

Shifted Partial Quadratics, their Pushforwards, and Signature Invariants for Tangles

Abstract. Following a general discussion of the computation of zombians of unfinished columbaria (with examples), I will tell you about my recent joint work w/ Jessica Liu on what we feel is the “textbook” extension of knot signatures to tangles, which for unknown reasons, is not in any of the textbooks that we know.



Jessica Liu



Columbaria in an East Sydney Cemetery



Jacobian, Hamiltonian, Zombian

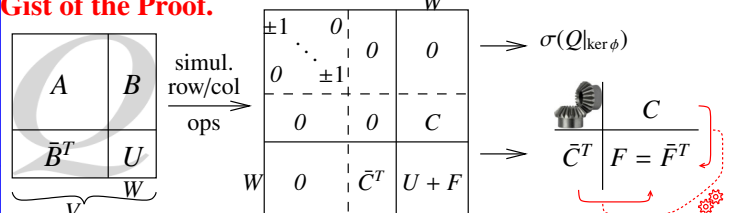
Kashaev’s Conjecture [Ka] For links, $\sigma_{Kas} = 2\sigma_{TL}$.
Liu’s Theorem [Li].

A **Partial Quadratic (PQ)** on V is a quadratic Q defined only on a subspace $\mathcal{D}_Q \subset V$. We add PQs with $\mathcal{D}_{Q_1+Q_2} := \mathcal{D}_{Q_1} \cap \mathcal{D}_{Q_2}$. Given a linear $\psi: V \rightarrow W$ and a PQ Q on W , there is an obvious **pullback** ψ^*Q , a PQ on V .

Theorem 1. Given a linear $\phi: V \rightarrow W$ and a PQ Q on V , there is a unique **pushforward** PQ ϕ_*Q on W such that for every PQU on W , $\sigma_V(Q + \phi^*U) = \sigma_{\ker\phi}(Q|_{\ker\phi}) + \sigma_W(U + \phi_*Q)$. (If you must, $\mathcal{D}(\phi_*Q) = \phi(\text{ann}_Q(\mathcal{D}(Q) \cap \ker\phi))$ and $(\phi_*Q)(w) = Q(v)$, where v is s.t. $\phi(v) = w$ and $Q(v, \text{rad } Q|_{\ker\phi}) = 0$).

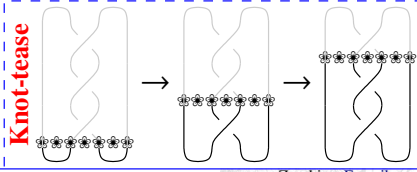
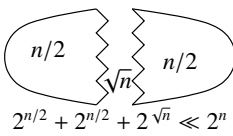
Prior Art on signatures for tangles / braids. Gambaudo and Ghys [GG], Cimasoni and Conway [CC], Conway [Co], Merz [Me]. All define signatures of tangles / braids by first closing them to links and then work hard to derive composition properties.

Gist of the Proof.



Why Tangles? • Faster!

- Conceptually clearer proofs of invariance (and of skein relations).
- Often fun and consequential:
 - The Jones Polynomial \rightsquigarrow The Temperley-Lieb Algebra.
 - Khovanov Homology \rightsquigarrow “Unfinished complexes”, complexes in a category.
 - The Kontsevich Integral \rightsquigarrow Associators.
 - HFK \rightsquigarrow OMG, type D , type A , $\mathcal{A}_\infty, \dots$



Computing Zombians of Unfinished Columbaria.

- Must be no slower than for finished ones.
- Future zombies must be able to complete the computation.
- Future zombies must not even know the size of the task that today’s zombies were facing.
- We must be able to extend to ZPUCs, Zombie Processed Unfinished Columbaria!

