

MAT 1100 Core Algebra.  
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To do. 1. Print "About"  
2. Print NCGE. (two sides)

on board Goal: Within your lifetime, understand  $G = \langle g_1 \dots g_m \rangle \subset S_n$ :  
1.  $|G| = ?$  2.  $\sigma \in G?$  3.  $\sigma = w(g_1, \dots, g_m)$  4. random

Two pre-requisites 1. Groups,  $S_n$ , silly uniquenesses, cancellation,  $(ab)^{-1} = b^{-1}a^{-1}$ , subgroups, the subgroup generated by  $\{\sigma_2\}$ .

2. Row reduction for real.

$$f \cdot g = f \circ g$$

Example  $\sigma_1 = (123)$   $\sigma_2 = (12)(34)$ , in  $S_4$   
2314 2143

11	I		
12	1	22	I
13	2	23	3
14	5	24	4
		33	I
		34	
		44	I

$\sigma_1 = 2314$   
 $\sigma_{12}^2 = 3124$   
 $\sigma_{12}^{-1} \sigma_2 = 1342$   
 $\sigma_{23} \sigma_{13} = 4132$   
 $\sigma_{13}^{-1} \sigma_{23} \sigma_{12} = 1423$

Feed  $\sigma_1 = 2314 \dots$  Feed @  $\sigma_{12}$

Feed  $\sigma_{12}^2 = 3124 \dots$  Feed @  $\sigma_{13}$

Feed  $\sigma_2 = 2143 \dots$  Feed  $\sigma_{12}^{-1} \sigma_2 = 1342 \dots$  Feed @  $\sigma_{23}$

Feed  $\sigma_{12} \sigma_{23} = 2143 \dots$  Feed  $\sigma_{12}^{-1} \sigma_{12} \sigma_{23} = \sigma_{23} \dots$

No point feeding  $\sigma_{ij} \sigma_{kl}$  if  $i=k$ !

Feed  $\sigma_{23} \sigma_{12} = 3412 \dots$  Feed  $\sigma_{13}^{-1} \sigma_{23} \sigma_{12} = 1423 \dots$  to  $\sigma_{24}$

Feed  $\sigma_{23} \sigma_{13} = 4132 \dots$  to  $\sigma_{14}$

Feed  $\sigma_{24} \sigma_{12} = 4213 \dots$  Feed  $\sigma_{14}^{-1} \sigma_{24} \sigma_{12} = 1423 \dots$  drop.

$\Rightarrow |G| = 4 \cdot 3 \cdot 1 \cdot 1 = 12$ . Is  $4123 \in G$ ?

Write  $2431$  in terms of  $\sigma_{i,j}$ .

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\* Go over the "about" handout.

$$\sigma_{8,j}(\sigma_{4,j_4} M_5) \stackrel{1}{=} (\sigma_{8,j} \sigma_{4,j_4}) M_5 \stackrel{2}{\subset} M_4 M_5$$

$$\stackrel{3}{=} \sigma_{4,j'_4} (M_5 M_5) \stackrel{4}{\subset} \sigma_{4,j'_4} M_5 \subset M_4$$

on board.

Read Along: Selick's notes 1.1, 1.2.1, 1.4; Lang's book I 1-3.

**Very quickly:** groups, uniqueness of  $1$ ,  $^{-1}$ ,  $(ab)^{-1} = b^{-1}a^{-1}$   
order of an element.

Group homomorphisms, "the category of groups"  
The group  $\text{Aut}(G)$

$$\text{Conjugation: } g^h = h^{-1}gh = C_h(g) \quad (g_1 g_2)^h = g_1^h g_2^h \quad g^{h_1 h_2} = (g^{h_1})^{h_2}$$

$h \mapsto C_h$  is an anti-homomorphism  $G \rightarrow \text{Aut}(G)$

Images, kernels, subgroups.

Example:  $S_3$  is an image of  $S_4$ , but not a kernel.

Normal subgroups, kernels are normal.

Question Is every normal subgroup the kernel of a homomorphism? Given  $N \trianglelefteq G$ , can we find a surjective homomorphism  $\phi: G \rightarrow H$ , with  $\ker \phi = N$ ?

Set Theoretic aside: Surjections are the same as equivalence relations.

(def'n, explanation ...)

Sol'n Suppose  $N$  had  $\phi$ , consider the resulting equiv:

done here.

Sol'n Suppose we had  $\phi$ , consider the resulting equiv:

$$g_1 \sim g_1 n \text{ or } g_1 \sim g_2 \text{ iff } g_1^{-1} g_2 \in N.$$

$$\text{Let } H = G/\sim = \{[g]\} \text{ where } [g] = gN$$

$$\text{with } \phi: G \rightarrow H \text{ being } \phi(g) = [g]$$

$$\text{Define } [g_1][g_2] = [g_1 g_2] \quad (\text{well defined!})$$

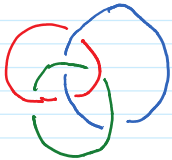
$$[g]^{-1} = [g^{-1}]$$

Claim  $H = G/\sim$  is a group &  $\phi$  is a morphism  
whose kernel is  $N$  ... we write  $H = G/N$ .

Theorem (The First Isomorphism Theorem) Given  
any morphism  $\phi: G \rightarrow H$ ,  $G/\ker \phi \cong \text{im } \phi$ .



Riddle Along.



Can you draw 4 linked loops, so that if you drop any one of them, the remaining 3 are not linked?

Read Along. Selick 1.1-1.4

Today's menu. Quotients and the isomorphism thms

Reminder: Given  $N \trianglelefteq G$  ( $\forall g \in G \quad N^g = g^{-1}Ng = N$ ), we seek  $\sim$  on  $G$  s.t.  $\phi: G \rightarrow G/\sim =: H$  will be a group homomorphism with  $\ker \phi = N$ .

$$\begin{aligned} g_1 \sim g_2 &\Leftrightarrow \phi(g_1) = \phi(g_2) \Leftrightarrow \phi(g_1 g_2^{-1}) = e \Leftrightarrow \\ &\Leftrightarrow g_1 g_2^{-1} \in N \Leftrightarrow g_1 \in g_2 N \Leftrightarrow g_1 N = g_2 N \end{aligned}$$

Let  $H = G/\sim = \{[g]\}$  where  $[g] = gN$   
with  $\phi: G \rightarrow H$  being  $\phi(g) = [g]$

$$\begin{aligned} \text{Define } [g_1][g_2] &= [g_1 g_2] \quad (\text{well defined!}) \\ [g]^{-1} &= [g^{-1}] \end{aligned}$$

Claim  $H = G/\sim$  is a group &  $\phi$  is a morphism whose kernel is  $N$  ... we write  $H = G/N$ .

Theorem (The First Isomorphism Theorem) Given any morphism  $\phi: G \rightarrow H$ ,  $G/\ker \phi \cong \text{im } \phi$ .

$$\begin{aligned} \text{pf construct } R: &\rightarrow \text{ by } [g] \mapsto \phi(g) \\ L: &\leftarrow \text{ by } h \mapsto [g] \text{ s.t. } \phi(g) = h. \end{aligned}$$

Aside  $G/H$  when  $H \leq G$  & Lagrange's thm.

Claim. For  $H, K \leq G$ ,  $HK \leq G$  iff  $HK = KH$ .

pf.  $\Leftarrow (h_1 k_1)(h_2 k_2) = h_1 h_2 k_2' k_2$

$\Rightarrow (hk)^{-1} = h'k' = k^{-1}h^{-1}$

Definition.  $C_G(X) := \{g \in G : \forall x \in X \quad g^{-1}xg = x\}$   
 $Z(G) := C_G(G)$   
 $N_G(X) := \{g \in G : g^{-1}Xg = X\}$  } all are subgroups

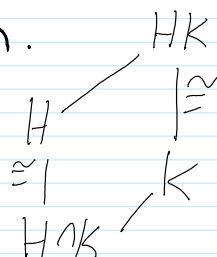
claim. IF  $H \leq N_G(K)$  then  $HK = KH$ ,  $K \trianglelefteq HK$ , &  $H \cap K \trianglelefteq H$ .

pf. trivial.

The 2<sup>nd</sup> Isomorphism Theorem.

IF  $H \leq N_G(K)$ , then

$HK/K \cong H/H \cap K$



pf.  $R: \rightarrow : [h]_K \rightarrow [h]_{H \cap K}$

$L: \leftarrow : \text{obvious.}$

The 3<sup>rd</sup> Isomorphism Thm.

IF  $K, H \trianglelefteq G$  &  $K \leq H$ , then

$\frac{G/K}{H/K} \cong G/H$

pf.  $R: \rightarrow : [[g]_K]_{H/K} \rightarrow [g]_H$

well defined?  $[[g]_K]_{H/K} = [[g_2]_K]_{H/K} \Rightarrow$

$\Rightarrow [g_1]_K [g_2]_K^{-1} = [h]_K \Rightarrow g_1 g_2^{-1} = hK = K$

The 4<sup>th</sup> Isomorphism Thm.

IF  $N \trianglelefteq G$  then  $\pi: G \rightarrow G/N$  induces a "faithful" bijection between subgroups of  $G/N$  and  $\{H : N \leq H \leq G\}$ :

\*  $A \leq B \Leftrightarrow \pi(A) \leq \pi(B)$  (& then,  $[B:A] = [\pi(B):\pi(A)]$ )

\*  $A \trianglelefteq B \Leftrightarrow \pi(A) \trianglelefteq \pi(B)$

\*  $\pi(A \cap B) = \pi(A) \cap \pi(B)$ .

Also did:  $\text{sign}(\sigma) = \text{sign}(\prod_{i < j} (\sigma(i) - \sigma(j)))$

Read Along. Pavel Etingof's "Groups Around Us", Lang's page 57.

Riddle Along. Your turn!

Today's Menu. Jordan-Holder.

Reminders.  $\phi: G \rightarrow H$ :

$$G/\ker \phi \cong \text{im } \phi$$

$$\dim V - \text{nullity } L = \text{rank } L$$

$$H < N_G(K):$$

$$HK/K \cong H/H \cap K$$

$$\underbrace{\hspace{2cm}}_{H} \underbrace{\hspace{2cm}}_{K} \dots$$

$$\frac{G/K}{H/K} \cong G/H$$

$$\frac{G}{N}: \left\{ \begin{array}{l} \text{subgroups} \\ 0 < G/N \end{array} \right\} \leftrightarrow \{H: N < H < G\}$$

Definition A simple group.

"A prime", (in fact,  $\mathbb{Z}/n$  is simple iff  $n$  is a prime)

$$\text{ykt } S_3 \triangleright A_3 = \langle (123) \rangle = \mathbb{Z}_3; S_3/A_3 = \mathbb{Z}_2$$

$$\mathbb{Z}_6 \triangleright 2\mathbb{Z}_6 = \langle 0, 2, 4 \rangle = \mathbb{Z}_3; \mathbb{Z}_6/2\mathbb{Z}_6 = \mathbb{Z}_2$$

The Jordan-Holder Theorem. Let  $G$  be a finite group. Then there exist a sequence

$$G = G_0 \triangleright G_1 \triangleright G_2 \triangleright \dots \triangleright G_n = \{e\} \text{ s.t. } H_i = G_i/G_{i-1}$$

is simple. Furthermore, the sequence  $(H_i)$ , the "composition series" of  $G$ , is unique up to a permutation.

$$\text{Example } S_4 \triangleright A_4 \triangleright \begin{pmatrix} (12)(34) \\ (13)(24) \\ (14)(23) \end{pmatrix} \triangleright \begin{pmatrix} (12)(34) \\ (13)(24) \end{pmatrix} \triangleright \{e\}$$

24      12      4      2      12      3

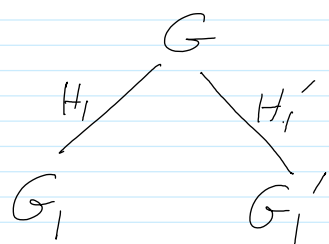
Proof by induction on  $|G|$ .

Existence: Let  $G_1$  be a maximal normal

on board

proper subgroup.

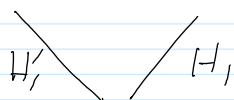
Uniqueness: Use the "Diamond principle":



$$G \supset G_1 \supset G_2 \dots$$

$$G \supset G_1' \supset G_2' \dots$$

claim  $G = G_1 G_1'$



pf  $G_1 G_1'$  is normal in  $G$  yet

bigger than  $G_1, G_1'$ .

done  
line

$$\text{sign}: S_n \rightarrow \{\pm 1\} \quad \text{by} \quad \text{sign}(\sigma) = (-1)^\sigma = \text{sign}\left(\prod_{i < j} (\sigma_i - \sigma_j)\right)$$

$$= \prod_{\{i,j\} \in \{1, \dots, n\}} S_{ij}(\sigma) \quad S_{ij}(\sigma) = \text{sign}\left(\frac{\sigma_i - \sigma_j}{i - j}\right)$$

$$(-1)^\sigma = \text{sign}\left[\prod_{\{i,j\}} \frac{\sigma_i - \sigma_j}{i - j}\right] = (-1)^\sigma (-1)^\tau$$

Every permutation is a product of transpositions,  
The parity is the parity of the number  
of transpositions.

**Theorem.**  $A_n$  is simple for  $n \neq 4$ . [Proof as in Lang's]

**Cycle Decomposition.**  $(12)(345) = [21453] = 21453$

claim If  $\sigma = (a_1 \dots a_k)$  and  $\tau = [\tau_1 \tau_2 \dots \tau_n]$ ,

then

$$\sigma^\tau = \tau^{-1} \sigma \tau = (\tau^{-1} a_1, \tau^{-1} a_2, \dots)$$

Corollary  $\sigma$  is conjugate to  $\sigma'$  iff they have  
the same cycle lengths

Corollary  $\#(\text{Conjugacy classes of } S_n) = P(n)$

**Lemma 1.** Every element of  $A_n$  is a product of 3-cycles.

PF  $(12)(23) = (123), (123)(234) = (12)(34) - \dots$

**Lemma 2.** IF  $N \triangleleft A_n$  contains a 3-cycle, then  $N = A_n$

PF WLOG,  $(123) \in N$ . claim For  $\sigma \in S_n$ ,  $(123)^\sigma \in N$   $\left( \begin{smallmatrix} \sigma \in A_n \checkmark \\ \sigma = (12)\sigma \checkmark \end{smallmatrix} \right)$

So  $N$  contains all 3-cycles...  $\square$

Now take  $N \triangleleft A_n$  w/  $N \neq \{1\}$

**Case 1.**  $N$  contains an element w/ cycle of length  $\geq 4$

$$\sigma = (123456)\sigma' \in N \quad \sigma^{-1}(123)\sigma(123)^{-1} = (136)$$

**Case 2.**  $N$  contains an element  $\sigma = (123)(456)\sigma'$

$$\text{consider } \sigma^{-1}(124)\sigma(124)^{-1} = (14263)$$

**Case 3.**  $N$  contains  $\sigma = (123)$  (product of pairs)

$$\text{Then } \sigma^2 = (132) - \dots$$

**Case 4.** Every element of  $N$  is a product of disjoint 2-cycles

$$\sigma = (12)(34)\sigma' \Rightarrow \sigma^{-1}(123)\sigma(123)^{-1} = (13)(24) = \tau \in N$$

$$\Rightarrow \tau^{-1}(125)\tau(125)^{-1} = (13452) \in N$$

HW1 is out!

Riddle Along. 1. Can you find uncountably many nearly-disjoint  
 $[\forall \alpha, \beta, |A_\alpha \cap A_\beta| < \infty]$  subsets of  $\mathbb{N}^{\mathbb{N}}$ ?

2. Can you find an uncountable chain  $[\forall \alpha, \beta, (A_\alpha \subset A_\beta) \vee (A_\beta \subset A_\alpha)]$   
 of subsets of  $\mathbb{N}^{\mathbb{N}}$ ?

Today's Menu. Simplicity of  $A_n$ , group actions.

