

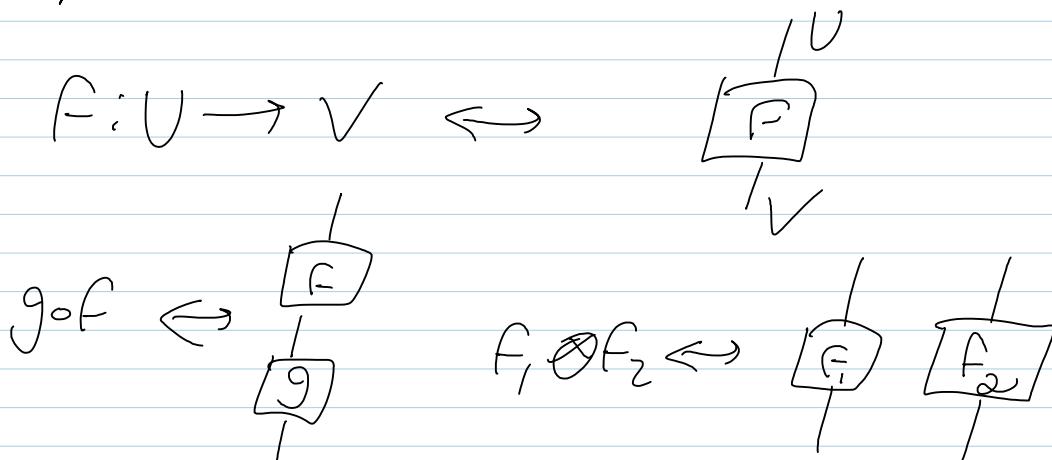
\* Relate "surface tangles" to some algebraic structures [monoidal categories]

$(\mathcal{C}, \otimes, \mathbb{1})$  with  $(X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$

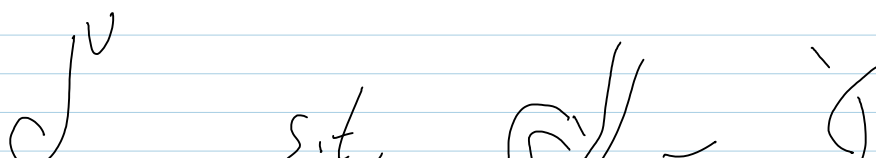
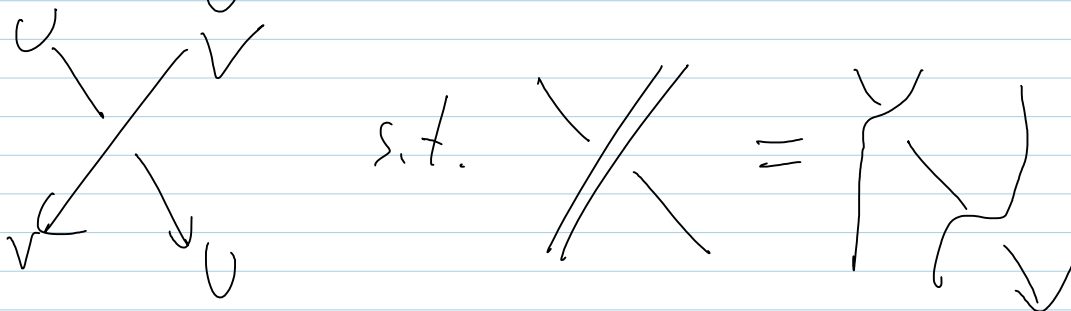
+ duality  $V \rightsquigarrow V^*$

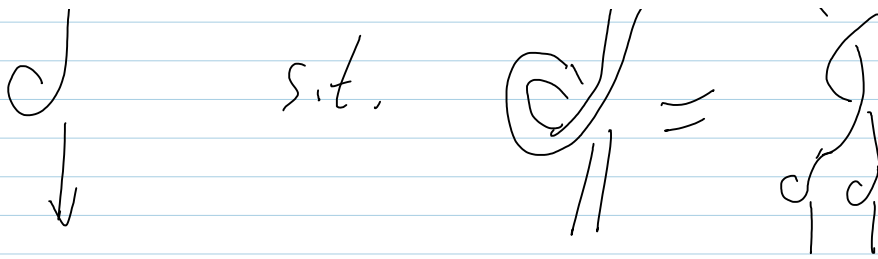
$$\begin{cases} V^* \otimes V \rightarrow \mathbb{1} \\ \mathbb{1} \rightarrow V \otimes V^* \end{cases} \quad \begin{array}{c} V \otimes V^* \\ \curvearrowright \\ V^* \otimes V \end{array}$$

Graphical calculus:

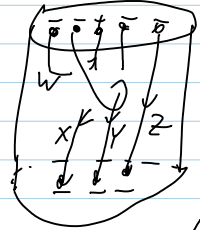


Def  $\mathcal{Z}$  is "ribbon" if it is equipped w/ braiding & twist:





Thm (Resh-Tur, Yetter...)



