, if $\zeta: \mathcal{K}_0^{bh} \rightarrow M$ is defined on the generators in the same way as ζ_0 , except extended by 0 to the CW factor,*

$$\zeta(\epsilon_u^t) := ((); 0), \quad \zeta(\epsilon_x^h) := ((x \rightarrow 0); 0), \quad \text{and} \quad \zeta(\rho_{ux}^\pm) := ((x \rightarrow \pm u); 0),$$

then it is well-defined; namely, the values above satisfy the relations in Definition [2.5](#).

Proof. MORE.

Thus we have a tree-and-wheel valued invariant ζ defined on \mathcal{K}_0^{bh} , and thus $\delta \parallel \zeta$ is a tree-and-wheel valued invariant of tangles and w-tangles.

5. SOME COMPUTATIONAL EXAMPLES

Part of the reason I am happy about the invariant ζ is that it is relatively easily computable. Cyclic words are easy to implement, and using the Lyndon basis (e.g. [Re, Chapter 5]), free Lie algebras are easy too. Hence I include here a demo-run of a rough implementation, written in *Mathematica*. The full source files are available at [KBH].

First we load the package `FreeLie.m`, which contains a collection of programs to manipulate series in a completed Lie algebra and series of cyclic words. We tell `FreeLie.m` to show series by default only up to degree 3, and that if two (infinite) series are compared, they are to be compared by default only up to degree 5:

```
<< FreeLie.m
$SeriesShowDegree = 3; $SeriesCompareDegree = 5;
```

Merely as a test of `FreeLie.m`, we tell it to set `t1` to be `bch(u, v)`. The computer's response is to print that series to degree 3:

```
t1 = BCH[⟨u⟩, ⟨v⟩]
LS[u + v,  $\frac{uv}{2}$ ,  $\frac{1}{12} \overline{u \overline{u v}} + \frac{1}{12} \overline{u \overline{v u}}$ ]
```

Note that by default Lie series are printed in “top bracket form”, which means that brackets are printed above their arguments, rather than around them. Hence $\overline{u \overline{u v}}$ means $[u, [u, v]]$. This practice is especially advantageous when it is used on highly-nested elements, when it becomes difficult for the eye to match left brackets with their corresponding right brackets.

Note also that that `FreeLie.m` utilizes *lazy evaluation*, meaning that when a Lie series (or a series of cyclic words) is defined, its definition is stored but no computation takes place until it is printed or until its value (at a certain degree) is explicitly requested. Hence `t1` is a reference to the entire Lie series `bch(u, v)`, and not merely to the degrees 1-3 parts of

that series, which are printed above. Hence when we request the value of `t1` at degree 6, the computer complies:

```
t1@6 // TopBracketForm
- $\frac{uuu\overline{uvv}}{1440}$  +  $\frac{1}{360}$   $\overline{uu\overline{uvv}v}$  +  $\frac{1}{240}$   $\overline{u\overline{uv}\overline{uvv}}$  +  $\frac{1}{720}$   $\overline{u\overline{u\overline{v}}\overline{uv}}$  -  $\frac{u\overline{uvv}vv}{1440}$ 
```

The package `FreeLie.m` know about various free Lie algebra operations, but not about our specific circumstances. Hence we have to make some further definitions. The first few are set-theoretic in nature. We define the “domain” of a function stored as a list of *key*→*value* pairs to be the set of “first elements” of these pair; meaning, the set of keys. We define what it means to remove a key (and its corresponding value), and likewise for a list of keys. We define what it means for two functions to be equal (their domains must be equal, and for every key #, we are to have # // $f_1 = \# // f_2$). We also define how to apply a Lie morphism `mor` to a function (apply it to each value), and how to compare (λ, ω) pairs (in $FL(T)^H \times CW(T)$):

```
Domain[f_List] := First /@ f;
f \ key_ := DeleteCases[f, key -> _];
f \ keys_List := Fold[#1 \ #2 &, f, keys];
f1_List == f2_List := Domain[f1] == Domain[f2] && (And @@ (
  ((# /. f1) == (# /. f2)) & /@ Domain[f1]
));
LieMorphism[mor_][f_List] := MapAt[LieMorphism[mor], f, {All, 2}];
M[\lambda_1, \omega_1] == M[\lambda_2, \omega_2] := (\lambda1 == \lambda2) && (\omega1 == \omega2);
```