Reminder.  $\text{sign}: S_n \rightarrow \{\pm 1\}$  by  $\text{sign}(\sigma) = (-1)^{\bar{\sigma}} = \text{sign}(\prod_{i < j} (\sigma_i - \sigma_j))$

$$= \prod_{\{i,j\} \subset \{1, \dots, n\}} S_{ij}(\sigma) \quad S_{ij}(\sigma) = \text{sign}\left(\frac{\sigma_i - \sigma_j}{i - j}\right)$$

$$(-1)^{\bar{\sigma}} = \text{sign}\left[\prod_{\{i,j\}} \frac{\sigma_i - \sigma_j}{i - j}\right] = (-1)^{\bar{\sigma}} (-1)^{\bar{\tau}}$$

Every permutation is a product of transpositions,  
 The parity is the parity of the number  
 of transpositions.

Theorem.  $A_n$  is simple for  $n \neq 4$ .

Cycle Decomposition.  $(12)(345) = [21453] = 21453$

Claim If  $\sigma = (a_1 \dots a_k)$  and  $\tau = [\tau_1 \tau_2 \dots \tau_n]$ ,

then

$$\sigma^\tau = \tau^{-1} \sigma \tau = (\tau^{-1}(a_1), \tau^{-1}(a_2), \dots)$$

Corollary  $\sigma$  is conjugate to  $\sigma'$  iff they have  
 the same cycle lengths

Corollary  $\#(\text{conjugacy classes of } S_n) = P(n)$

Now follow handout....

Jordan-Hölder for  $S_n$ :  $S_n \triangleright A_n \triangleright \{e\}$  ( $n \geq 5$ )

Definition A  $G$ -set (left- $G$ -set)  $G \times X \rightarrow X$

s.t.  $(g_1 g_2)x = g_1(g_2 x)$ ,  $e x = x$ . Same as  $\alpha: G \rightarrow S(X)$ .

$G$ -sets are a category!

Examples. 0.  $G$  itself, under mult. on the left.

1.  $G$  itself, under conjugation.

2.  $\text{Subgroups}(G)$ , under conjugation.

Examples: 1.  $G/H$  when  $H$  is not-necessarily normal

sub-example:  $S_n/S_{n-1}$   $\sigma S_{n-1} = \sigma' S_{n-1}$  iff

$\sigma(n) = \sigma'(n)$ . Let  $\tau_i(n) = i$ , then

$\sigma \tau_i S_{n-1} = \tau_{\sigma(i)} S_{n-1}$ . So  $S_n/S_{n-1}$  is  $\{1, \dots, n\}$ ...

2. If  $X_1, X_2$  are  $G$ -sets, then so is  $X_1 \sqcup X_2$ .

3.  $S^2 = \text{SO}(3)/\text{SO}(2)$

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Theorem. 1. Every  $G$ -set is a disjoint union of "transitive  $G$ -sets"

2. If  $X$  is a transitive  $G$  set and  $x \in X$ , then  $X \cong G/\text{stab}_x(x)$ . (So  $|X| \mid |G|$ )

Theorem. If  $X$  is a  $G$  set and  $x_i$  are representatives of the orbits, then

$$|X| = \sum_i \frac{|G|}{|\text{stab}_x(x_i)|}$$

Example. If  $G$  is a  $p$ -group, the centre of  $G$  is more than  $\{e\}$ .

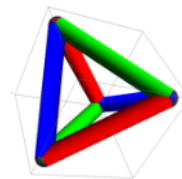
Suggestion for a good deed:  
TeX this up nicely!

## The Simplicity of the Alternating Groups

This handout is to be read twice: first read **red** only, to ascertain that everything in **red** is easy and boring, then read black and **red**, to actually understand the proof.

**Theorem.** The alternating group  $A_n \trianglelefteq S_n$  is simple for  $n \neq 4$ .

**Remark.** Easy for  $n \leq 3$ , false for  $n=4$  as there is  $\phi: A_4 \twoheadrightarrow A_3$ , so assume  $n \geq 5$ .



**Lemma 1.** Every element of  $A_n$  is a product of 3-cycles.  
**PF.** Every  $\sigma \in A_n$  is a product of an even number of 2-cycles, and  $(12)(23) = (123) \ \& \ (123)(234) = (12)(34)$ .

**Lemma 2.** If  $N \triangleleft A_n$  contains a 3-cycle, then  $N = A_n$ .  
**PF.** WLOG,  $(123) \in N$ . Then for all  $\sigma \in S_n$ ,  $(123)^\sigma \in N$ : if  $\sigma \in A_n$ , this is clear. Otherwise  $\sigma = (12)\sigma'$  w/  $\sigma' \in A_n$ , and then as  $(123)^{(12)} = (123)^2$ ,  $(123)^\sigma = ((123)^2)^{\sigma'} \in N$ . So  $N$  contains all 3-cycles.

**Case 1.**  $N$  contains an element w/ cycle of length  $\geq 4$ .

**Resolution.**  $\sigma = (123456)\sigma' \in N \Rightarrow \sigma^{-1}(123)\sigma(123)^{-1} = (136) \in N$

**Case 2.**  $N$  contains an element w/ 2 cycles of length 3.

**Res.**  $\sigma = (123)(456)\sigma' \in N \Rightarrow \sigma^{-1}(124)\sigma(124)^{-1} = (14263) \in N$ .

**Case 3.**  $N$  contains  $\sigma = (123) \cdot (\text{a product of disjoint 2-cycles})$ .

**Res.**  $\sigma^2 = (132) \in N$

**Case 4.** Every element of  $N$  is product of disjoint 2-cycles.

**Res.**  $\sigma = (12)(34)\sigma' \Rightarrow \sigma^{-1}(123)\sigma(123)^{-1} = (13)(24) = \tau \in N$   
 $\Rightarrow \tau^{-1}(125)\tau(125)^{-1} = (13452) \in N$

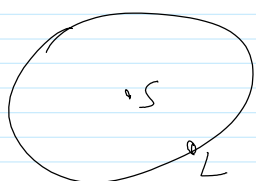




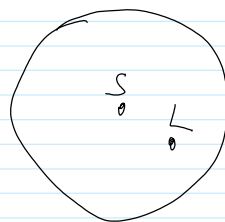
class photo at end!

HW1 is out!

Riddle Along.



$$V_L = 4V_S$$



$$V_S = V_L$$

Today's Menu. Group actions.

Reminder.  $G \curvearrowright X$ ,  $X \curvearrowright G$ , both are categories!  $G/H$  start line

Theorem. 1. Every  $G$ -set is a disjoint union of "transitive  $G$ -sets"

2. If  $X$  is a transitive  $G$  set and  $x \in X$ , then  $X \cong G/\text{stab}_x(x)$ . (So  $|X| \mid |G|$ )

Theorem. If  $X$  is a  $G$  set and  $x_i$  are representatives of the orbits, then

$$|X| = \sum_i \frac{|G|}{|\text{stab}_x(x_i)|}$$

The class equation:

the centre of  $G$

the centralizer of  $y_i$  in  $G$

$$|G| = |Z(G)| + \sum_i (G : C_G(y_i))$$

Where  $\{y_i\}$  are representatives from the non-central conjugacy classes of  $G$ .

Example. If  $G$  is a  $p$ -group, the centre of  $G$  is more than  $\{e\}$ .

done  
lim

is more than  $\{e\}$ .

## THE SYLOW THEOREMS.

Lovely notation:  $p^\alpha \parallel |G|$

$|G| = p^\alpha m$ ,  $p$  prime,  $p \nmid m$ ;  $\text{Syl}_p(G) := \{P \leq G : |P| = p^\alpha\}$  are "Sylow  $p$ -subgroups of  $G$ ". A " $p$ -subgroup" in general, is any subgroup of  $G$  of order a power of  $p$ .

Sylow 1  $\text{Syl}_p(G) \neq \emptyset$ .

Proof. By induction on  $|G|$ , if  $G$  has a normal subgroup of order  $p$  (or  $p^k$ ) or if  $G$  has a subgroup of order divisible by  $p^\alpha$ , we are done. The existence of one of the said types follows from the class equation:

$$|G| = |Z(G)| + \sum_i (G : C_G(y_i))$$

} Either both are divisible by  $p$ , or neither. Do 2<sup>nd</sup> case first.

Where  $\{y_i\}$  are representatives from the non-central conjugacy classes of  $G$ . □

Theorem. If  $G$  is a finite Abelian group of order divisible by a prime  $p$ , then  $G$  contains an element of order  $p$ . "Cauchy's Thm" DLF pp 102

Proof. Enough to find an element of order divisible by  $p$ ; if  $z$  is of order  $p \cdot n$ ,  $z^n$  would be of order  $p$ . Pick  $x \in G$ ,  $x \neq 1$ . If  $p \mid |x|$ , we're done. Otherwise  $p \nmid |x|$ , so by induction,  $\exists y \in \langle x \rangle$  s.t.

$|y| = p$  in  $G/\langle x \rangle$ . Now use the following claim.  $\square$

claim. if  $\phi: G \rightarrow H$  is a morphism &  $y \in G$ ,  
then  $|\phi(y)| \mid |y|$ .

Proof. If  $|\phi(y)| = n$ ,  $|y| = m$ ,  $m = nq + r$ , then  
 $e = \phi(y^m) = \phi(y^{nq})\phi(y^r) = (\phi(y)^n)^q \phi(y)^r = \phi(y)^r$   
so  $r = 0$ .

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Stronger Sylow 1. If  $p^B \mid |G|$ , then  $G$   
has a subgroup of order  $p^B$ .

Proof. Let  $X = \{ \underset{\text{subset}}{S} \subseteq G : |S| = p^B \}$ , and write

$|G| = p^{\alpha+B} m$  w/ maximal  $\alpha$ . By counting  
& binomial nonsense,  $p^\alpha \mid |X|$  yet  $p^{\alpha+1} \nmid |X|$ .

$G$  acts on  $X$  by translations, so there must  
be  $S_0 \in X$  s.t.  $p^{\alpha+1} \nmid |G \cdot S_0|$ , hence

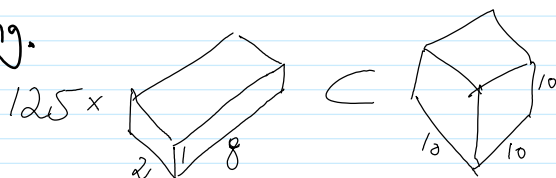
$p^B \mid |H = \text{stab}_G(S_0)|$ . Yet if  $x \in S_0$  then

$g \mapsto gx$  is an injection  $H \rightarrow S_0$ , so

$|H| \leq |S_0| = p^B$ , so  $|H| = p^B$ .

class photo on web!

Riddle Along.



Today's Menu. Sylow 1 2 3, some classification.  
Reminders.

$$G \curvearrowright X \Rightarrow |X| = \sum_i \frac{|G|}{|\text{Stab}_G(x_i)|}$$

$$|G| = |Z(G)| + \sum_i (G : C_G(y_i))$$

$G$  a  $p$ -group  $\Rightarrow Z(G)$  non-trivial

## THE SYLOW THEOREMS. Lovely notation: $p^x \parallel |G|$

$|G| = p^x m$ ,  $p$  prime,  $p \nmid m$ ;  $\text{Syl}_p(G) := \{P \leq G : |P| = p^x\}$   
are "Sylow  $p$ -subgroups of  $G$ ". A " $p$ -subgroup" in general, is any subgroup of  $G$  of order a power of  $p$ .

Sylow 1  $\text{Syl}_p(G) \neq \emptyset$ .

Proof. By induction on  $|G|$ , if  $G$  has a normal subgroup of order  $p$  (or  $p^x$ ) or if  $G$  has a subgroup of order divisible by  $p^x$ , we are done. The existence of one of the said types follows from the class equation:

$$|G| = |Z(G)| + \sum_i (G : C_G(y_i))$$

} Either both are divisible by  $p$ ,  
or neither.  
Do 2nd case first

$$|G| = |Z(G)| + \sum_i (G \cdot \langle G(y_i) \rangle) \quad \text{or 1st case. Do 2nd case first.}$$

Where  $\{y_i\}$  are representatives from the non-central conjugacy classes of  $G$ .  $\square$

Theorem. If  $G$  is a finite Abelian group of order divisible by a prime  $p$ , then  $G$  contains an element of order  $p$ . "Cauchy's Thm" D&F pp 102

Proof. Enough to find an element of order divisible by  $p$ , if  $z$  is of order  $p \cdot n$ ,  $z^n$  would be of order  $p$ . Pick  $x \in G, x \neq 1$ . If  $p \mid |x|$ , we're done. Otherwise  $p \nmid |x|$ , so by induction,  $\exists y \in G$  s.t.

$|y| = p$  in  $G/\langle x \rangle$ . Now use the following claim.  $\square$

claim. if  $\phi: G \rightarrow H$  is a morphism &  $y \in G$ , then  $|\phi(y)| \mid |y|$ .

Proof. If  $|\phi(y)| = n, |y| = m, m = nq + r$ , then

$$e = \phi(y^m) = \phi(y^{nq})\phi(y^r) = (\phi(y)^n)^q \phi(y)^r = \phi(y)^r$$

So  $r = 0$ .

Stronger Sylow 1. If  $p^B \mid |G|$ , then  $G$  has a subgroup of order  $p^B$ .

Proof. Let  $X = \{S \subseteq G : |S| = p^B\}$ , and write  
subset

$|G| = p^{\alpha+B} m$  w/ maximal  $\alpha$ . By counting & binomial nonsense,  $p^{\alpha} \mid |X|$  yet  $p^{\alpha+1} \nmid |X|$ .

$G$  acts on  $X$  by translations, so there must be  $s_0 \in X$  s.t.  $p^{\alpha+1} \nmid |G \cdot s_0|$ , hence  $p^B \mid |H = \text{stab}_G(s_0)|$ . Yet if  $x \in s_0$  then  $g \mapsto gx$  is an injection  $H \rightarrow s_0$ , so  $|H| \leq |s_0| = p^B$ , so  $|H| = p^B$ .

- Theorem.**
1. Sylow  $p$ -groups always exist;  $\text{Syl}_p(G) \neq \emptyset$ .
  2. Every  $p$ -group is contained in a Sylow- $p$  group.
  3. All Sylow- $p$  subgroups of  $G$  are conjugate, and  $n_p(G) := |\text{Syl}_p(G)| \equiv 1 \pmod p$  &  $n_p(G) \mid |G|$

### Groups of order 15.

$P_5$  is normal in  $G$ ,  $P_3$  is normal in  $G$ . Any  $y \in P_3$  commutes with  $P_5$  [otherwise,  $|y| \mid |\text{Aut } P_5| = 4$ ],

(Aside.  $\text{Aut}(\mathbb{Z}/p) = (\mathbb{Z}/p)^*$  so  $|\text{Aut}(\mathbb{Z}/p)| = p-1$ )

So  $G = x^i y^j = y^j x^i$  for generators  $x \in P_5, y \in P_3$ .

Aside. If  $G = G_1 \cdot G_2$ ,  $G_1 \cap G_2 = \langle e \rangle$ ,  $[G_1, G_2] = \langle e \rangle$ , then

$$G = G_1 \times G_2$$

Aside.  $\mathbb{Z}/p \times \mathbb{Z}/q = \mathbb{Z}_{pq}$

So  $G_{15} = \mathbb{Z}/15$ .

In fact, if  $(a, b) = 1$ , then  $\mathbb{Z}/a \times \mathbb{Z}/b \cong \mathbb{Z}/ab$

Proof. Find  $s, t$  s.t.  $as + bt = 1$ , and write

$$\begin{array}{ccccc} & & \cdot t & \rightarrow & \mathbb{Z}/a & \xrightarrow{\cdot b} & \\ & \nearrow & & & \searrow & & \\ \mathbb{Z}/ab & & & \times & & & \mathbb{Z}/ab \\ & \searrow & \cdot s & \rightarrow & \mathbb{Z}/b & \xrightarrow{\cdot a} & \end{array}$$

This also works for order  $p^2$ ,  $p < q$  primes,  $p \nmid q-1$ .

Groups of order 21.  $P_7$  is normal,  $P_3$  might not be  
 $P_3$  may act on  $P_7$ . IF  $P_7 = \langle x \rangle$ ,  $P_3 = \langle y \rangle$ , we have  $x^y = x$ , or  $x^2$ , or  $x^4$ . done line (but only 5k 21, not general p.q)

Def. What does this mean?