The Free ...  $T_{\mathcal{S}}$  Framed  
is the category of tangles whose objects  
are coloured w/  $\mathcal{S}$

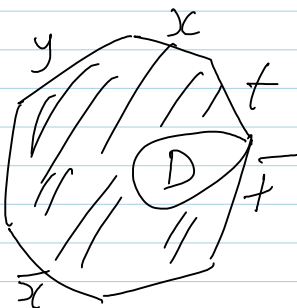
This language used to intimidate me and I'm  
sure it intimidates many still.

The category of tangles on a surface

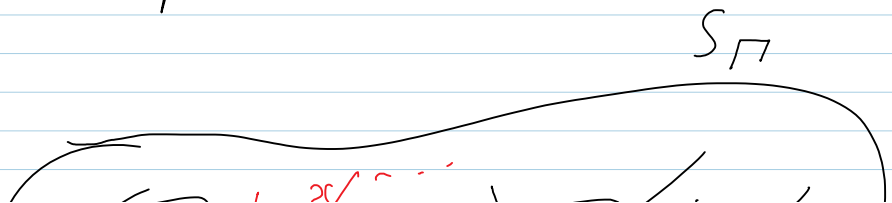
Combinatorial surface:  
an ordered finite set w/ <sup>F.p. free</sup> involution

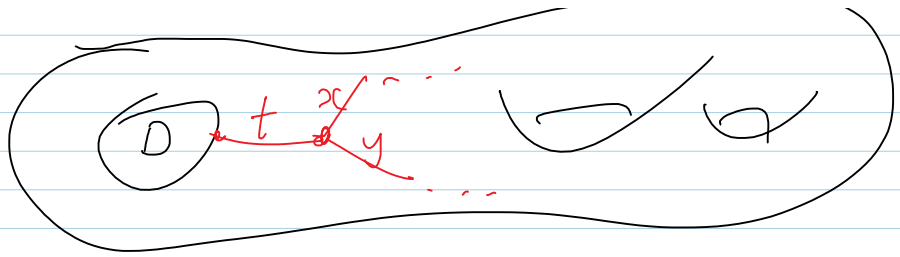
$$(\Gamma, \leq) \ \& \ (x \mapsto \bar{x})$$

s.t. if  $t = \min(\Gamma)$  then  $\bar{t} = \max(\Gamma)$



"a once-bordered fat graph"  
represents:



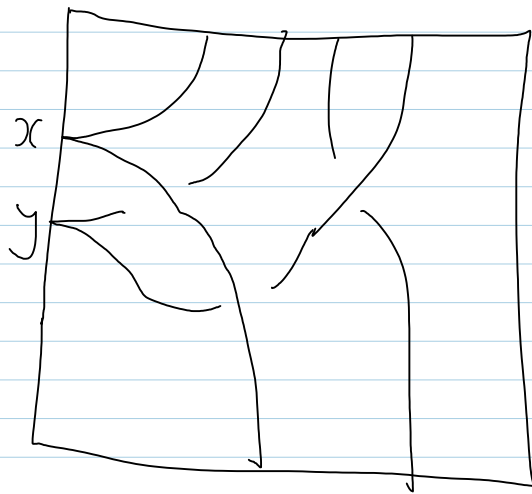


the vertices of the graph are the cycles  
of  $\sigma$ , where  $\sigma(e) = \overline{\text{pred}(e)}$

$T(\Gamma) :=$  tangles on  $S_n \cup \text{Disc}$   
(in  $(S_n \cup D) \times I$ )

There is a functor  $T \rightarrow T(\Gamma)$   
by inclusion of the disk.

Representing surface tangles by "beak diagrams"



where  
 $x, y \in \Gamma$

/ isotopy

obvious composition

$D(\Gamma) :=$  beak diagrams category.