Next we enter some free-Lie definitions that are not a part of `FreeLie.m`. Namely we define $R_{u,\bar{u}}^{\lambda_x}(s)$ to be the result of “stable application” of the morphism $u \rightarrow e^{\text{ad}(\lambda_x)}(\bar{u})$ to s (namely, apply the morphism repeatedly until things stop changing; at any fixed degree this happens after a finite number of iterations). We define $R_u^{\lambda_x}$ to be $R_{u,\bar{u}}^{\lambda_x} // (\bar{u} \rightarrow u)$. Finally, we define J as in Equation (13):

```
RC[u_, \lambda_x LieSeries, ub_][s_] :=
  StableApply[LieMorphism[<u> -> Ad[\lambda_x][<ub>]], s];
RC[u_, \lambda_x LieSeries][s_] := s // RC[u, \lambda_x, <u>] // LieMorphism[<u> -> <u>];
J[u_, \lambda_x] := Module[{s},
  IntegrateCWSeries[
    div[u, \lambda_x // RC[u, s \lambda_x]] // LieMorphism[u -> Ad[-s \lambda_x][u]],
    {s, 0, 1}];
```

Next is a series of definitions that implement the definitions of $*$, tm , hm , and tha following Sections 3.2 and 4.2:

```

M := M[λ1, ω1] ∪ M[λ2, ω2] := M[λ1 ∪ λ2, ω1 + ω2];
tm[u_, v_, w_][λ_List] := λ // LieMorphism[⟨u⟩ → ⟨w⟩, ⟨v⟩ → ⟨w⟩];
tm[u_, v_, w_][M[λ_, ω_]] := LieMorphism[⟨u⟩ → ⟨w⟩, ⟨v⟩ → ⟨w⟩] /@ M[λ, ω];
hm[x_, y_, z_][λ_List] := Union[λ \ {x, y}, {z → BCH[x /. λ, y /. λ]}];
hm[x_, y_, z_][M[λ_, ω_]] := M[λ // hm[x, y, z], ω];
tha[u_, x_][λ_List] := MapAt[RC[u, x /. λ], λ, {All, 2}];
tha[u_, x_][M[λ_, ω_]] :=
  M[λ // tha[u, x], (ω + J[u, x /. λ]) // RC[u, x /. λ]];
ρ+[u_, x_] := M[{x → MakeLieSeries[⟨u⟩]}, MakeCWSeries[0]];
ρ-[u_, x_] := M[{x → MakeLieSeries[-⟨u⟩]}, MakeCWSeries[0]];
Print /@ {{u = ⟨"u"⟩, v = ⟨"v"⟩, w = ⟨"w"⟩}};
1 → (t1 = M[{
  x → MakeLieSeries[u + v + w],
  y → MakeLieSeries[b[u, v] + b[v, w]]
}, MakeCWSeries[CW["uvw"]]]),
2 → (t1 // tm[u, v, u]),
3 → (t2 = t1 // tm[u, v, u] // tm[u, w, u]),
4 → (t1 // tm[v, w, v]),
5 → (t3 = t1 // tm[v, w, v] // tm[u, v, u]),
6 → (t2 ≡ t3)
];

1 → M[{x → LS[u + v + w, 0, 0], y → LS[0, u v + v w, 0]}, CWS[0, 0, CW[uvw]]]
2 → M[{x → LS[2 u + w, 0, 0], y → LS[0, u w, 0]}, CWS[0, 0, CW[uuw]]]
3 → M[{x → LS[3 u, 0, 0], y → LS[0, 0, 0]}, CWS[0, 0, CW[uuu]]]
4 → M[{x → LS[u + 2 v, 0, 0], y → LS[0, u v, 0]}, CWS[0, 0, CW[uvv]]]
5 → M[{x → LS[3 u, 0, 0], y → LS[0, 0, 0]}, CWS[0, 0, CW[uuu]]]
6 → True