Aside.  $\text{Aut}(\mathbb{Z}/p)$  is cyclic;

$$\text{Aut}(\mathbb{Z}/7) = \langle x \mapsto x^3 \rangle$$

1 3 2 6 4 5

This also works for order  $p^2$ ,  $p < q$  primes,  $p \nmid q-1$ .

Riddle Along. Your turn again!

Today's Menu.  $G_{pq}$ , proofs of Sylow 2-3.

Reminders. **Theorem.** 1. Sylow  $p$ -groups always exist;  $\text{Syl}_p(G) \neq \emptyset$

2. Every  $p$ -group is contained in a Sylow- $p$  group.

3. All Sylow- $p$  subgroups of  $G$  are conjugate, and

$$n_p(G) := |\text{Syl}_p(G)| \equiv 1 \pmod{p} \quad \& \quad n_p(G) \mid |G|$$

Groups of order 15.  $P_3 \triangleleft G$ ,  $P_5 \triangleleft G$ ,  $G = P_3 \times P_5 = \mathbb{Z}/3 \times \mathbb{Z}/5 = \mathbb{Z}/15$

Groups of order 21.  $P_7 \triangleleft G$ ,  $P_3$  may not be normal

IF normal,  $G = P_3 \times P_7 = \mathbb{Z}/21$ .

Otherwise,  $P_7 = \langle x \rangle$ ,  $P_3 = \langle y \rangle$ ,  
we have  $x^y = x$ , or  $x^2$ , or  $x^4$ .

Dedt. What does this mean?

Aside.  $\text{Aut}(\mathbb{Z}/p)$  is cyclic;  
 $(\mathbb{Z}/p)^*$

$$\text{Aut}(\mathbb{Z}/7) = \langle x \mapsto x^3 \rangle$$

skip

Groups of order  $pq$ .  $n_p \mid pq \Rightarrow n_p \mid q$ , (or  $n_p = 1$ )

$$n_p \equiv 1 \pmod{p} \Rightarrow q \equiv 1 \pmod{p} \Rightarrow p \mid q-1$$

IF  $p < q$ ,  $p \nmid q-1 \Rightarrow G = \mathbb{Z}/pq$

if  $p \mid q-1$ , small may act on big ----.

The "extension lemma":

**Lemma.** 1. IF  $P \in \text{Syl}_p(G)$  &  $H < N_G(P)$  is a  $p$ -group,

then  $H \subset P$

2. IF  $P \in \text{Syl}_p(G)$ ,  $|x| = p^B$ ,  $x \in N_G(P)$ , then  $x \in P$ .

Reformulation:  $P \in \text{Syl}_p(G)$ ,  $|H| = p^B \Rightarrow N_H(P) = H \cap P$

**Proposition.** IF  $P \in \text{Syl}_p(G)$ , then  $|\text{conjugates of } P| \equiv 1 \pmod{p}$ .

**Proof.**  $P$  acts on the (and  $n_p \mid |G|$ , of course)

Set of its conjugates by conjugation. The orbit



$\{P\}$  is a singleton; by lemma, the sizes of all other orbits are divisible by  $p$ .

done line

**Proposition.** If  $H$  is a  $p$ -subgroup &  $P \in \text{Syl}_p(G)$ , then  $H$  is contained in a conjugate of  $P$ . [in particular, all Sylow- $p$  subgroups are conjugates]

**Proof.**  $H$  acts on the set of conjugates of

$P$  by conjugation. There must be a singleton orbit — a  $P'$  s.t.  $H \leq N_G(P')$ .

HW1 due!

Riddle Along.

$$\forall x \in \mathbb{R} \exists a_i \in \mathbb{Q} \text{ s.t. } a_i \rightarrow x$$

$$\mathbb{Q} \cap [-\infty, x] \text{ so what?}$$

Today's Menu. Finish Sylow, semi-direct products

Reminders. **Theorem.** 1. Sylow  $p$ -groups always exist;  $\text{Syl}_p(G) \neq \emptyset$ .

2. Every  $p$ -group is contained in a Sylow- $p$  group.

3. All Sylow- $p$  subgroups of  $G$  are conjugate, and

$$n_p(G) := |\text{Syl}_p(G)| \equiv 1 \pmod{p} \quad \& \quad n_p(G) \mid |G|$$

The extension trick: Can't extend a Sylow by something of order  $p$ .

**Proposition.** If  $P \in \text{Syl}_p(G)$ , then  $|\text{conjugates of } P| \equiv 1 \pmod{p}$ .  
(and  $n_p \mid |G|$ , of course)

**Proposition.** If  $H$  is a  $p$ -subgroup &  $P \in \text{Syl}_p(G)$ , then  $H$  is contained in a conjugate of  $P$ . In particular, all Sylow- $p$  subgroups are conjugate.

**Proof.**  $H$  acts on the set of conjugates of

$P$  by conjugation. There must be a singleton orbit — a  $P'$  s.t.  $H \leq N_G(P')$ .

**Semi-Direct Products.** If  $N \leq G$ ,  $H \leq G$ , compare  $N \rtimes H$  with  $NH$ .

There's always  $\mu: N \rtimes H \rightarrow NH$  by  $(n, h) \mapsto nh$ .

In general, nothing to say.

If  $N \cap H = \{e\}$ , injective but image might not be a group.

Example:  $\langle (123) \rangle, \langle (345) \rangle \subset S_5$

If  $N \cap H = \{e\}$  &  $N \trianglelefteq G$  &  $H \trianglelefteq G$ , then  $[N, H] = \{e\}$  &

$$NH \cong N \times H.$$

The interesting case is when  $N \cap H = \{e\}$ ,  $N \trianglelefteq G$ ,  $H$  <sup>maybe not</sup>.

Get  $H \hookrightarrow \text{Aut}(N)$  by  $h \mapsto (n \mapsto n^{h^{-1}} = h n h^{-1})$

$$\text{or } \phi_h(n) = h n h^{-1}$$

$$n_1 h_1 n_2 h_2 = n_1 h_1 n_2 h_1^{-1} h_1 h_2 = n_1 \phi_{h_1}(n_2) h_1 h_2$$

$$(nh)^{-1} = h^{-1} n^{-1} = h^{-1} n^{-1} h h^{-1} = \phi_{h^{-1}}(n^{-1}) \cdot h^{-1}$$

**Definition.** Given abstract  $N, H$  &  $\phi: H \rightarrow \text{Aut}(N)$ ,  
the semi-direct product  $N \rtimes H$ .

**Prop.** 1. In the above case,  $\mu: N \rtimes H \rightarrow NH$  is  
an isomorphism.

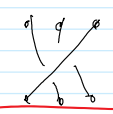
$$2. H < N \rtimes H, N \triangleleft (N \rtimes H) \text{ and } N \rtimes H / N \cong H.$$

**Small Examples.** 1.  $D_{2n} \cong \mathbb{Z}/n \rtimes \{\pm 1\}$

$$2. \{ax+b\} = \mathbb{R}_b^+ \rtimes \mathbb{R}_a^\times$$

$$3. \{Ax+b: A \in GL(V), b \in V\} = V_b \rtimes GL(V)_A$$

$$4. \text{"The Poincare/Relativity Group"} = \mathbb{R}^4 \rtimes SO(3,1)$$

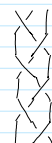
**Big Example.**  $B_n = \pi_1((\mathbb{C}^2 - \{\text{diag}\})/S_n) =$   I should have started the discussion of PB with an intro to free groups and w/  $\pi_1(\text{xxx}) = F_n$

$$B_n = \langle \sigma_1, \dots, \sigma_{n-1} : \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}, \sigma_i \sigma_j = \sigma_j \sigma_i \text{ if } |i-j| > 1 \rangle$$

an aside on free groups, generators & relations.

$$\pi: B_n \rightarrow S_n \quad PB_n = \ker \pi$$

$$PB_n \triangleleft B_n \text{ yet not } B_n = PB_n \rtimes S_n$$



Two reasons why I like this one:  
1. knotted  $\mathbb{R}^2$ 's  
2. Borromean.

$$\rho: PB_n \rightarrow PB_{n-1} \quad \ker \rho = F_{n-1} \text{ and}$$

$$PB_n = F_{n-1} \rtimes PB_{n-1} = F_{n-1} \rtimes (F_{n-2} \rtimes (\dots (F_2 \rtimes \mathbb{Z}) \dots))$$

**Groups of order 21.**  $\mathbb{Z}/21, \mathbb{Z}/7 \rtimes \mathbb{Z}/3 = \langle x \rangle \rtimes \langle y \rangle$

$$\text{Aut}(\mathbb{Z}/7) = \mathbb{Z}/6 = \langle \phi_3 \rangle; \phi_3(x) = x^3; x^y = x \text{ or } x^2 \text{ or } x^4$$

(iso: if  $x^y = x^2$  &  $y^2 = y^2$  then  $x^{\bar{y}} = x^4$ ) ↑ isomorphic

**Groups of order 12.** If  $|G| = 12$ ,  $P_4 = \mathbb{Z}/4$  or  $(\mathbb{Z}/2)^2$ ,  $P_3 = \mathbb{Z}/3$ ,

and at least one of  $P_4$  is normal, for there's not enough room for 4  $P_3$  & 3  $P_4$ 's. So  $G$  is a semi-direct

Product:  $\mathbb{Z}/4 \rtimes \mathbb{Z}/3$  : must be  $\mathbb{Z}/4 \rtimes \mathbb{Z}/3 = \mathbb{Z}/12$  ( $\text{Aut}(\mathbb{Z}/4) = \mathbb{Z}/2$ !)

$(\mathbb{Z}/2 \times \mathbb{Z}/2) \rtimes \mathbb{Z}/3$ : Either direct;  $\mathbb{Z}/2 \times \mathbb{Z}/6$

or the fun action of  $\mathbb{Z}/3$  on  $(\mathbb{Z}/2)^2$ , giving  $A_4$

$\langle (234) \rangle$

$\emptyset$   
 $(12)(34)$   
 $(13)(24)$   
 $(14)(23)$

$\mathbb{Z}/3 \rtimes (\mathbb{Z}/2 \times \mathbb{Z}/2)$ : Either direct or  $D_6 \rtimes \mathbb{Z}/2 = D_{12}$

$\mathbb{Z}/3 \rtimes \mathbb{Z}/4$ : Either direct or  $\mathbb{Z}/3 \rtimes \mathbb{Z}/4$

# Scratch 141005

October-05-14 8:10 PM

$$(12) \cdot (123)$$

$$9 / 120$$

$$(123)(345)$$

$$(12)(237)$$

$$6 / 24$$

HW2 discussion.

I should have added to HW2:  $G \rtimes G \cong G \times G$  conj. action

Aside,



Two reasons why I like this one:  
1. knotted \$20's  
2. Borromean.  
3. It is a commutator.

Q. Can you find a 4-component Brunnian link?

Today's Menu. semi-direct products, groups of order 12

Reminders. Given  $N, H, \phi: H \rightarrow \text{Aut}(N)$ ,

$$N \rtimes H := \langle nh \rangle; \quad n_1 h_1 \cdot n_2 h_2 = n_1 \phi_{h_1}(n_2) h_1 h_2$$

Thm 1.  $N \rtimes H$  is a group,  $H \leq N \rtimes H$ ,  $N \triangleleft N \rtimes H$ ,

$$N \cap H = \{e\} \quad (N \rtimes H / N = H)$$

2. In general, if  $G = NH$ ,  $N \triangleleft G$ ,  $H \leq G$ ,  $N \cap H = \{e\}$ ,

$$\text{Then } G \cong N \rtimes_{\phi} H \text{ w/ } \phi_h(n) = h n h^{-1}$$

$$PB_n := \pi_1(\mathbb{C}^n \setminus \text{diags}) = \text{"pure braids on } n \text{ strands"}$$

$$\rho: PB_n \rightarrow PB_{n-1} \quad \ker \rho = F_{n-1} \text{ and}$$

$$PB_n = F_{n-1} \rtimes PB_{n-1} = F_{n-1} \rtimes (F_{n-2} \rtimes (\dots (F_2 \rtimes \mathbb{Z}) \dots))$$

Aside,

$$B_n = \langle \sigma_1, \dots, \sigma_{n-1} : \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}, \sigma_i \sigma_j = \sigma_j \sigma_i \text{ if } |i-j| > 1 \rangle$$

an aside on free groups, generators & relations.

Groups of order 21.  $\mathbb{Z}/21$ ,  $\mathbb{Z}/7 \rtimes \mathbb{Z}/3 = \langle x \rangle \rtimes \langle y \rangle$

$$\text{Aut}(\mathbb{Z}/7) = \mathbb{Z}/6 = \langle \phi \rangle; \quad \phi(x) = x^3; \quad y \mapsto \phi^0 \text{ or } \phi^2 \text{ or } \phi^4$$

$$yxy^{-1} = \underbrace{x}_{\mathbb{Z}/7} \text{ or } \underbrace{x^2 \text{ or } x^4}_{\text{isomorphic to } \mathbb{Z}/21}$$

iso: if  $yxy^{-1} = x^2$  & then  $y^2xy^{-2} = x^4$ , so

$$G_2 = \langle x \rangle \rtimes_{\phi^2} \langle y \rangle \longrightarrow \langle \bar{x} \rangle \rtimes_{\phi^4} \langle \bar{y} \rangle = G_4$$

$$\begin{pmatrix} x \\ y^2 \end{pmatrix} \longmapsto \begin{pmatrix} \bar{x} \\ \bar{y} \end{pmatrix} \text{ is iso,}$$

skipped

Groups of order 12. If  $|G|=12$ ,  $P_4 = \mathbb{Z}/4$  or  $(\mathbb{Z}/2)^2$ ,  $P_3 = \mathbb{Z}/3$ ,

and at least one of  $P_2$  is normal, for there's not enough room for 4  $P_3$  & 3  $P_4$ 's. So  $G$  is a semi-direct Product:  $\mathbb{Z}_4 \rtimes \mathbb{Z}_3$  : must be  $\mathbb{Z}_4 \times \mathbb{Z}_3 = \mathbb{Z}_{12}$  (not  $\mathbb{Z}_6$ !) ( $\text{Aut}(\mathbb{Z}_4) = \mathbb{Z}_2$ !)

$(\mathbb{Z}_2 \times \mathbb{Z}_2) \rtimes \mathbb{Z}_3$ : Either direct;  $\mathbb{Z}_2 \times \mathbb{Z}_6$  done  
or the faithful action of  $\mathbb{Z}_3$  on  $(\mathbb{Z}_2)^2$ , giving  $A_4$  skipped  
 $\langle (234) \rangle$   
 $\begin{matrix} \emptyset \\ (12)(34) \\ (13)(24) \\ (14)(23) \end{matrix}$

$\mathbb{Z}_3 \rtimes (\mathbb{Z}_2 \times \mathbb{Z}_2)$ : Either direct or  $D_6 \times \mathbb{Z}_2 = D_{12}$  done

$\mathbb{Z}_3 \rtimes \mathbb{Z}_4$ : Either direct or  $\mathbb{Z}_3 \rtimes \mathbb{Z}_4$  done

**Solvable Groups.** Def  $G$  is solvable if all quotients in its Jordan-Hölder series are Abelian.

Thm 1. IF  $N \trianglelefteq G$ ,  $G$  is solvable iff  $N$  &  $G/N$  are.

2. IF  $H \leq G$  and  $G$  is solvable, so is  $H$ .

$A \trianglelefteq B$   $H \cap A \trianglelefteq H \cap B$ ?  $\checkmark$   $\frac{H \cap B}{H \cap A} \rightarrow \frac{B}{A}$  by  $[b]_{H \cap A} \rightarrow [b]_A$  is injective.

Cor. IF a group contains  $A_n$   $n \geq 4$ , it is not solvable.

Further return HW1

HW2 due!

TT next class, Mon. Oct 20 1<sup>st</sup> - 3PM here

Material: Everything on groups. See oldies.

Final exam: Mon Dec 8 Top; Wed Dec 10 Analysis; Thu Dec 11 PDE

Algebra: Fri Dec 12 or Mon Dec 15? Time?

