

colours missing esp. green.

spell check!

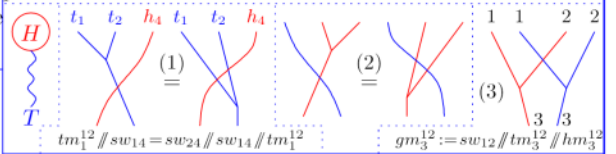
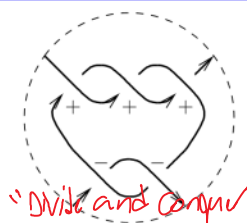
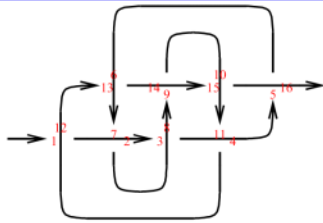
**Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 1**

Dror Bar-Natan at Knots in Washington XXXIV  
<http://www.math.toronto.edu/~drobna/Talks/GWU-1203/>



**Abstract.** A straightforward proposal for a group-theoretic invariant of knots fails if one really means groups, but works once generalized to meta-groups (to be defined). We will construct one complicated but elementary meta-group as a meta-bicrossed-product (to be defined), and explain how the resulting invariant is a not-yet-understood generalization of the Alexander polynomial, while at the same time being a specialization of a somewhat-understood “universal finite type invariant of w-knots” and of an elusive “universal finite type invariant of v-knots”.

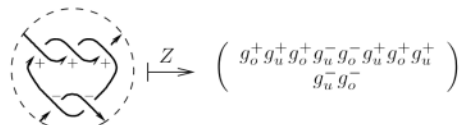
**Bicrossed Products.** If  $G = HT$  is a group presented as a product of two of its subgroups, with  $H \cap T = \{e\}$ , then also  $G = TH$  and  $G$  is determined by  $H, T$ , and the “swap” map  $sw^{th} : (t, h) \mapsto (h', t')$  defined by  $th = h't'$ . The map  $sw$  satisfies (1) and (2) below; conversely, if  $sw : T \times H \rightarrow H \times T$  satisfies (1) and (2) (+ lesser conditions), then (3) defines a group structure on  $H \times T$ , the “bicrossed product”.



A **Meta-Bicrossed-Product** is a collection of sets  $\beta(H, T)$  and operations  $tm_z^{xy}, hm_z^{xy}$  and  $sw^{th}$  (and lesser ones), such that  $tm$  and  $hm$  are “associative” and (1) and (2) hold (+ lesser conditions). A meta-bicrossed-product defines a meta-group with  $G_X := \beta(X, X)$  and  $gm$  as in (3).

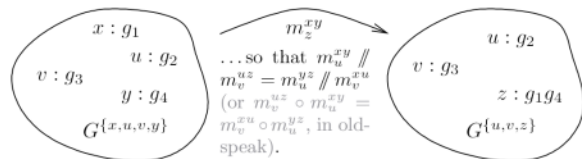


**Idea.** Given a group  $G$  and two pairs  $R^\pm = (g_o^\pm, g_u^\pm) \in G^2$ , map them to crossings and “multiply along”, so that



**This Fails!** R2 implies that  $g_o^+ g_u^- = e$  and then R3 implies that  $g_o^+$  and  $g_u^+$  commute, so the result is a simple counting invariant.

**A Group Computer.** Given  $G$ , can store group elements and perform operations on them:



(Also has  $S_x$  for inversion,  $e_x$  for unit insertion,  $d_x$  for register deletion,  $\Delta_{xy}^z$  for element cloning, and  $(D_1, D_2) \mapsto D_1 \cup D_2$  for merging, and very many obvious composition axioms relating these.)

**A Meta-Group.** Is a similar “computer”, only its internal structure is unknown to us. Namely it is a collection of sets  $\{G_X\}$  indexed by all finite sets  $X$ , and a collection of operations  $m_z^{xy}, S_x, e_x, d_x, \Delta_{xy}^z$ , and  $\cup$ , satisfying the exact same properties.

- Example 1.** The non-meta example,  $G_X := G^X$ .
- Example 2.**  $G_X := M_{X \times X}(\mathbb{Z})$ , with simultaneous row and column operations.