```

commentary missing
same

Commentary

```
Print /@ {
  1 → (t1 = ρ*[u, x] ∪ ρ*[v, y] ∪ ρ*[w, z]),
  2 → (t1 // hm[x, y, x]),
  3 → (t2 = t1 // hm[x, y, x] // hm[x, z, x]),
  4 → (t1 // hm[y, z, y]),
  5 → (t3 = t1 // hm[y, z, y] // hm[x, y, x]),
  6 → (t2 ≡ t3)
};
```

add a wheels part

```
1 → M[{x → LS[u, 0, 0], y → LS[v, 0, 0], z → LS[w, 0, 0]}, CWS[0, 0, 0]]
2 → M[{x → LS[u + v,  $\frac{u\bar{v}}{2}$ ,  $\frac{1}{12} \overline{u\bar{u}\bar{v}} + \frac{1}{12} \overline{u\bar{v}\bar{v}}$ ], z → LS[w, 0, 0]}, CWS[0, 0, 0]]
3 → M[{x → LS[u + v + w,  $\frac{u\bar{v}}{2} + \frac{u\bar{w}}{2} + \frac{v\bar{w}}{2}$ ,  $\frac{1}{12} \overline{u\bar{u}\bar{v}} + \frac{1}{12} \overline{u\bar{u}\bar{w}} + \frac{1}{3} \overline{u\bar{v}\bar{w}} + \frac{1}{12} \overline{v\bar{v}\bar{w}} + \frac{1}{12} \overline{u\bar{v}\bar{v}} + \frac{1}{6} \overline{u\bar{w}\bar{v}} + \frac{1}{12} \overline{u\bar{w}\bar{w}} + \frac{1}{12} \overline{v\bar{w}\bar{w}}$ ], CWS[0, 0, 0]]
4 → M[{x → LS[u, 0, 0], y → LS[v + w,  $\frac{v\bar{w}}{2}$ ,  $\frac{1}{12} \overline{v\bar{v}\bar{w}} + \frac{1}{12} \overline{v\bar{w}\bar{w}}$ ], CWS[0, 0, 0]]
5 → M[{x → LS[u + v + w,  $\frac{u\bar{v}}{2} + \frac{u\bar{w}}{2} + \frac{v\bar{w}}{2}$ ,  $\frac{1}{12} \overline{u\bar{u}\bar{v}} + \frac{1}{12} \overline{u\bar{u}\bar{w}} + \frac{1}{3} \overline{u\bar{v}\bar{w}} + \frac{1}{12} \overline{v\bar{v}\bar{w}} + \frac{1}{12} \overline{u\bar{v}\bar{v}} + \frac{1}{6} \overline{u\bar{w}\bar{v}} + \frac{1}{12} \overline{u\bar{w}\bar{w}} + \frac{1}{12} \overline{v\bar{w}\bar{w}}$ ], CWS[0, 0, 0]]
6 → True
```

test
fraction,
h action

```
Print /@ {
  1 → (t1 = ρ*[u, x] ∪ ρ*[v, y] ∪ ρ*[w, z]),
  2 → (t2 = t1 // tm[v, w, v] // hm[x, y, x] // tha[u, z]),
  3 → (t3 = ρ*[v, x] ∪ ρ*[w, z] ∪ ρ*[u, y]),
  4 → (t4 = t3 // tm[v, w, v] // hm[x, y, x]),
  5 → (t2 ≡ t4)
};
```

```
1 → M[{x → LS[u, 0, 0], y → LS[v, 0, 0], z → LS[w, 0, 0]}, CWS[0, 0, 0]]
2 → M[{x → LS[u + v,  $-\frac{u\bar{v}}{2}$ ,  $\frac{1}{12} \overline{u\bar{u}\bar{v}} + \frac{1}{12} \overline{u\bar{v}\bar{v}}$ ], z → LS[v, 0, 0]}, CWS[0, 0, 0]]
3 → M[{x → LS[v, 0, 0], y → LS[u, 0, 0], z → LS[w, 0, 0]}, CWS[0, 0, 0]]
4 → M[{x → LS[u + v,  $-\frac{u\bar{v}}{2}$ ,  $\frac{1}{12} \overline{u\bar{u}\bar{v}} + \frac{1}{12} \overline{u\bar{v}\bar{v}}$ ], z → LS[v, 0, 0]}, CWS[0, 0, 0]]
5 → True
```

817
Barromeny

6. THE RELATION WITH THE BF TOPOLOGICAL QUANTUM FIELD THEORY

7. THE SIMPLEST NON-COMMUTATIVE REDUCTION AND AN ULTIMATE ALEXANDER INVARIANT

8. THE RELATION WITH ALEKSEEV-TOROSSIEN AND WITH ^{WKO}[BND]

9. ODDS AND ENDS

ubsec:Hopf

9.1. **Linking Numbers and Signs.** If x is an oriented S^1 and u is an oriented S^2 in an oriented S^4 (or \mathbb{R}^4) and the two are disjoint, their linking number l_{ux} is defined as follows. Pick a ball B whose oriented boundary is u (using the "outward pointing normal" convention