Groups of order 12.  $P_2 = \mathbb{Z}_4$  or  $(\mathbb{Z}_2)^2$ ,  $P_3 = \mathbb{Z}_3$ , at least one of  $P_i$  is normal, so: on board

$$G = \mathbb{Z}_3 \rtimes \mathbb{Z}_4: \text{Aut}(\mathbb{Z}_3) = \mathbb{Z}_2 \text{ so}$$

$$\mathbb{Z}_2 \text{ or no-name } \mathbb{Z}_3 \rtimes_{\text{parity}} \mathbb{Z}_4$$

$$G = \mathbb{Z}_3 \rtimes (\mathbb{Z}_2 \times \mathbb{Z}_2): \mathbb{Z}_6 \times \mathbb{Z}_2 \text{ or } S_3 \times \mathbb{Z}_2 = D_6 \quad \text{Diagram: } \triangle \text{ with internal lines}$$

$$G = \mathbb{Z}_4 \rtimes \mathbb{Z}_3: \text{Aut}(\mathbb{Z}_4) = \mathbb{Z}_2 \Rightarrow \mathbb{Z}_2$$

$$G = (\mathbb{Z}_2 \times \mathbb{Z}_2) \rtimes \mathbb{Z}_3: \text{Aut}(\mathbb{Z}_2 \times \mathbb{Z}_2) = S_3 \Rightarrow$$

(direct)  $\mathbb{Z}_6 \times \mathbb{Z}_2$

or the fun action of  $\mathbb{Z}_3$  on  $(\mathbb{Z}_2)^2$ , giving  $A_4$

$\langle (234) \rangle$

$\begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}$	$\begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}$
$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$	$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix}$
$\begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix}$	$\begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}$

Groups of order 21.  $\mathbb{Z}_{21}$ ,  $\mathbb{Z}_7 \rtimes \mathbb{Z}_3 = \langle x \rangle \rtimes \langle y \rangle$

$$\text{Aut}(\mathbb{Z}_7) = \mathbb{Z}_6 = \langle v \rangle; v(x) = x^3; y \text{ acts } v^0 \text{ or } v^2 \text{ or } v^4$$

$$yxy^{-1} = \underbrace{x}_{\mathbb{Z}_7} \text{ or } \underbrace{x^2 \text{ or } x^4}_{\text{isomorphic}}$$

Exercise:  $\phi: H \rightarrow \text{Aut}(N); \eta \in \text{Aut} H, \nu \in \text{Aut}(N)$

$$\phi\eta: H \rightarrow \text{Aut}(N) \quad (\phi^\nu)_h = \nu^{-1} \circ \phi_h \circ \nu$$

$$\phi^\nu \in \text{Hom}(H, \text{Aut}(N))$$

$$\text{Then } N \rtimes_{\phi} H \cong N \rtimes_{\phi\eta} H \cong N \rtimes_{\phi^\nu} H.$$

In our case  $\phi_4 = \phi_2 \circ \eta$  where  $\eta: \mathbb{Z}_3 \rightarrow \mathbb{Z}_3$



In our case  $\phi_4 = \phi_2 \circ \eta$  done line where  $\eta: \mathbb{Z}/3 \rightarrow \mathbb{Z}/3$  is multiplication by 2.

**Solvable Groups.** Def  $G$  is solvable if all quotients in its Jordan-Hölder series are Abelian.

Thm 1. IF  $N \trianglelefteq G$ ,  $G$  is solvable iff  $N$  &  $G/N$  are.

2. IF  $H \leq G$  and  $G$  is solvable, so is  $H$ .

$A \trianglelefteq B$   $H \cap A \trianglelefteq H \cap B$  ?  $\checkmark$   $\frac{H \cap B}{H \cap A} \rightarrow \frac{B}{A}$  by  $[b]_{H \cap A} \rightarrow [b]_A$  is injective.

Cor. IF a group contains  $A_n$   $n \geq 4$ , it is not solvable.

Term test line.

## Rings.

**Definition 2.1.1.** A **ring** consists of a set  $R$  together with binary operations  $+$  and  $\cdot$  satisfying:

1.  $(R, +)$  forms an abelian group,
2.  $(a \cdot b) \cdot c = a \cdot (b \cdot c) \forall a, b, c \in R$ ,
3.  $\exists 1 \neq 0 \in R$  such that  $a \cdot 1 = 1 \cdot a = a \forall a \in R$ , and
4.  $a \cdot (b + c) = a \cdot b + a \cdot c$  and  $(a + b) \cdot c = a \cdot c + b \cdot c \forall a, b, c \in R$ .

Also define:  
Commutative ring.

**Examples.**  $\mathbb{Z}$ ,  $R[x]$ ,  $M_{n \times n}(R)$

Morp isms,  $\left( \begin{array}{l} \text{Examples: } 1. \mathbb{Z} \rightarrow \mathbb{Z}/n \\ 2. R \rightarrow R[x] \text{ at deg } 0 \\ 3. R \rightarrow M_{n \times n}(R) \text{ as diag} \\ 4. \text{ev}_a: R[x] \rightarrow R \text{ (if } R \text{ is commutative)} \\ 5. M_{n \times n}(R[x]) \cong M_{n \times n}(R)[x] \end{array} \right)$

im, subring, ker, ideal.

**Q.** Is every ideal a quotient.

**Ans.** Define  $R/I$ .

Good luck w/ term test!

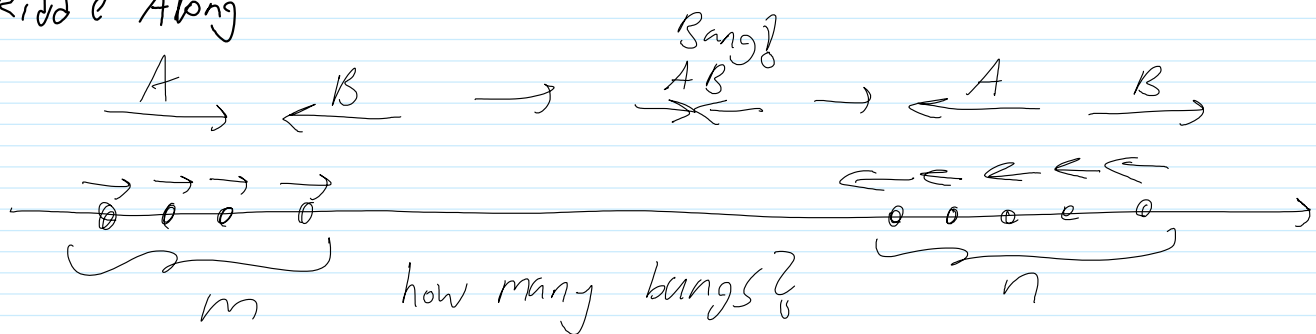
# 14-1100 Oct 20, hours 17-18: Term Test

October-22-14 4:06 PM

Return TT, etc.

HW3 on web, more may be added next week.

Riddle Along



Solvable Groups. Def  $G$  is solvable if all quotients in its Jordan-Hölder series are Abelian.

Thm 1. IF  $N \triangleleft G$ ,  $G$  is solvable iff  $N$  &  $G/N$  are.

2. IF  $H \leq G$  and  $G$  is solvable, so is  $H$ .

$A \triangleleft B$   $H \cap A \triangleleft H \cap B$ ?  $\checkmark$   $\frac{H \cap B}{H \cap A} \rightarrow \frac{B}{A}$  by  $[b]_{H \cap A} \rightarrow [b]_A$  is injective.

Cor. IF a group contains  $A_n$   $n \geq 5$ , it is not solvable.

## Rings.

**Definition 2.1.1.** A ring consists of a set  $R$  together with binary operations  $+$  and  $\cdot$  satisfying:

1.  $(R, +)$  forms an abelian group,
2.  $(a \cdot b) \cdot c = a \cdot (b \cdot c) \forall a, b, c \in R$ ,
3.  $\exists 1 \neq 0 \in R$  such that  $a \cdot 1 = 1 \cdot a = a \forall a \in R$ , and
4.  $a \cdot (b + c) = a \cdot b + a \cdot c$  and  $(a + b) \cdot c = a \cdot c + b \cdot c \forall a, b, c \in R$ .

Also define:  
Commutative ring.

Examples.  $\mathbb{Z}$ ,  $R[x]$ ,  $M_{n \times n}(R)$ ,  $RG$

Morphisms,

Examples: 1.  $\mathbb{Z} \rightarrow \mathbb{Z}/n$

2.  $R \rightarrow R[x]$  at  $\deg 0$

3.  $R \rightarrow M_{n \times n}(R)$  as diag

4.  $\ell u: R[x] \rightarrow R$   
(if  $R$  is commutative)

done  
link

5.  $M_{n \times n}(R[x]) \cong M_{n \times n}(R)[x]$

6. IF  $\varphi: G \rightarrow H$ ,  $\varphi_*: RG \rightarrow RH$

Cayley-Hamilton A matrix annihilates its characteristic poly:  
Let  $A \in M_{n \times n}(R)$ ,  $R$  commutative. Set

$$\chi_A(t) = \det(tI - A). \text{ Then } \chi_A(A) = 0$$

Wrong proof.  $\chi_A(A) = \det(AI - A) = \det(0) = 0$

Nonsense! Would have worked for trace just as well!

$$\chi_A^{\text{tr}} = \text{tr}(tI - A) = nt - \text{tr}(A)$$

$$\text{So } A = \frac{\text{tr} A}{n} I$$

The issue:

$$M_{n \times n}(R)[t] \xrightarrow{\det} R[t]$$

$$\downarrow \text{ev}_A$$

$$\downarrow \text{ev}_A$$

$$M_{n \times n}(R) \xrightarrow{?} M_{n \times n}(R)$$

Right proof.

in  $M_{n \times n}(R[t])$

in  $M_{n \times n}(R)[t]$

$$\det(tI - A) \cdot I = \text{adj}(tI - A)(tI - A) = \left(\sum B_i t^i\right)(tI - A) \text{ in } M_{n \times n}(R[t])$$

now substitute  $t = A$ . The  $B_i$ 's commute with  $A$

$$\text{because } (tI - A)\text{adj}(tI - A) = \text{adj}(tI - A)(tI - A).$$

im, subring, ker, ideal.

Q. Is every ideal a quotient.

Ans. Define  $R/I$ .

HW3 2 questions added!

Riddle Along 1 2 3 4 5 6 7 8 9

Two players alternate drawing cards from the above deck. The first player to have 3 cards that add up to 15, wins. Would you like to be the first to move or the second?

Reminders 1. Rings:  $(R, +, \times, 0 \neq 1)$

2.  $R[x]$ ,  $M_{n \times n}(R)$ ,  $RG$

3. Morphisms (make rings a "category")  $[F(1)=1]$

Further examples.

1. If  $\varphi: G \rightarrow H$ ,  $\varphi_*: RG \rightarrow RH$

2.  $M_{n \times n}(R[x]) \cong M_{n \times n}(R)[x]$

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so  $A = \frac{\text{tr } A}{n} I$

The issue:  $M_{n \times n}(R)[t] \xrightarrow{\det} R[t] \left\{ \begin{array}{l} \text{not} \\ \text{mentioned.} \end{array} \right.$   
 $\downarrow \text{ev}_A$   
 $M_{n \times n}(R) \xrightarrow{?} M_{n \times n}(R)$   
 $\downarrow \text{ev}_A$

Right proof.

$$\begin{array}{ccc} \text{in } M_{n \times n}(R[t]) & & \text{in } M_{n \times n}(R)[t] \\ \downarrow & & \downarrow \\ \det(tI - A) \cdot I = \text{adj}(tI - A)(tI - A) = (\sum B_i t^i)(tI - A) & \text{in} & \end{array}$$

now substitute  $t = A$ . The  $B_i$ 's commute with  $A$

because  $(tI - A) \text{adj}(tI - A) = \text{adj}(tI - A)(tI - A)$ .

see 2015-12

Im, subring, ker, ideal. (ideals are subrings but never subrings  $\subseteq$ )  
 Q. Is every <sup>proper</sup> ideal a kernel?  
 Ans. Define  $R/I$ .

Example.  $\mathbb{R}[x]/\langle x^2+1 \rangle = \mathbb{C}$ ,

The Isomorphism Theorems. 1.  $\psi: R \rightarrow S \Rightarrow R/\ker \psi = \text{im } \psi$ .

(Example:  $\text{ev}_i: \mathbb{R}[x] \rightarrow \mathbb{C} \Rightarrow \mathbb{R}_i \cong \mathbb{C}$ )

done  
line

2.  $\frac{A+I}{I} \cong \frac{A}{A \cap I}$   $A \subseteq R$  subring,  $I \subseteq R$  proper ideal.

3.  $I \subseteq J \subseteq R$  ideals  $\Rightarrow \frac{R/I}{J/I} \cong R/J$

4. Given an ideal  $I$  of  $R$ , there's a bijection between  
 ideals  $I \subseteq J \subseteq R$  & ideals of  $R/I$ .

From this point, our goal  
 is "modules over PID"

Better Rings. 1. The ultimate:

Field [commutative,  $F$  of a group]

("division ring", if not commutative)

Example:  $\mathbb{H} = \{a+bi+cj+dk\} / \begin{matrix} i^2=j^2=k^2=-1 \\ ij=k \\ \text{useful for 3D rotations, etc.} \end{matrix}$

[almost all of  
 high-school &  
 freshman algebra  
 carries through]

2. (Integral) domains: commutative, has no 0-divisors.

How make? For ideals which,  $R/I$  is a field or a domain?

... From now on,  $R$  is commutative.

Maximal Ideals. 1. Definition.

2.  $I \subseteq R$  is maximal  $\Leftrightarrow R/I$  is a field.

Fishy proof: Use the 4th isomorphism theorem.

Honest proof:  $\Rightarrow: x \notin I \Rightarrow Rx+I = R \Rightarrow \exists y \in R \ yx+I = 1+I$

$\Leftarrow J \not\supseteq I, x \in J \setminus I \Rightarrow [x]_I \neq 0 \Rightarrow \exists y \ xy-1 \in I \Rightarrow 1 \in J$

Examples. 1.  $p\mathbb{Z}$  is a maximal ideal in  $\mathbb{Z}$ .

2.  $S = \ell^\infty = \{ \text{bndd seq's in } \mathbb{R} \}$   $A_n = \{ (a_i) : a_n = 0 \}$

<sup>Fishy</sup> Theorem. Every ideal is contained in a maximal ideal.

Proof using Zorn's Lemma.

Theorem There exists a function

$\text{Lim}: \{\text{bdd seq's in } \mathbb{R}\} \rightarrow \mathbb{R} \text{ s.t.}$

1. IF  $(a_n)$  is convergent,  $\lim a_n = \text{Lim } a_n$ .
2.  $\text{Lim } (a_n + b_n) = \text{Lim } (a_n) + \text{Lim } (b_n)$
3.  $\text{Lim } (a_n b_n) = \text{Lim } (a_n) \cdot \text{Lim } (b_n)$  + more....

Proof.  $S = \{\text{bdd seq's in } \mathbb{R}\}$   $I = \{(a_n) : a_n \neq 0 \text{ for finitely many } n\}$

$J$  - a maximal ideal containing  $I$ .

$\text{Lim}: S \rightarrow S/J \cong \mathbb{R}$

Prime Ideals. 1. Definition  $P \subset R$  is prime if  $ab \in P \Rightarrow a \in P \text{ or } b \in P$ .

2. Theorem.  $R/P$  is a domain iff  $P$  is prime.

Proof.  $\Rightarrow ab \in P \Rightarrow [ab] = 0 \Rightarrow [a][b] = 0 \Rightarrow \begin{matrix} [a] = 0 \Rightarrow a \in P \\ [b] = 0 \Rightarrow b \in P \end{matrix}$

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Theorem. A maximal ideal is prime.

target line

From this point on,  $R$  is a commutative integral domain.

" $a, b$  are associates"

Primes. 1.  $a|b$  [ $a \neq 0, \exists q \text{ s.t. } aq = b$ ] ( $a|b \wedge b|a \Rightarrow a = ub$ )

2.  $\gcd(a, b) = q$  ;  $\gcd = q$  &  $\gcd = q' \Rightarrow q' = uq$ .