**$\beta$  Calculus.** Let  $\beta(H, T)$  be

$$\left\{ \begin{array}{c|ccc} \omega & h_1 & h_2 & \cdots \\ t_1 & \alpha_{11} & \alpha_{12} & \cdot \\ t_2 & \alpha_{21} & \alpha_{22} & \cdot \\ \vdots & \cdot & \cdot & \cdot \end{array} \middle| \begin{array}{c} h_j \in H, t_i \in T, \text{ and } \omega \text{ and the } \alpha_{ij} \text{ are Laurent polynomials in variables } T_i, \text{ in bijection with the } t_i\text{'s} \end{array} \right\}$$

with operations  $tm_z^{xy} : \begin{array}{c|ccc} \omega & \cdots \\ t_x & \alpha & \beta \\ t_y & \gamma & \gamma \\ \vdots & & \end{array} \mapsto \begin{array}{c|ccc} \omega & \cdots \\ t_z & \alpha + \beta & \gamma \end{array}$ ,

$hm_z^{xy} : \begin{array}{c|ccc} \omega & h_x & h_y & \cdots \\ \vdots & \alpha & \beta & \gamma \end{array} \mapsto \begin{array}{c|ccc} \omega & h_z & \cdots \\ \vdots & \alpha + \beta + \langle \alpha \rangle \beta & \gamma \end{array}$ ,

$sw_{xy}^{th} : \begin{array}{c|ccc} \omega & h_y & \cdots \\ t_x & \alpha & \beta \\ \vdots & \gamma & \delta \end{array} \mapsto \begin{array}{c|ccc} \omega \epsilon & h_y & \cdots \\ t_x & \alpha(1 + \langle \gamma \rangle / \epsilon) & \beta(1 + \langle \gamma \rangle / \epsilon) \\ \vdots & \gamma / \epsilon & \delta - \gamma \beta / \epsilon \end{array}$ ,

where  $\epsilon := 1 + \alpha$ ,  $\langle \alpha \rangle := \sum_i \alpha_i$ , and  $\langle \gamma \rangle := \sum_{i \neq x} \gamma_i$ , and let

$$R_{xy}^p := \begin{array}{c|cc} 1 & h_x & h_y \\ t_x & 0 & T_x - 1 \\ t_y & 0 & 0 \end{array} \quad R_{xy}^m := \begin{array}{c|cc} 1 & h_x & h_y \\ t_x & 0 & T_x^{-1} - 1 \\ t_y & 0 & 0 \end{array}$$

**Theorem.**  $Z^\beta$  is a tangle invariant (and much more). Restricted to knots, the  $\omega$  part is the Alexander polynomial. Restricted to links, it contains the multivariable Alexander polynomial. Restricted to braids, it is equivalent to the Burau representation.

**Why Happy?** • Applications to w-knots. • Everything that I know about the Alexander polynomial can be expressed cleanly in this language (even if without proof), except HF, but including genus, ribboness, cabling, v-knots, knotted graphs, etc., and there’s potential for vast generalizations.