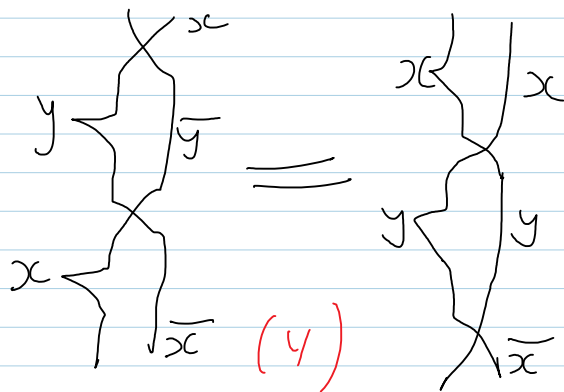
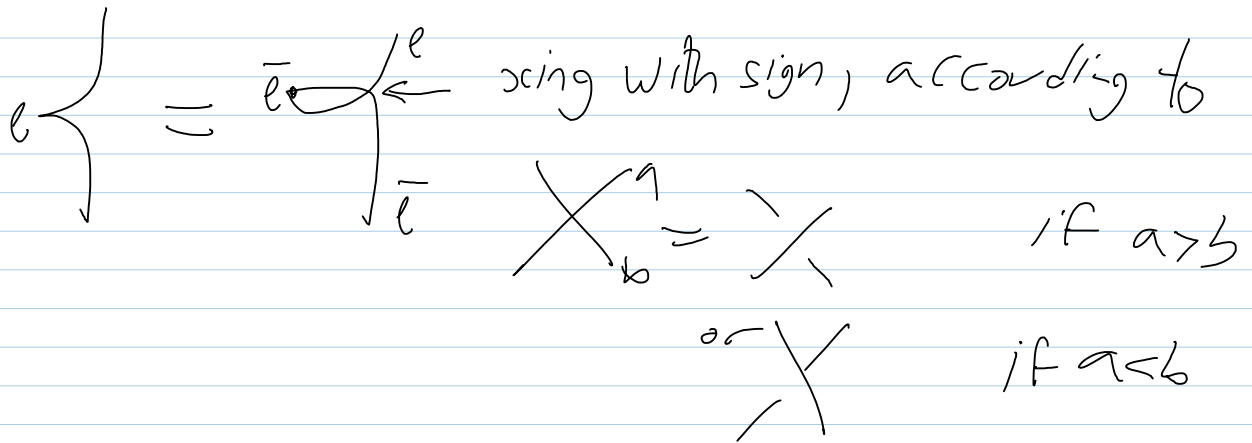
# Functor

$$\varphi: D(\Gamma) \longrightarrow T(\Gamma)$$

by sending an  $x$ -beak to a path dual to  $x$ . We want to describe  $\varphi$ :

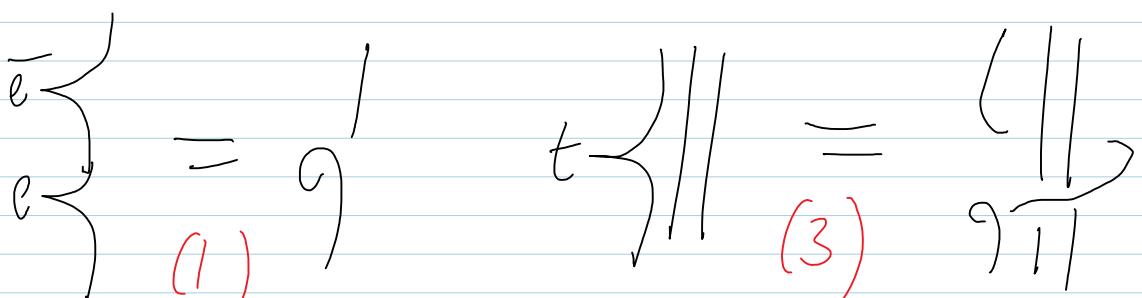
\* usual Reidemeister-moves. (R2 & R3 &  $a^d = 1$ )

\* S-moves (slide)



Same convention.

\* T-moves (topology)



$$\nu \quad (1) \quad | \quad \nu \quad (3) \quad \eta \quad |$$

if  $e_1, \dots, e_k$  is a vertex,  $\left. \begin{matrix} e_k \\ \vdots \\ e_1 \end{matrix} \right\} = \eta \quad |$   
(2)

$$\underline{\text{Thm}} \quad \frac{D(\Gamma)}{RST(\Gamma)} \xrightarrow{\sim} T(\Gamma)$$

Def A  $\Gamma$ -structure on  $\mathcal{C}$  is a functor

$$\mathcal{C} \xrightarrow{e} \mathcal{C}_\Gamma \quad \text{equipped with}$$

$$e \in \Gamma \rightsquigarrow (e) \in \text{Aut}(\{ \cdot \otimes \cdot \})$$

intimidation at its best!

meaning, For any  $X, Y \in \text{obj}(\mathcal{C})$

$$(e)_{X,Y} \in \text{Aut}(\{ X \otimes Y \})$$

(think  $\left. \begin{matrix} e \\ \vdots \\ X \otimes Y \end{matrix} \right\}$ )

s.t. (1)-(3) hold, and

(5):

$$\left( \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \right) \left| \right. = \left( \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \right) \left| \right. \quad \left( \text{w/ xing -info} \right)$$



Poor marketing when it comes to applications.