3. Primes:  $p \neq 0$  non-unit  $p|ab \Rightarrow p|a \text{ or } p|b$

$p$  is prime iff  $\langle p \rangle$  is prime ideal.

4. Irreducible  $x = ab \Rightarrow a \in R^* \vee b \in R^*$

Claim. prime  $\Rightarrow$  irreducible

$p = ab \Rightarrow p|a \Rightarrow a = pc$

$\Rightarrow p = pcb \Rightarrow cb = 1 \Rightarrow b \in R^*$

Counterexample: in  $\mathbb{Z}[\sqrt{-5}]$ ,  
2 is irred (for norm reasons)  
but not prime, as

$2|(1-\sqrt{-5})(1+\sqrt{-5}) = 6$

UFDs. Def. Every non-zero element can be factored into primes.

Thm. Uniqueness up to units & a permutation.

Thm. In a UFD, Prime  $\Leftrightarrow$  irreducible.

pf If an irred. is decomposed, the decomposition must

have length 1.

Thm.  $UFD \Leftrightarrow$  evry  $x \neq 0$  has a unique decomposition  
into irreducibles.  $\text{pf}$  need  $\text{irred} \Rightarrow \text{prime}$ . If  $x$  is  $\text{irred}$  &  $x|ab$ , then  
 $ax = a_1 \dots a_n b_1 \dots b_m \Rightarrow x \sim a_i$  or  $x \sim b_j \Rightarrow x|a$  or  $x|b$   
 $\text{irreds}$

Thm. In a  $UFD$  gcd's always exist.



Reminders Ideal:  $0 \in I$ ,  $I + I \subset I$ ,  $-I \subset I$ ,  $R \cap I \subset I$ ,  $I \cap R \subset R$   
 $R/I$ , Iso 1: Given  $\varphi: R \rightarrow S$ ,  $R/\ker \varphi \cong \text{im } \varphi$

2.  $\frac{A+I}{I} \cong \frac{A}{A \cap I}$   $A \subset R$  subring,  $I \subset R$  proper ideal.

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Example.  $S = \{ \text{bdd seq's in } \mathbb{R} \}$   $I = \{ (a_n) : a_n \rightarrow 0 \}$   $[a_n = 0 \text{ a.e.}]$   
 $J$  - a maximal ideal containing  $I$ .

Lim:  $S \rightarrow S/J \cong \mathbb{R}$   $[R \rightarrow S/J \text{ is obvious; } \dots]$

$\text{Lim}: S \rightarrow S/J = \mathbb{R}$  [ $\mathbb{R} \rightarrow S/J$  is obvious; the other direction is not]  
**Theorem**  $\text{Lim}$  satisfies:

1. If  $(a_n)$  is convergent,  $\lim a_n = \text{Lim } a_n$ .
2.  $\text{Lim}(a_n + b_n) = \text{Lim}(a_n) + \text{Lim}(b_n)$
3.  $\text{Lim}(a_n b_n) = \text{Lim}(a_n) \cdot \text{Lim}(b_n)$  + more.... □

**Prime Ideals.** 1. Definition  $P \subset R$  is prime if  $ab \in P \Rightarrow a \in P$  or  $b \in P$ .

2. Theorem.  $R/P$  is a domain iff  $P$  is prime.  
 Proof:  $\Rightarrow ab \in P \Rightarrow [ab] = 0 \Rightarrow [a][b] = 0 \Rightarrow \begin{matrix} [a]=0 \Rightarrow a \in P \\ [b]=0 \Rightarrow b \in P \end{matrix}$   
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**Theorem.** A maximal ideal is prime.

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**Primes.** 1.  $a|b$  [ $a \neq 0, \exists q$  s.t.  $aq = b$ ] ( $a|b \wedge b|a \Rightarrow a = ub$ ) "a, b are associates"

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4. Irreducible  $x = ab \Rightarrow a \in R^* \vee b \in R^*$

<p><b>Claim.</b> prime <math>\Rightarrow</math> irreducible</p> <p><math>p = ab \Rightarrow p a \Rightarrow a = pc</math></p> <p><math>\Rightarrow p = pcb \Rightarrow cb = 1 \Rightarrow b \in R^*</math></p>	<p>Counterexample: in <math>\mathbb{Z}[\sqrt{-5}]</math>,  <math>2</math> is irred (for norm reasons)              but not prime, as  <math>2 (1-\sqrt{-5})(1+\sqrt{-5}) = 6</math></p>
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**UFDs.** Def: Every non-zero element can be factored into primes.

Thm. Uniqueness up to units & a permutation.

Thm. In a UFD, Prime  $\Leftrightarrow$  irreducible.

pf If an irred. is decomposed, the decomposition must have length 1.

Thm. UFD  $\Leftrightarrow$  evry  $x \neq 0$  has a unique decomposition  
as prod. irred  $\Rightarrow$  prime. If  $x$  is irred &  $x|ab$ , then

into irreducibles.  $\frac{x}{z} = \underbrace{a_1 \dots a_n b_1 \dots b_m}_{\text{irreds}} \Rightarrow x \sim a_i \text{ or } x \sim b_j \Rightarrow x|a_i \text{ or } x|b_j$

Thm. In a UFD gcd's always exist.

(141102) Assaf's riddle: <sup>50</sup>~~5~~ kids share a loot of <sup>50</sup>~~n~~ in-wrapping hal-loween candies. The first kid proposes a way to split the loot; if it is not accepted by a strict majority (her included), she's ~~left out~~ <sup>goes home</sup> and the second proposes a split, etc. How is the loot split?

Global goal: IT3CSW:  $M$  f.g. module over a PID  $R \Rightarrow$  Uniquely  
 $M \cong R^k \oplus \bigoplus R/(p_i^{s_i})$   $p_i$  prime  $s_i \geq 1$

Cor 1.  $A$  f.g. Abelian  $\Rightarrow A \cong \mathbb{Z}^k \oplus \bigoplus \mathbb{Z}/p_i^{s_i}$

Cor 2.  $A \in M_{n \times n}(\mathbb{C})$  has a "Jordan form"

No Joy Agenda. Euc  $\Rightarrow$  PID  $\Rightarrow$  UFD.

Reminders  $R/I$  a field  $\Leftrightarrow I$  is maximal.

$R/I$  a domain  $(ab=0 \Rightarrow (a=0) \vee (b=0))$  start line  
 $\Leftrightarrow I$  is prime.

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Theorem. A maximal ideal is prime.

From this point on,  $R$  is a commutative integral domain.

Divisibility & Primes. 1.  $a|b$   $[a \neq 0, \exists q \text{ s.t. } aq=b]$   $a|b \wedge b|a \Rightarrow a=ub$   $\leftarrow$  " $a, b$  are associates"

2.  $\gcd(a, b) = q$   $\downarrow$   $\gcd = q$  &  $\gcd = q' \Rightarrow q = uq'$

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Thm. UFD  $\Leftrightarrow$  every  $x \neq 0$  has a unique decomposition

into irreducibles. pf need irred  $\Rightarrow$  prime. If  $x$  is irred &  $x|ab$ , then  
 $zx = \underbrace{a_1 \dots a_n b_1 \dots b_m}_{\text{irreds}} \Rightarrow x \sim a_i \text{ or } x \sim b_j \Rightarrow x|a \text{ or } x|b$

Thm. In a UFD gcd's always exist.

How show UFD? Norm  $\Rightarrow$  "PID"  $\Rightarrow$  UFD.

Def. Euclidean domain: has a "norm"  $e: R \setminus \{0\} \rightarrow \mathbb{N}$  s.t.

1.  $e(ab) \geq e(a)$
2.  $\forall a, b \exists q, r$  s.t.  $a = qb + r$  &  $r = 0$  or  $e(r) < e(b)$

Example. 1.  $\mathbb{Z}$

$$\text{Example } \frac{a = x^3 - 2x^2 - 5x + 12}{b = x^2 + 1}$$

2.  $F[x]$

$$\dots r = -6x + 14 \quad \left. \begin{array}{l} a(1) = 14 - 6 \end{array} \right\} \text{why?}$$

theorem. A Euclidean domain is a "PID" (def).

(Thm: a PID is a UFD, later)

Proposition. In a PID, every prime ideal is maximal.

pf.  $I = \langle p \rangle$  prime,  $I \subset J = \langle x \rangle \subset R \Rightarrow p = ax \Rightarrow$

$$(a \in R^* \Rightarrow I = J) \vee (x \in R^* \Rightarrow J = R)$$

done  
line

theorem. PID  $\Rightarrow$  UFD.

What proof. Take  $x = x_1$  unless  $x_1 \in R^*$ ,  $x_1 \in M_1$  where  $M_1$  is a maximal ideal containing  $\langle x \rangle$ .  $M_1 = \langle p_1 \rangle$ ,

$p_1$  prime. So  $x_1 = p_1 x_2$  unless  $x_2 \in R^*$   $x_2 \in \langle x_3 \rangle \subset M_2$  maximal

$M_2 = \langle p_2 \rangle$ ,  $x_1 = p_1 x_2$ ,  $x_2 = p_2 x_3 \dots$  if process was infinite,

$M_2 = \langle p_2 \rangle$ ,  $x_2 = p_2 x_3, \dots$  if process was infinite,

$$\langle x_1 \rangle \subsetneq \langle x_2 \rangle \subsetneq \langle x_3 \rangle \subsetneq \dots$$

But a PID is "Noetherian",

so the process must terminate.

$$\text{So } x = x_1 = p_1 x_2 = p_1 p_2 x_3 = \dots = p_1 p_2 \dots p_n u$$

$\langle x_n \rangle \subsetneq \langle x_{n+1} \rangle$  as  $x_n = p_n x_{n+1}$   
if  $x_{n+1} \in \langle x_n \rangle$ ,  $x_{n+1} = a x_n$  so  
 $x_n = p_n a x_n$  &  $p$ 's not prime.

Theorem. In a PID  $\langle a, b \rangle = \langle \gcd(a, b) \rangle$ . (so  $\gcd(a, b) = sa + tb$ )

target  
line

The Euclidean Algorithm. In a Euc. Domain, a practical algorithm for finding  $s(a, b)$  &  $t(a, b)$  as above: WLOG,  $\ell(a) \geq \ell(b)$

If  $\langle a, b \rangle = \langle b \rangle$ , take  $(s, t) = (0, 1)$ . Otherwise

$$a = bq + r, \ell(r) < \ell(b),$$

$\langle a, b \rangle = \langle b, r \rangle$  so if  $g = s'b + t'r$ , then

$$g = s'b + t'(a - bq) = \underbrace{t'}_s a + \underbrace{(s' - t'q)}_t b$$

Theorem.  $R$  is a PID iff it has a "Dedekind-Hasse"

norm:  $d: R - \{0\} \rightarrow \mathbb{N}_{>0}$  [or add  $d(0) = 0$ ]

s.t. if  $a, b \neq 0$  either  $a \in \langle b \rangle$  or  $\exists 0 \neq x \in \langle a, b \rangle$

w/  $d(x) < d(b)$ .

pf.  $\Leftarrow$  as before.  $\Rightarrow$  Replace every prime by 2, get

even a "multiplicative" D-H norm.

HW. HW3 due, HW4 on wed

Riddle along: A game: Player A writes the numbers 1-18 on the faces of three blank dice, to her liking. Player B takes one of the 3 dice. Player B takes one of the remaining two, and throws away the third. Player A and B then play 1,000 rounds of "dice war" with the dice they hold. Whom would you rather be, player A or player B?

Global goal:  $M$  f.g. module over a PID  $R \Rightarrow$  uniquely  
IT3C4W  
 $M \cong R^k \oplus \bigoplus R/(p_i^{s_i})$   $p_i$  prime  
 $s_i \geq 1$

Cor1.  $A$  f.g. Abelian  $\Rightarrow A \cong \mathbb{Z}^k \oplus \bigoplus \mathbb{Z}/p_i^{s_i}$

Cor2.  $A \in M_{n \times n}(\mathbb{C})$  has a "Jordan form"

Today. Finish rings, start modules.

Reminders. Euc  $\Rightarrow$  PID  $\nRightarrow$  UFD

theorem. PID  $\Rightarrow$  UFD.

What proof. Take  $x = x_1$  unless  $x_1 \in R^\times$ ,  $x_1 \in M_1$  where  $M_1$  is a maximal ideal containing  $\langle x \rangle$ .  $M_1 = \langle p_1 \rangle$ ,  $p_1$  prime. So  $x_1 = p_1 x_2$  unless  $x_2 \in R^\times$   $x_2 \in \langle x_3 \rangle \subset M_2$  maximal  $M_2 = \langle p_2 \rangle$ ,  $x_2 = p_2 x_3, \dots$  if process was infinite,

$$\langle x_1 \rangle \subsetneq \langle x_2 \rangle \subsetneq \langle x_3 \rangle \subsetneq \dots$$

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so the process must terminate.

$$\text{So } x = x_1 = p_1 x_2 = p_1 p_2 x_3 = \dots = p_1 p_2 \dots p_n u$$

theorem. In a PID  $\langle a, b \rangle = \langle \gcd(a, b) \rangle$ . (so  $\gcd(a, b) = sa + tb$ )

The Euclidean Algorithm. In a Euc. Domain, a practical algorithm for finding  $s(a, b)$  &  $t(a, b)$  as above: WLOG,  $\ell(a) \geq \ell(b)$

If  $\langle a, b \rangle = \langle b \rangle$ , take  $(s, t) = (0, 1)$ . Otherwise

$$a = bq + r, \ell(r) < \ell(b),$$

$$\langle a, b \rangle = \langle b, r \rangle \text{ so if } g = s'b + t'r, \text{ then}$$

$$g = s'b + t'(a - bq) = \underbrace{t'}_s a + \underbrace{(s' - t'q)}_r b$$

**Theorem.**  $R$  is a PID iff it has a "Dedekind-Hasse"

norm:  $d: R \setminus \{0\} \rightarrow \mathbb{N}_{>0}$  [or add  $d(0) = 0$ ]

s.t. if  $a, b \neq 0$  either  $a \in \langle b \rangle$  or  $\exists 0 \neq x \in \langle a, b \rangle$   
w/  $d(x) < d(b)$ .

skipped.

**pf.**  $\Leftarrow$  as before.  $\Rightarrow$  Replace every prime by 2, get  
even a "multiplicative" D-H norm.

target line

**Definition.** An  $R$ -module: "A vector space over a ring".

**Examples.** 1. V.S. over a field.

2. Abelian groups over  $\mathbb{Z}$ .

3. Given  $T: V \rightarrow V$ ,  $V$  over  $F[x]$ .

4. Given ideal  $I \subset R$ ,  $R/I$  over  $R$ .

5. Column vectors  $R^n$  over  $M_{n \times n}$  (Left module  $R$ -mod)  
row vectors  $(R^n)^T$  over  $M_{n \times n}$  (right module  $\text{mod-}R$ )

done  
line

**Def/claim.**  $R$ -mod &  $\text{mod-}R$  are categories.

**Def/claim.** Submodules,  $\ker \phi$ ,  $\text{im } \phi$ ,  $M/N$

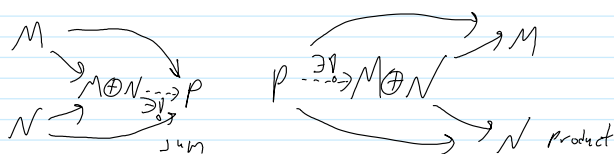
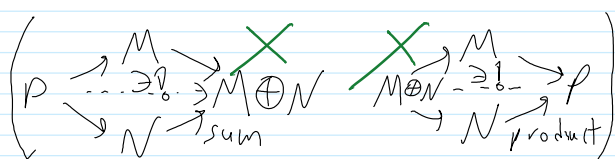
**Boring Theorems.** 1.  $\phi: M \rightarrow N \Rightarrow M/\ker \phi \cong \text{im } \phi$

$$2. A, B \subset M \Rightarrow \frac{A+B}{B} \cong \frac{A}{A \cap B}$$

$$3. A \subset B \subset M \Rightarrow \frac{M/A}{B/A} \cong \frac{M}{B}$$

4. Also dual.

**Direct sums.**  $M, N \Rightarrow M \oplus N$



differ for infinite families!



$$\text{Hom}(\bigoplus N_i, \bigoplus M_i) = \left\{ \begin{pmatrix} \cdot & \cdot \\ a_{m1} & a_{nn} \end{pmatrix} : a_{ij} \in \text{Hom}(M_i, N_j) \right\}$$

Example:  $\dim(V \oplus W) = \dim V + \dim W.$

Example: if  $\gcd(a,b)=1$   $1=sa+tb$  [e.g., if  $R$  is a PID]

$$\frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle} \cong \frac{R}{\langle ab \rangle} \quad \text{via} \quad \begin{array}{ccc} R/\langle a \rangle & \xrightarrow{t \cdot b} & R/\langle ab \rangle \\ \oplus & & \uparrow \\ R/\langle b \rangle & \xrightarrow{s \cdot a} & R/\langle ab \rangle \end{array} \begin{array}{c} \xrightarrow{1} R/\langle a \rangle \\ \oplus \\ \xrightarrow{1} R/\langle b \rangle \end{array}$$

$$\mathbb{Z}/7 \oplus \mathbb{Z}/11 \oplus \mathbb{Z}/13 \cong \mathbb{Z}/77 \oplus \mathbb{Z}/13 \cong \mathbb{Z}/1001 \quad \text{"the chinese remainder theorem"}$$

The inclusions

$$\text{UFD} \stackrel{1}{\not\subset} \text{PID} \stackrel{2}{\not\subset} \text{Euc}$$

are strict.