\* Fits on one sheet, including implementation

# Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 2

```
I mean business! << Utilities.m
The key implementation trick is the bijection
      ω | h_j
      t_i | α_ij  ↔ B(ω, ∑ α_ij t_i h_j) :
⟨μ_⟩ := μ / . t_ → 1;
tm_x,y → z [β_] := β / . {t_x|y → t_z, T_x|y → T_z};
hm_x,y → z [B[ω_, Δ_]] := Module[
  {α = D[Δ, h_x], β = D[Δ, h_y], γ = Δ / . h_x|y → 0},
  B[ω, (α + (1 + ⟨α⟩) β) h_x + γ] // βCollect];
sw_x,y [B[ω_, Δ_]] := Module[α, β, γ, δ, ε,
  α = Coefficient[Δ, h_y t_x]; β = D[Δ, t_x] / . h_y → 0;
  γ = D[Δ, h_y] / . t_x → 0; δ = Δ / . h_y | t_x → 0;
  ε = 1 + α;
  B[ω * ε, α (1 + ⟨γ⟩ / ε) h_y t_x + β (1 + ⟨γ⟩ / ε) t_x
    + γ / ε h_y + δ - γ * β / ε
  ] // βCollect];
gm_x,y → z [β_] := β // sw_x,y // hm_x,y → z // tm_x,y → z;
B / : B[ω1_, Δ1_] B[ω2_, Δ2_] := B[ω1 * ω2, Δ1 + Δ2];
Rp_x,y_ := B[1, (T_x - 1) t_x h_y];
Rm_x,y_ := B[1, (T_x^-1 - 1) t_x h_y];
```

```
{β = B[ω, Sum[α10,1,3 t1 h3, {i, {1, 2, 3}}, {j, {4, 5}}]},
  β // tm1,2 → 1 // sw1,4,
  β // sw2,4 // sw1,4 // tm1,2 → 1
} // ColumnForm
Some testing...
(ω h4 h5
 t1 α14 α15
 t2 α24 α25
 t3 α34 α35)
(ω (1 + α14 + α24) h4 h5
 t1 (α14 + α24) (1 - α14 - α24) (1 - α14 - α24)
 t2 α34 (1 - α14 - α24) -α15 α34 - α25 α34 - α14 α35 - α24 α35
 t3 α34 (1 - α14 - α24) (1 - α14 - α24)
(ω (1 + α14 + α24) h4 h5
 t1 (α14 + α24) (1 - α14 - α24) (1 - α14 - α24)
 t2 α34 (1 - α14 - α24) -α15 α34 - α25 α34 - α14 α35 - α24 α35
 t3 α34 (1 - α14 - α24) (1 - α14 - α24)
```

```
{Rm5,1 Rm6,2 Rp3,4 // gm1,4 → 1 // gm2,5 → 2 // gm3,6 → 3,
  Rp6,1 Rm2,4 Rm3,5 // gm1,4 → 1 // gm2,5 → 2 // gm3,6 → 3}
{ 1 h1 h2 | 1 h1 h2
 t2 -1/T2 0 | t2 -1/T2 0
 t3 -1/T3 -1/T3 | t3 -1/T3 -1/T3 }
```

```
β = Rm12,1 Rm2,7 Rm8,3 Rm4,11 Rp16,5 Rp6,13 Rp14,9 Rp10,15
1 h1 h3 h5 h7 h9 h11 h13 h15
t2 0 0 0 -1/T2 0 0 0 0
t4 0 0 0 0 0 0 -1/T4 0 0
t6 0 0 0 0 0 0 0 -1 + T6 0
t8 0 -1/T8 0 0 0 0 0 0 0
t10 0 0 0 0 0 0 0 0 -1 + T10
t12 -1/T12 0 0 0 0 0 0 0 0
t14 0 0 0 0 0 -1 + T14 0 0 0
t16 0 0 -1 + T16 0 0 0 0 0 0
```

Propaganda

```
Do[β = β // gm1,k → 1, {k, 2, 10}]; β
817, cont.
h1 h11 h13 h15
t1 (-1-T1) T14 (T1-T14) (-1-T1) (1-T1-T14) T14 T16 (-1-T1) (1-T1-T14) T14 -1 + T1
t12 (-1-T1) (T1-T14) T16 0 0 0 0 0 0
t14 (-1-T14) (-1-T1-T14) (-1-T1) (1-T1-T14) T16 (-1-T1) (1-T1-T14) (-1-T14) 0
t16 (-1-T1) T1 (-1-T14) (-1-T1) T1 (-1-T14) (-1-T1) T1 (-1-T14) 0
Do[β = β // gm1,k → 1, {k, 11, 16}]; β
(-1-4 T1 + 8 T1^2 - 11 T1^3 + 8 T1^4 - 4 T1^5 + T1^6) h1
t1 0
<< KnotTheory
Alexander[Knot[8, 17]][T1] // Factor
Loading KnotTheory` version of August 22, 2010, 13:36:57.55.
Read more at http://katlas.org/wiki/KnotTheory.
KnotTheory:loading: Loading precomputed data in PD4Knots.
1-4 T1 + 8 T1^2 - 11 T1^3 + 8 T1^4 - 4 T1^5 + T1^6
T1^3
```

Where does it come from? The accidental<sup>1</sup> answer is that it is a symbolic calculus for a natural reduction<sup>4</sup> of the unique homomorphic expansion<sup>2</sup> of w-tangles<sup>3</sup>.

1. "Accidental" for it's only how I came about it. There ought to be a better answer.
2. A "homomorphic expansion", aka as a homomorphic universal finite type invariant, is a completely canonical construct whose presence implies that the objects in questions are susceptible to study using graded algebra.
3. "v-Tangles" are the meta-group generated by crossings modulo Reidemeister moves. "w-Tangles" are a natural quotient of v-tangles. They are at least related and perhaps identical to a certain class of 1D/2D knots in 4D.
4. To "only what is visible by the 2D Lie algebra".

A certain generalization will arise by not reducing as in 4. A vast generalization may arise when homomorphic expansions for v-tangles are understood, a task likely equivalent to the Etingof-Kazhdan quantization of Lie bialgebras.

### A Partial To Do List.

1. Where does it more simply come from?
2. Remove all the denominators.
3. How do determinants arise in this context (×2)?
4. Understand links.
5. Find the "reality condition".
6. Do some "Algebraic Knot Theory".
7. Understand antipodes.
8. Categorify.

Montpellier w-pictures