1. Many examples; especially polynomial rings in several variables and  $\mathbb{Z}[x]$ . (In general,  $R \text{ UFD} \Rightarrow R[x] \text{ UFD}$ ).
2. Examples are hard. The easiest seems to be  $\mathbb{Z}\left[\frac{1+\sqrt{-19}}{2}\right]$ .

A sequence of exercises leading to a proof is in eprints/Bergman:

Math 250A, G. Bergman, 2002

**A principal ideal domain that is not Euclidean**  
developed as a series of exercises

Office Hour this week Wed 1<sup>30</sup> - 2<sup>30</sup> (not at 2<sup>30</sup>)

Global goal:  $M$  f.g. module over a PID  $R \Rightarrow$  Uniquely  
 IT3C4W:  $M \cong R^k \oplus \bigoplus R/(p_i^{s_i})$   $p_i$  prime  $s_i \geq 1$

Cor 1.  $A$  f.g. Abelian  $\Rightarrow A \cong \mathbb{Z}^k \oplus \bigoplus \mathbb{Z}/p_i^{s_i}$

Cor 2.  $A \in M_{n \times n}(\mathbb{C})$  has a "Jordan Form"

Today. Further dull technicalities, then proof of existence side of Thm.

Euc  $\Rightarrow$  PID  $\Rightarrow$  UFD.

Many UFD's are not PID's.

$\mathbb{Z}[\frac{1+\sqrt{-19}}{2}]$  is a PID but is not Euclidean.

Theorem.  $R$  is a PID iff it has a "Dedekind-Hasse"

norm:  $d: R \setminus \{0\} \rightarrow \mathbb{N}_{>0}$  [or add  $d(0)=0$ ]

s.t. if  $a, b \neq 0$  either  $a \in \langle b \rangle$  or  $\exists 0 \neq x \in \langle a, b \rangle$   
 w/  $d(x) < d(b)$ .

pf.  $\Leftarrow$  as before.  $\Rightarrow$  Replace every prime by 2, get  
 even a "multiplicative" D-H norm.

Reminder. Modules.

Def/claim.  $R$ -mod & mod- $R$  are categories.

Def/claim. Submodules,  $\ker \phi$ ,  $\text{im } \phi$ ,  $M/N$

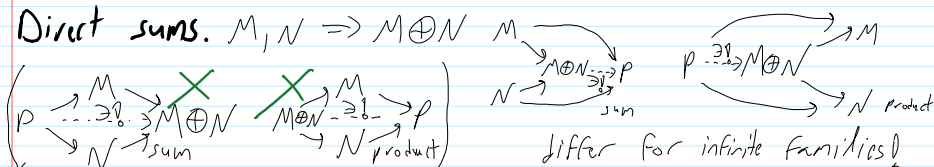
Boxing Theorems. 1.  $\phi: M \rightarrow N \Rightarrow M/\ker \phi \cong \text{im } \phi$

2.  $A, B \subset M \Rightarrow \frac{A+B}{B} \cong \frac{A}{A \cap B}$

3.  $A \subset B \subset M \Rightarrow \frac{M/A}{B/A} \cong M/B$

4. Also dull.

Direct sums.  $M, N \Rightarrow M \oplus N$



$\text{Hom}(\bigoplus N_i, \bigoplus M_i) = \left\{ \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix} : a_{ij} \in \text{Hom}(M_i, N_j) \right\}$

Example: if  $\gcd(a, b) = 1$   $1 = sa + tb$  [e.g., if  $R$  is a PID] PM. I should have done  $1 = \text{lcm}$   
 $\frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle} \cong \frac{R}{\langle 9 \rangle} \oplus \frac{R}{\langle 7 \rangle} \cong \frac{R}{\langle 63 \rangle}$

Example: If  $\gcd(a,b)=1$   $1=sa+tb$  [e.g., if  $R$  is a PID]

PM. I should have done  $1=lc_m$   
 $\frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle} \cong \frac{R}{\langle g \rangle} \oplus \frac{R}{\langle 0 \rangle} = \frac{R}{\langle g \rangle}$   
 in a way compatible w/  
 $\frac{R}{\langle ab \rangle} \rightarrow \frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle}$  by (1)  
 in the case of  $\gcd=1$ .

$$\frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle} \cong \frac{R}{\langle ab \rangle} \text{ via } \begin{matrix} (1b \ 5a) & \begin{pmatrix} 1 \\ 1 \end{pmatrix} \\ R/\langle a \rangle \xrightarrow{tb} R/\langle ab \rangle \xrightarrow{1} R/\langle a \rangle \\ \oplus \\ R/\langle b \rangle \xrightarrow{sa} R/\langle ab \rangle \xrightarrow{1} R/\langle b \rangle \end{matrix}$$

$\mathbb{Z}_7 \oplus \mathbb{Z}_{11} \oplus \mathbb{Z}_{13} \cong \mathbb{Z}_{77} \oplus \mathbb{Z}_{13} \cong \mathbb{Z}_{1001}$  "the chinese remainder theorem"

Let  $R$  be a PID...

Sketch  $\{ \text{matrices} \} / \text{row \& col ops} \xrightarrow{\text{onto}} \{ \text{f.g. modules} \}$   
 finite by infinite but the infinite is just a nuisance

PM. I should have gone  
 1.  $M_{n \times m} \rightarrow$  modules w/  $n$  genes &  $m$  reals.  
 2.  $M_{n \times X} \rightarrow$  f.g. modules  
 3. The above is surjective.

So we're back to Gaussian elimination!

Def  $M$  is "finitely generated" if  $\exists g_1, \dots, g_n \in M$   
 s.t.  $M = \langle \sum a_i g_i : a_i \in R \rangle$ .

$$R^X \xrightarrow{A} R^g \xrightarrow{\pi} M \quad \ker \pi = \langle r_x : x \in X \rangle$$

$$A = \left( \underbrace{\quad}_x \right) \} g \quad A \in M_{g \times X}(R)$$

... In general, every  $g \times X$  matrix determines a f.g. module, and every f.g. module arises in this way.

Examples.  $(1)$ ,  $(a)$ ,  $(0)$

Exercise. If  $C = \begin{pmatrix} A & 0 \\ 0 & B \end{pmatrix}$ , then  $M_C = M_A \oplus M_B$

Comment. May add/remove 0 columns.

don't  
line

$$\begin{matrix} R^X \xrightarrow{A} R^g \\ \uparrow Q \quad \downarrow P \\ R^X \xrightarrow{A'} R^g \end{matrix} \quad \begin{matrix} \text{Claim if } P, Q \text{ are invertible} \\ \text{on the left, then} \\ M = R^g / \text{im } A \text{ and } M' = R^g / \text{im } A' \\ \text{are isomorphic.} \end{matrix}$$

PF  $\phi: M \rightarrow M'$  by  $[\alpha]_{\text{im } A} \rightarrow [P\alpha]_{\text{im } A'}$

$P$  can be interpreted as  $g \times g$  matrix

$Q$  can be interpreted as an  $X \times X$  column-finite matrix;  $A' = PAQ$

... Can do arbitrary row operations on  $A$ ,  
 and arbitrary invertible column ops, provided each column is touched finitely many times.

each column is touched finitely many times.

Of all the matrices reachable from  $A$ , let  $A'$  be the one having an entry with the smallest D-H norm; wlog, that entry is  $a_{11}$ .

Claim  $a_{11}$  divides all other entries in its row & column.

PF1 for a Euclidean domain.

PF2 In a PID, if  $q = \gcd(a, b) = sa + tb$ , then

$$(a \ b) \begin{pmatrix} s & -b/q \\ t & a/q \end{pmatrix} = (q \ 0), \text{ while } \begin{pmatrix} s & -b/q \\ t & a/q \end{pmatrix}^{-1} = \begin{pmatrix} a/q & b/q \\ -t & s \end{pmatrix} \quad \square$$

$\Rightarrow$  w.l.o.g., the row & column of  $a_{11}$  are 0 (except for  $a_{11}$ )

$\Rightarrow$  all entries of  $A$  are divisible by  $a_{11}$ :

$$A = \begin{pmatrix} a_{11} & \cdots & 0 & \cdots \\ 0 & & & \\ \vdots & & A_1 & \\ 0 & & & \end{pmatrix} \quad \begin{matrix} \text{all entries} \\ \text{divisible} \\ \text{by } a_{11} \end{matrix}$$

Continue to get  $A \sim \left( \begin{array}{c|c} a_{11} & a_{22} \\ \hline 0 & 0 \end{array} \right) \quad \left( \begin{matrix} \text{w.l.o.g., } A \\ \text{is square} \end{matrix} \right)$

$$\text{so } M \cong \bigoplus_{i=1}^g R / \langle a_{ii} \rangle \cong R^k \oplus \bigoplus_{a_1 | a_2 | \dots | p_n} R / \langle a_i \rangle$$

Goal:  $M$  f.g. / PID  $R \Rightarrow$

$$M = R^k \oplus \bigoplus_{i=1}^n R / \langle p_i^{s_i} \rangle \quad \begin{array}{l} p_i \text{ prime} \\ s_i \in \mathbb{Z}_{>0} \end{array}$$

There is a map from  $n \times m$  matrices to f.g. modules.

$$A \mapsto R^n \xrightarrow{A} R^n \longrightarrow R^n / \text{im } A =: M_A$$

Equally well,  $n \times \infty$  matrices to f.g. modules:

$$A \mapsto R^\infty \xrightarrow{A} R^n \longrightarrow R^n / \text{im } A =: M_A$$

$M_{n \times \infty}(R) \rightarrow \text{f.g. modules}$  is surjective.

Examples.  $(1)$ ,  $(a)$ ,  $(0)$

Exercise. If  $C = \begin{pmatrix} A & 0 \\ 0 & B \end{pmatrix}$ , then  $M_C = M_A \oplus M_B$

Comment. May add/remove  $0$  columns.

$$R^X \xrightarrow{A} R^g$$

$$\begin{array}{c} \uparrow Q \quad \downarrow P \\ R^X \xrightarrow{A'} R^g \end{array}$$

Claim if  $P, Q$  are invertible  
on the left, then

$M = R^g / \text{im } A$  and  $M' = R^g / \text{im } A'$   
are isomorphic.

Def  $\phi: M \rightarrow M'$  by  $[\alpha]_{\text{im } A} \mapsto [P\alpha]_{\text{im } A'}$

$P$  can be interpreted as  $g \times g$  matrix

$Q$  can be interpreted as an  $X \times X$  column-finite matrix;  $A' = PAQ$

... Can do arbitrary invertible row operations on  $A$ ,

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be the one having an entry with the smallest D-H norm; wlog, that entry is  $a_{11}$ .

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$\Rightarrow$  w.l.o.g, the row & column of  $a_{11}$  are 0 (except for  $a_{11}$ )

$\Rightarrow$  all entries of  $A$  are divisible by  $a_{11}$ :

$$A = \begin{pmatrix} a_{11} & \cdots & 0 & \cdots \\ 0 & & & \\ \vdots & & A_1 & \\ 0 & & & \end{pmatrix} \quad \begin{matrix} \text{all entries} \\ \text{divisible} \\ \text{by } a_{11} \end{matrix}$$

Continue to get  $A \sim \begin{pmatrix} a_{11} & a_{22} & & 0 \\ & & & \\ 0 & & & \\ & & & 0 \end{pmatrix} \quad \left( \begin{matrix} \text{w.l.o.g., } A \\ \text{is square} \end{matrix} \right)$

$$\text{so } M \cong \bigoplus_{i=1}^g R / \langle a_{ii} \rangle \cong R^k \oplus \bigoplus_{a_1 | a_2 | \dots | a_n} R / \langle a_i \rangle$$

Plan. JCF abstractly & in practice.

HW4 due, HW5 on web.

Riddle. 1. A spherical loaf of bread goes into a bread cutting machine which slice has the most crust?

2. Can you cover  $\bigcirc_{100}$  with  $99 \times \boxed{1}_{100}$ ?

**Corollary 2.** Over an algebraically closed field  $\mathbb{F}$ , every square matrix

$A$  is conjugate to a block diagonal matrix  $B = \begin{pmatrix} B_1 & 0 & \cdots & 0 \\ 0 & B_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & B_n \end{pmatrix}$ ,

where each  $B_i$  is either a  $1 \times 1$  matrix ( $\lambda_i$ ) for some  $\lambda_i \in \mathbb{F}$ , or an  $s_i \times s_i$  matrix with  $\lambda_i$ 's on the diagonals, 1's right below the diagonal, and 0's elsewhere,

$$\begin{pmatrix} \lambda_i & 0 & \cdots & \cdots & 0 & 0 \\ 1 & \lambda_i & \ddots & & & 0 \\ 0 & \ddots & \ddots & \ddots & & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & & \ddots & \ddots & \lambda_i & 0 \\ 0 & 0 & \cdots & 0 & 1 & \lambda_i \end{pmatrix},$$

for some  $\lambda_i \in \mathbb{F}$  and for some  $s_i \geq 2$ . Furthermore,  $B$  is unique up to a permutation of its blocks  $B_i$ .

(Corollary: good old diagonalization.)

on  
projector  
screen

JCF.  $V$  a f.d.v.s,  $A: V \rightarrow V$  linear, makes  $V$  a module over  $R := \mathbb{F}[x]$  v'ia  $xu = Au$ . Then

$$V \cong \bigoplus \frac{\mathbb{F}[x]}{(x-\lambda_i)^{s_i}}. \quad \text{What's } \frac{\mathbb{F}[x]}{(x-\lambda_i)^{s_i}}?$$

Basis:  $1, x-\lambda, (x-\lambda)^2, \dots, (x-\lambda)^{s-1}$

$A-\lambda$  acts by "shift to the right"  $\begin{pmatrix} 0 & 0 & & \\ & 0 & 1 & \\ & & \ddots & \ddots \\ & & & 0 & 1 \end{pmatrix}$

so  $A$  acts by  $\begin{pmatrix} \lambda & & & \\ & \lambda & & \\ & & \ddots & \\ & & & \lambda \end{pmatrix}$

Now let's do that in practice....

step 1. Find a presentation matrix for  $V \in R\text{-mod}$ .

w.l.o.g  $V = F^n$  and  $A \in M_{nn}(\mathbb{F})$ .  $\ker \pi = \zeta_0$

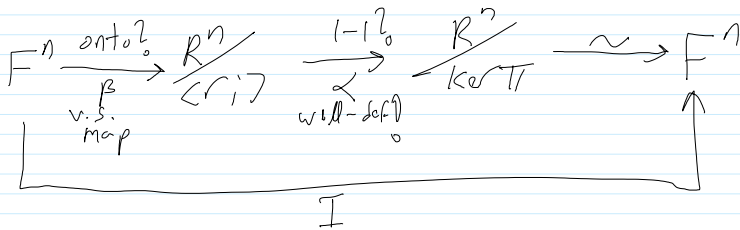
$$r_i = x e_i - A e_i \in \ker \pi \quad \left| \quad R^n \xrightarrow{xI - A} R^n \xrightarrow{\pi} F^n$$

claim  $\langle r_i \rangle = \ker \pi$

$$\begin{aligned} e_i &\longrightarrow e_i \\ x^k e_i &\longrightarrow A^k e_i \end{aligned}$$

pf Consider





We want to know if  $\alpha$  is 1-1; it is enough to show that  $\beta$  is onto; i.e., that any  $x^k e_i$  can be written, modulo  $\langle r_i \rangle$ , as a combination of  $e_j$ 's. Indeed,

$$x^k e_i = x^{k-1}(x e_i) = x^{k-1} A e_i = \dots = A^k e_i$$

Go over handout along with "run 1"

Dror Bar-Natan: Classes: 2014-15: Math 1100 Algebra I:

## JCF Tricks and Programs

### Row and Column Operations

Row operations are performed by left-multiplying  $N$  by some properly-positioned  $2 \times 2$  matrix and at the same time left-multiplying the "tracking matrix"  $P$  by the same  $2 \times 2$  matrix. Column operations are similar, with left replaced by right and  $P$  by  $Q$ .

```
RowOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[TT.NN]; PP = Simplify[TT.PP];
];
ColOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[NN.TT]; QQ = Simplify[QQ.TT];
];
```

### Swapping Rows and Columns

```
SwapRows[i_, j_] := RowOp[i, j, {{0 1}, {1 0}}];
SwapColumns[i_, j_] := ColOp[i, j, {{0 1}, {1 0}}];
SwapBoth[i_, j_] := {SwapRows[i, j]; SwapColumns[i, j];}
```

### The "GCD" Trick

If  $q = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} q \\ 0 \end{pmatrix}$  allows us to replace pairs of entries in the same column by their greatest common divisor (and a zero!), using invertible row operations. A similar trick works for rows.

? PolynomialExtendedGCD

PolynomialExtendedGCD[poly1, poly2, x] gives the extended GCD of poly1 and poly2 treated as univariate polynomials in x. PolynomialExtendedGCD[poly1, poly2, x, Modulus -> p] gives the extended GCD over the integers mod prime p. >>

```
GCDTrick[{i_, j_}, k_] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[i, k]],
    b = NN[[j, k], x];
  RowOp[i, j, {{s/q, t/q}, {-b/q, a/q}}];
];
GCDTrick[k_, {i_, j_}] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[k, i]],
    b = NN[[k, j], x];
  ColOp[i, j, {{s/q, -b/q}, {t/q, a/q}}];
];
```

### Factoring Diagonal Entries

If  $1 = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} s & a & 1 \\ -tb & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & ab \end{pmatrix} \begin{pmatrix} a & -b \\ t & s \end{pmatrix} = \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$  is an invertible row-column-operations proof of the isomorphism  $\frac{R}{(a)} \oplus \frac{R}{(b)} \cong \frac{R}{(ab)}$ .

```
SplitToSum[i_, j_, a_, b_] := Module[
  {q, s, t, T1, T2},
  {q, {s, t}} = PolynomialExtendedGCD[a, b, x];
  If[q == 1,
    RowOp[i, j, {{s a, 1}, {-t b, 1}}]; ColOp[i, j, {{a, -b}, {t, s}}];
  ]
];
```

### The Jordan Trick

A repeated application of the identity

$$\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{-k} & 0 \\ 1 & p \end{pmatrix} \text{ will bring a matrix like } \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^k \end{pmatrix}$$

to the "Jordan" form of  $\begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}$ , using invertible row and column operations.

```
JordanTrick[i_, j_, p_, s_] :=
  {RowOp[i, j, {{p^{s-1}, -1}, {1, 0}}]; ColOp[i, j, {{1, p}, {0, 1}}];
```

along with:

# Running the JCF Programs

```
In[2]:= SetDirectory["C:\\drorbn\\AcademicPensieve\\Classes\\14-1100"];
<< JCF-Program.m
```

Matrix I - 3x3, 3 eigenvalues.

```
In[4]:= n = 3; AA =  $\begin{pmatrix} 3 & 0 & 0 \\ 4 & -2 & -6 \\ -2 & 0 & 1 \end{pmatrix}$ ;
PP = QQ = II = IdentityMatrix[n];
MM = x II - AA;
NN = PP.MM.QQ;
```

done to  $\begin{pmatrix} 1 & & \\ & 1 & \\ & & (x) \end{pmatrix}$

0 0 0 0

$$\begin{array}{c|c}
 \text{Recovering } C \text{ from } P? & \\
 \hline
 \begin{array}{ccc}
 R^n \xrightarrow[M]{Ix-A} R^n \xrightarrow{T_A} F^n & & \\
 \uparrow Q & \downarrow P & \downarrow C \\
 R^n \xrightarrow[N]{Ix-B} R^n \xrightarrow{T_B} F^n & & 
 \end{array} & 
 \begin{array}{l}
 C e_i = T_B(P e_i) \\
 = T_B(\sum x^k p_k e_i) \\
 = \sum x^k T_B(p_k e_i) \\
 = \sum B^k p_k e_i \\
 \Rightarrow C = \sum B^k p_k \quad \dots \text{complete run 1}
 \end{array}
 \end{array}$$

Go through run 2 until stuck, then

The "Jordan Trick":  $R\langle p^s \rangle = \langle x \rangle / p^s x = 0$   
 $x_0 = x$   
 $x_1 = -p x$   
 $x_2 = p^2 x$   
 $= \langle x_0, \dots, x_{s-1} \rangle / \begin{array}{l} p x_i + x_{i+1} = 0 \\ p x_{s-1} = 0 \end{array}$

so  $(p^s) \sim \begin{pmatrix} p & & \\ & p & \\ & & p \end{pmatrix} \sim \begin{pmatrix} 1 & & \\ & 1 & \\ & & p^s \end{pmatrix} \sim \begin{pmatrix} p & & \\ & p & \\ & & p \end{pmatrix}$

more precisely:

Explicitly:

A repeated application of the identity  $\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{1-k} & 0 \\ 1 & p \end{pmatrix}$  will bring a matrix like

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix} \text{ to the "Jordan" form of } \begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}, \text{ using invertible row and column operations.}$$

`JordanTrick[i_, j_, p_, s_] := {RowOp[i, j,  $\begin{pmatrix} p^{s-1} & -1 \\ 1 & 0 \end{pmatrix}$ ], ColOp[i, j,  $\begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix}]}$`

Then go through the rest of run 2 & through run 3 . . . .

# The Jordan Trick

November-20-14 10:04 AM

$$R \langle p^s \rangle = \langle x \rangle / p^s x = 0$$

$$x_0 = x$$

$$= \langle x_0 \dots x_{s-1} \rangle / \begin{matrix} p x_i + x_{i+1} = 0 \\ p x_{s-1} = 0 \end{matrix}$$

$$x_1 = -p x$$

$$x_2 = p^2 x$$

so  $(p^s) \sim \begin{pmatrix} p & & & \\ 1 & p & & \\ & 1 & p & \\ & & 1 & p \end{pmatrix} \sim \begin{pmatrix} 1 & & & \\ & 1 & & \\ & & \ddots & \\ & & & p^s \end{pmatrix} \sim \begin{pmatrix} p & & & \\ 1 & p & & \\ & 1 & p & \\ & & 1 & p \\ & & & 1 & p \end{pmatrix}$

more precisely:

course evals: 2/17

Riddles as in Nov24Riddles.nb

$$\begin{array}{ccc}
 R^n & \xrightarrow[\mathcal{M}]{I\mathcal{X}-A} & R^n \xrightarrow{T_A} F^n \\
 \uparrow Q & & \downarrow P \\
 R^n & \xrightarrow[\mathcal{N}]{I\mathcal{X}-B} & R^n \xrightarrow{T_B} F^n
 \end{array}$$

Finish last week's material:  
Go over handout along with "run 1"

Dror Bar-Natan: Classes: 2014-15: Math 1100 Algebra I:

## JCF Tricks and Programs

## Row and Column Operations

Row operations are performed by left-multiplying  $N$  by some properly-positioned  $2 \times 2$  matrix and at the same time left-multiplying the "tracking matrix"  $P$  by the same  $2 \times 2$  matrix. Column operations are similar, with left replaced by right and  $P$  by  $Q$ .

```

RowOp[i_, j_, mat_] := Module[{TT = II},
  TT[[i, j], {i, j}] = mat;
  NN = Simplify[TT.NN]; PP = Simplify[TT.PP];
];
ColOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[NN.TT]; QQ = Simplify[QQ.TT];
];

```

## Swapping Rows and Columns

```

SwapRows[i_, j_] := RowOp[i, j,  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ];
SwapColumns[i_, j_] := ColOp[i, j,  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ];
SwapBoth[i_, j_] := {SwapRows[i, j]; SwapColumns[i, j]};

```

## The "GCD" Trick

If  $q = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} q \\ 0 \end{pmatrix}$  allows us to replace pairs of entries in the same column by their greatest common divisor (and a zero!), using invertible row operations. A similar trick works for rows.

? PolynomialExtendedGCD

PolynomialExtendedGCD[pol1, pol2, x] gives the extended GCD of pol1 and pol2 treated as univariate polynomials in x.  
PolynomialExtendedGCD[pol1, pol2, x, Modulus -> p] gives the extended GCD over the integers mod prime p. >>

along with:

```

GCDTrick[{i_, j_}, k_] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[i, k],
    b = NN[[j, k], x];
  RowOp[i, j,  $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix}$ ];
];
GCDTrick[k_, {i_, j_}] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[k, i],
    b = NN[[k, j], x];
  ColOp[i, j,  $\begin{pmatrix} s & -b/q \\ t & a/q \end{pmatrix}$ ];
];

```

## Factoring Diagonal Entries

If  $1 = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} sa & 1 \\ -tb & 1 \end{pmatrix} \begin{pmatrix} a & -b \\ 0 & ab \end{pmatrix} \begin{pmatrix} t & s \\ 0 & b \end{pmatrix} = \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$  is an invertible row-column-operations proof of the isomorphism  $\frac{R}{(a)} \oplus \frac{R}{(b)} \cong \frac{R}{(ab)}$ .

```

SplitToSum[i_, j_, a_, b_] := Module[
  {q, s, t, T1, T2},
  {q, {s, t}} = PolynomialExtendedGCD[a, b, x];
  If[q == 1,
    RowOp[i, j,  $\begin{pmatrix} sa & 1 \\ -tb & 1 \end{pmatrix}$ ]; ColOp[i, j,  $\begin{pmatrix} a & -b \\ t & s \end{pmatrix}$ ];
  ];
];

```

## The Jordan Trick

A repeated application of the identity

$$\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{-1-k} & 0 \\ 1 & p \end{pmatrix}$$

will bring a matrix like  $\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix}$

to the "Jordan" form of  $\begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}$ , using invertible row and column operations.

```

JordanTrick[i_, j_, p_, s_] :=
  {RowOp[i, j,  $\begin{pmatrix} p^{s+1} & -1 \\ 1 & 0 \end{pmatrix}$ ]; ColOp[i, j,  $\begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix}$ ];
};

```

# Running the JCF Programs

```
In[2]:= SetDirectory["C:\\drorbn\\AcademicPensieve\\Classes\\14-1100"];
<< JCF-Program.m
```

## Matrix I - 3x3, 3 eigenvalues.

```
In[4]:= n = 3; AA = {{3, 0, 0},
                     {4, -2, -6},
                     {-2, 0, 1}};
PP = QQ = II = IdentityMatrix[n];
MM = x II - AA;
NN = PP.MM.QQ;
```

0 0 0 0

$$\begin{array}{c|c}
 \text{Recovering } C \text{ from } P? & \\
 \hline
 \begin{array}{ccc}
 R^n \xrightarrow[\text{M}]{Ix-A} R^n \xrightarrow{\pi_A} F^n & & \\
 \uparrow Q & \downarrow P & \downarrow C \\
 R^n \xrightarrow[\text{N}]{Ix-B} R^n \xrightarrow{\pi_B} F^n & & 
 \end{array}
 \end{array}
 \left|
 \begin{array}{l}
 C e_i = \pi_B(P e_i) \\
 = \pi_B(\sum x^k P_k e_i) \\
 = \sum x^k \pi_B(P_k e_i) \\
 = \sum B^k P_k e_i
 \end{array}
 \right.$$

$$\Rightarrow C = \sum B^k P_k \quad \dots \text{complete run 1}$$

Go through run 2 until stuck, then

The "Jordan Trick":  $R\langle p^s \rangle = \langle x \rangle / p^s x = 0$   
 $x_0 = x$   
 $x_1 = -p x$   
 $x_2 = p^2 x$   
 $= \langle x_0 \dots x_{s-1} \rangle / p x_i + x_{i+1} = 0$   
 $p x_{s-1} = 0$

so  $(p^s) \sim \begin{pmatrix} p & & \\ & p & \\ & & p \end{pmatrix} \sim \begin{pmatrix} 1 & & \\ & 1 & \\ & & p^s \end{pmatrix} \sim \begin{pmatrix} p & & \\ & p & \\ & & p \end{pmatrix}$

more precisely:

Explicitly:

A repeated application of the identity  $\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{-1+k} & 0 \\ 1 & p \end{pmatrix}$  will bring a matrix like

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix} \text{ to the "Jordan" form of } \begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}, \text{ using invertible row and column operations.}$$

JordanTrick[i\_, j\_, p\_, s\_] := {RowOp[i, j, {p^{s-1} -1; 1 0}], ColOp[i, j, {1 p; 0 1}]},

Then go through the rest of run 2 & through run 3. . . .

Goal: The "ring" of modules.

Recall that  $(R\text{-mod}, \oplus)$  is an "Abelian group" (really, an Abelian semi-group, and even this is not precise)

Tensor Products. Given  $M, N$

Definition A "tensor product"  $M \otimes N$  is a module  $M \otimes N$  along with a bilinear  $\gamma: M \times N \rightarrow M \otimes N$  s.t.

$$\begin{array}{ccc} M \times N & \xrightarrow[\text{bilinear}]{\gamma} & M \otimes N \\ & \searrow \text{bilinear} & \\ & p \in \text{---} \exists! \text{ linear} & \end{array}$$

Thm  $M \otimes N$  exists & is unique up to isomorphism.

pf First uniqueness, Then

$$M \otimes_R N := \left\{ \sum_{i=1}^n a_i (m_i \otimes n_i) : n_i \in N, a_i \in R \right\} / \begin{array}{l} (am) \otimes n = a(m \otimes n) = m \otimes (an) \\ (m_1 + m_2) \otimes n = \dots \\ m \otimes (n_1 + n_2) = \dots \end{array}$$

$\uparrow$  bilinear  
 $M \times N$

Example.  $\dim V \otimes W = (\dim V)(\dim W)$

Example. If  $q \in \gcd(a, b)$ ,  $\frac{R}{\langle a \rangle} \otimes \frac{R}{\langle b \rangle} \simeq \frac{R}{\langle q \rangle}$

$$\begin{array}{ll} \text{pf. } [r_1]_a \otimes [r_2]_b \longrightarrow [r_1 \cdot r_2]_q & [q] \otimes [1] = [a+tb] \otimes [1] = 0 \\ [r]_a \longrightarrow [r]_a \otimes [1]_b & [r_1 r_2] \otimes [1] = [r_1] [r_2] \end{array}$$

example.  $\gamma: F(X) \otimes F(Y) \rightarrow F(X \times Y)$

1. Always injective! [not so] [easy!]
2. Isomorphism if  $X$  or  $Y$  are finite.
3. Not surjective if  $R = \mathbb{Z}$ ,  $X, Y$  are infinite.  
[not at all obvious!]

Theorem.  $(R\text{-mod}, \oplus, \otimes, 0, R)$  is a "ring".

Theorem.  $(M, N) \mapsto M \otimes N$  is a "bifunctor".

Return HW4!

Course evals: 2/17. Vote and warn others!

Definition A "tensor product"  $M \otimes N$  is a module  $M \otimes N$  along with a bilinear  $\gamma: M \times N \rightarrow M \otimes N$  s.t.

$$\begin{array}{ccc} M \times N & \xrightarrow[\text{bilinear}]{\gamma} & M \otimes N \\ & \searrow \text{bilinear} & \\ & \rho \in \text{---} \exists! \text{ linear} & \end{array}$$

Thm  $M \otimes N$  exists & is unique up to isomorphism. today

Example.  $\dim V \otimes W = (\dim V)(\dim W)$

Proof of uniqueness.

Example. If  $q \in \langle a, b \rangle$ ,  $\frac{R}{\langle a \rangle} \otimes \frac{R}{\langle b \rangle} \cong \frac{R}{\langle q \rangle}$   
 $q = sa + tb$

$$\begin{array}{ll} \text{pf. } [r_1]_a \otimes [r_2]_b \rightarrow [r_1 \cdot r_2]_q & [q] \otimes [1] = [sa + tb] \otimes [1] = 0 \\ [r]_a \rightarrow [r]_a \otimes [1]_b & [r_1 r_2] \otimes [1] = [r_1] [r_2] \end{array}$$

example.  $\gamma: F(X) \otimes F(Y) \rightarrow F(X * Y)$

1. Always injective! [not so easy!]
2. Isomorphism if  $X$  or  $Y$  are finite.
3. Not surjective if  $R = \mathbb{Z}$ ,  $X, Y$  are infinite.  
[not at all obvious!]

Theorem.  $(R\text{-mod}, \otimes, \otimes, 0, R)$  is a "ring".

Theorem.  $(M, N) \mapsto M \otimes N$  is a "bifunctor".

Example.  $\mathbb{Q} \otimes_{\mathbb{Z}} \mathbb{Z}^n \cong \mathbb{Q}^n$  "Extension of scalars".  
 $\leftarrow \text{an } \mathbb{Q}\text{-module}$

done line

Example.  $\mathbb{Q} \otimes_{\mathbb{Z}} \mathbb{Z}^n \cong \mathbb{Q}^n$  "Extension of scalars".  
 $\leftarrow$  a  $\mathbb{Q}$ -module!

In general, given  $\phi: R \rightarrow S$  a ring morphism,  $S$  is an  $R$  module & set  $M_S := S \otimes_R M$ . Then  $M_S$  is an  $S$ -module and  $R_S^n = S^n$ .

Prop. For any domain  $R$  there is a unique Field  $\mathbb{Q}(R)$

s.t.  $R \xrightarrow{\iota} \mathbb{Q}(R)$  "The Field of Fractions"  
 $\searrow \exists!$   
 $F$  Proof later.

Claim IF  $M$  is torsion  $\left[ \forall m \in M \exists r \in R \setminus \{0\} : rm = 0 \right]$  then  $M_{\mathbb{Q}(R)} = 0$ .

$$a \otimes m = r \left( \frac{a}{r} \otimes m \right) = \frac{a}{r} \otimes rm = 0$$

Prop IF  $M \cong R^K \oplus \bigoplus R/\langle p_i, s_i \rangle$ , then

1.  $\dim_{\mathbb{Q}(R)} M_{\mathbb{Q}(R)} = K$

2.  $\dim_{R/\langle p \rangle} M_{R/\langle p \rangle} = K + |\{i : p_i \sim p\}|$

3.  $\dim_{R/\langle p \rangle} \text{im}(m \mapsto p^s m)_{R/\langle p \rangle} = K + |\{i : p_i \sim p \ \& \ s < s_i\}|$

$R/\langle p \rangle$  is a Field  
 because in a PID  
 $\langle p \rangle$  is maximal

$$\text{as } \text{im}(m \mapsto p^s m) \cong \begin{cases} p^s R \cong R & \text{on } R \\ R/\langle q^t \rangle & \text{on } R/\langle q^t \rangle \ q \neq p \\ 0 & \text{on } R/\langle p^t \rangle \ s \geq t \\ R/\langle p^{t-s} \rangle & \text{on } R/\langle p^t \rangle \ s < t \end{cases}$$

$$\text{and } \ker(m \mapsto p^s m) \cong \begin{cases} 0 & \text{on } R \\ 0 & \text{on } R/\langle q^t \rangle \ q \neq p \\ R/\langle p^t \rangle & \text{on } R/\langle p^t \rangle \ s \geq t \\ R/\langle p^s \rangle & \text{on } R/\langle p^t \rangle \ s < t \end{cases}$$

$R/\langle p^s \rangle \hookrightarrow \ker$  by  $[r]_{p^s} \mapsto [p^{t-s}r]_{p^t}$

So such a decomposition is unique!

Localization & Fields of fractions. Let  $R$  be a commutative domain

Def A multiplicative subset  $S$  of  $R \setminus \{0\}$ . (contains 1, closed under  $\times$ )

Examples  $R \setminus \{0\}$ ,  $R \setminus P$  ( $P$  prime), Powers of  $a \neq 0$ .

Definition  $S^{-1}R = \{ \frac{r}{s} \} / \sim$



$$\frac{1}{s_1} \sim \frac{r_2}{s_2} \text{ if } r_1 s_2 = r_2 s_1$$

$$\left[ \frac{r_1}{s_1} \sim \frac{r_2}{s_2}, \frac{r_2}{s_2} \sim \frac{r_3}{s_3} \Rightarrow r_1 s_2 = r_2 s_1, r_2 s_3 = r_3 s_2 \Rightarrow \right.$$

$$\left. r_1 s_2 s_3 = r_2 s_1 s_3 = s_1 r_3 s_2 \Rightarrow r_1 s_3 = r_3 s_1 \right]$$

$$\frac{r_1}{s_1} + \frac{r_2}{s_2} = \dots$$

$$\frac{r_1}{s_1} \cdot \frac{r_2}{s_2} = \dots$$

$R\text{-}f\text{of}$  - "Field of Fractions  $\mathbb{Q}(K)^\sim$ "

$R\text{-}P$  - "localization at  $P^\sim$ "

$\{2^n\}$  - "dyadic rationals".

$$R \rightarrow S^{-1}R$$

is injective

next class: Wed 1-3 OH 3-4.

Course evals: 4/17 Vote and warn others!

Goal. Uniqueness in the structure thm.

Theorem.  $(R\text{-mod}, \oplus, \otimes, 0, R)$  is a "ring".

Theorem.  $(M, N) \mapsto M \otimes N$  is a "bifunctor".

start line

Example.  $\mathbb{Q} \otimes_{\mathbb{Z}} \mathbb{Z}^n \cong \mathbb{Q}^n$  "Extension of scalars".  
 $\leftarrow \text{a } \mathbb{Q}\text{-module}$

In general, given  $\phi: R \rightarrow S$  a ring morphism,  $S$  is an  $R$  module & set  $M_S := S \otimes_R M$ . Then  $M_S$  is an  $S$ -module and  $R_S^n = S^n$ .

Prop. For any domain  $R$  there is a unique field  $\mathbb{Q}(R)$   
 s.t.  $R \xrightarrow{\iota} \mathbb{Q}(R)$  "the field of fractions"  
 $\searrow \exists!$   
 $F$  proof: later.

Claim IF  $M$  is torsion  $\left[ \forall m \in M \exists r \in R \setminus 0 \text{ s.t. } rm = 0 \right]$  then  $M_{\mathbb{Q}(R)} = 0$ .  
 $a \otimes m = r \left( \frac{a}{r} \otimes m \right) = \frac{a}{r} \otimes rm = 0$

Prop IF  $M \cong R^K \oplus \bigoplus R/\langle p_i, s_i \rangle$ , then

1.  $\dim_{\mathbb{Q}(R)} M_{\mathbb{Q}(R)} = K$
2.  $\dim_{R/\langle p \rangle} M_{R/\langle p \rangle} = K + |\{i : p_i \sim p\}|$
3.  $\dim_{R/\langle p \rangle} \text{im}(m \mapsto p^s m)_{R/\langle p \rangle} = K + |\{i : p_i \sim p \ \& \ s < s_i\}|$

$R/\langle p \rangle$  is a field because in a PID  $\langle p \rangle$  is maximal

$$\text{as } \text{im}(m \mapsto p^s m) \cong \begin{cases} p^s R \cong R & \text{on } R \\ R/\langle q^t \rangle & \text{on } R/\langle q^t \rangle \ q \neq p \\ 0 & \text{on } R/\langle p^t \rangle \ s \geq t \\ R/\langle p^{t-s} \rangle & \text{on } R/\langle p^t \rangle \ s < t \end{cases}$$

$$\text{and } \ker(m \mapsto p^s m) \cong \begin{cases} 0 & \text{on } R \\ 0 & \text{on } R/\langle q^t \rangle \ q \neq p \\ R/\langle p^t \rangle & \text{on } R/\langle p^t \rangle \ s \geq t \\ R/\langle p^s \rangle & \text{on } R/\langle p^t \rangle \ s < t \end{cases}$$

$$\begin{cases} \text{on } R/\langle p^t \rangle \text{ s.t.} \\ R/\langle p^{t+s} \rangle \text{ on } R/\langle p^t \rangle \text{ s.t.} \end{cases}$$

$$\begin{cases} R/\langle p^t \rangle \text{ or } R/\langle p^{t-1} \rangle \text{ s.t.} \\ R/\langle p^s \rangle \text{ on } R/\langle p^t \rangle \text{ s.t.} \\ R/\langle p^s \rangle \mapsto \ker \text{ by } [r]_{p^s} \mapsto [p^{t-s}r]_{p^t} \end{cases}$$

So such a decomposition is unique!

**Localization & Fields of fractions.** Let  $R$  be a commutative

Def A multiplicative subset  $S$  of  $R \setminus \{0\}$ . (contains 1, closed under  $\times$ )

Examples  $R \setminus \{0\}$ ,  $R \setminus P$  ( $P$  prime), Powers of  $a \neq 0$ .

Definition  $S^{-1}R = \left\{ \frac{r}{s} \right\} / \frac{r_1}{s_1} \sim \frac{r_2}{s_2} \text{ if } r_1 s_2 = r_2 s_1$

$$\left[ \frac{r_1}{s_1} \sim \frac{r_2}{s_2}, \frac{r_2}{s_2} \sim \frac{r_3}{s_3} \Rightarrow r_1 s_2 = r_2 s_1, r_2 s_3 = r_3 s_2 \Rightarrow \right. \quad \left. \frac{r_1}{s_1} + \frac{r_2}{s_2} = \dots \right.$$

$$\left. r_1 s_2 s_3 = r_2 s_1 s_3 = s_1 r_3 s_2 \Rightarrow r_1 s_3 = r_3 s_1 \right] \quad \frac{r_1}{s_1} \cdot \frac{r_2}{s_2} = \dots$$

$R \setminus \{0\}$  — "Field of fractions  $\mathbb{Q}(R)$ "

$R \setminus P$  — "localization at  $P$ "

$\{2^n\}$  — "dyadic rationals".

$R \rightarrow S^{-1}R$   
is injective

*all done*

Pam needs volunteers to mark olympiad questions!  
Course evals: 5/17

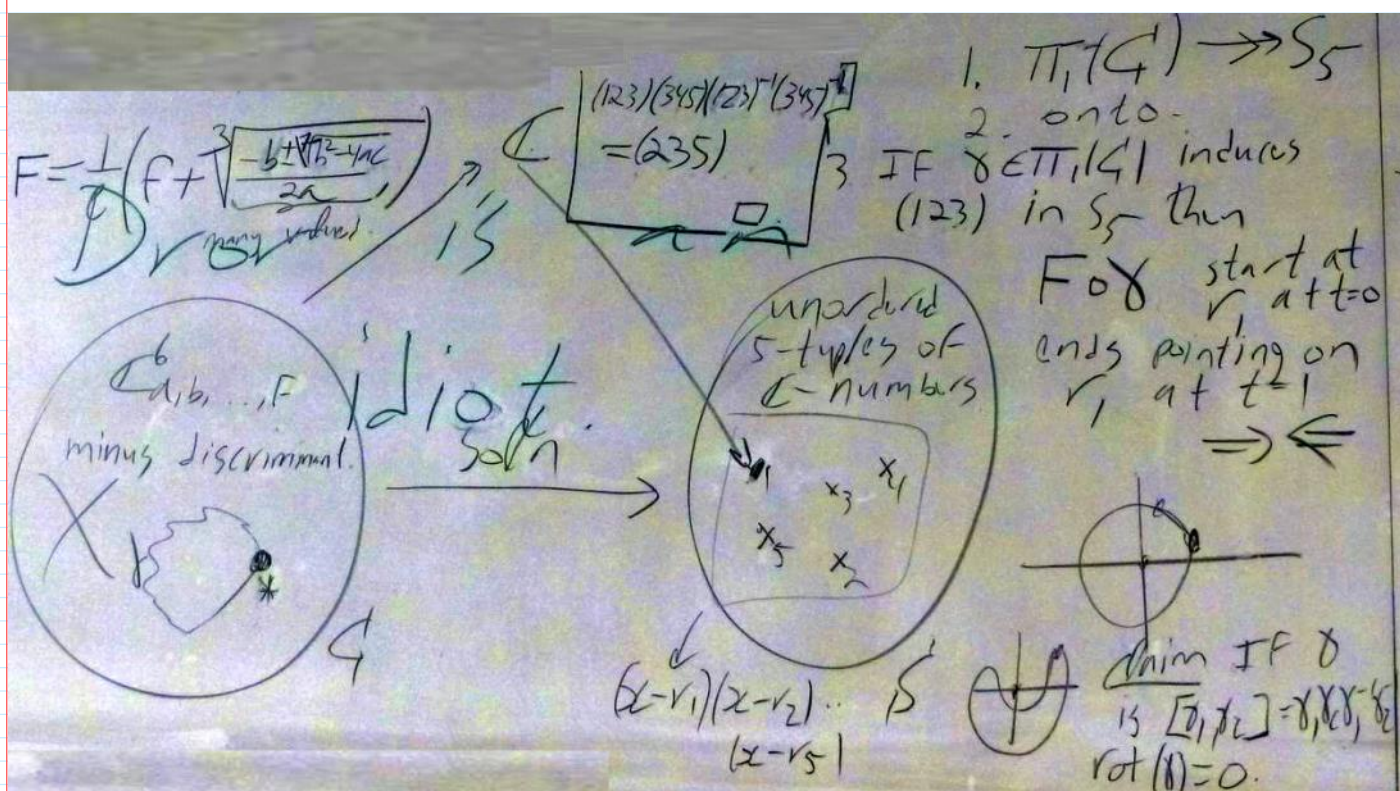
The Final: All is included, same style as term test & as previous years.

The key: Understand EVERYTHING.

Today: Not solving the quintic, more on JCF.

Tomorrow: Riddles session! Baker 6/83, 10 AM.

Following <http://drorbn.net/dbnvp/AKT-140314.php>:



## Some JCF tricks

If  $q = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} q \\ 0 \end{pmatrix}$  allows us to replace pairs of entries in the same column by their greatest common divisor (and a zero!), using invertible row operations. A similar trick works for rows.

If  $1 = \gcd(a, b) = sa + tb$ , the equality  $\begin{pmatrix} sa & 1 \\ -tb & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & ab \end{pmatrix} \begin{pmatrix} a & -b \\ t & s \end{pmatrix} = \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$  is an invertible row-column-operations proof of the isomorphism  $\frac{R}{(a)} \oplus \frac{R}{(b)} \simeq \frac{R}{(ab)}$ .

A repeated application of the identity  $\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{-1+k} & 0 \\ 1 & p \end{pmatrix}$  will bring a matrix like

$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix}$  to the "Jordan" form of  $\begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}$ , using invertible row and column operations.

$$\langle x, y \rangle / \begin{matrix} y=0 \\ p^k x=0 \end{matrix} \cong \langle x, z \rangle / \begin{matrix} p^{k-1}x + z = 0 \\ pz=0 \end{matrix}$$

$$\begin{array}{ccc} y & \xrightarrow{\quad} & p^{k-1}x + z \\ x & \xrightarrow{\quad} & -x \\ -x & \xleftarrow{\quad} & x \\ y + p^{k-1}x & \xleftarrow{\quad} & z \end{array}$$