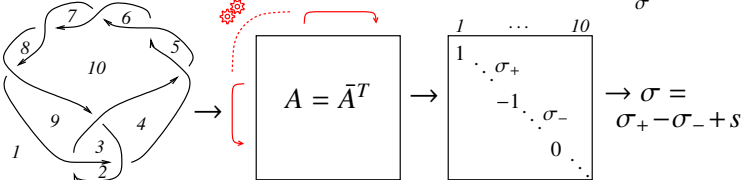


Columbarium near Assen

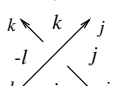
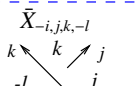
Example / Exercise. Compute the determinant of a $1,000 \times 1,000$ matrix in which 50 entries are not yet given.

Homework / Research Projects. • What with ZPUCs? • Use this to get an Alexander tangle invariant.

Reminders. {links} \rightleftharpoons {matrices / quadratic forms} $\xrightarrow{\text{signature}} \mathbb{Z}$:

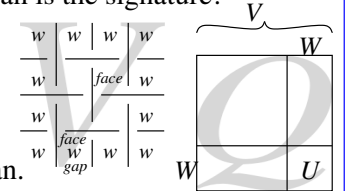


With $|\omega| = 1, t = 1 - \omega, r = t + \bar{t}, v = \text{Re}(\omega),$ and $u = \text{Re}(\omega^{1/2})$:

$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} -r & -t & 2t & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ 2\bar{t} & t & -r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A += \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$s = 0$		$s = -1$
$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} r & -t & -2\bar{t} & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ -2t & t & r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A -= \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$s = 0$		$s = +1$

... and the quadratic $F := \phi_*Q$ is well-defined only on $D := \ker C$. Exactly what we want, if the Zombian is the signature!

- V : The full space of faces.
 - W : The boundary, made of gaps.
 - Q : The known parts.
 - U : The part yet unknown.
 - $\sigma_V(Q + \phi^*(U))$: The overall Zombian.
 - $\sigma(Q|_{\ker\phi})$: An internal bit. $U + \phi_*Q$: A boundary bit.
- And so our ZPUC is the pair $S = (\sigma(Q|_{\ker\phi}), \phi_*Q)$.



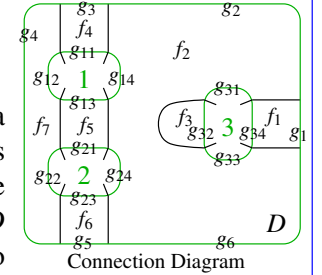
A **Shifted Partial Quadratic (SPQ)** on V is a pair $S = (s \in \mathbb{Z}, Q \text{ a PQ on } V)$. addition also adds the shifts, pullbacks keep the shifts, yet $\phi_*S := (s + \sigma_{\ker\phi}(Q|_{\ker\phi}), \phi_*Q)$ and $\sigma(S) := s + \sigma(Q)$.

Theorem 1' (Reciprocity). Given $\phi: V \rightarrow W$, for SPQs S on V and U on W we have $\sigma_V(S + \phi^*U) = \sigma_W(U + \phi_*S)$ (and this characterizes ϕ_*S). Note. ψ^* is additive but ϕ_* is not.

Theorem 2. ψ^* and ϕ_* are functorial.
Theorem 3. “The pullback of a pushforward scene is a pushforward scene”: If, on the right, β and δ are arbitrary, $Y = EQ(\beta, \gamma) = V \oplus_Z W = \{(v, w) : \beta v = \gamma w\}$ and μ and ν are the obvious projections, then $\gamma^*\beta_* = \nu_*\mu^*$.

Definition. $S \left(\begin{matrix} g_2 \\ g_3 \dots g_1 \end{matrix} \right) := \left\{ \begin{matrix} \text{SPQ } S \\ \text{on } \langle g_i \rangle \end{matrix} \right\}$.

Theorem 4. $\{S(\text{cyclic sets})\}$ is a planar algebra, with compositions $S(D)((S_i)) := \phi_*^D(\psi_D^*(\bigoplus_i S_i))$, where $\psi_D: \langle f_i \rangle \rightarrow \langle g_{\alpha i} \rangle$ maps every face of D to the sum of the input gaps adjacent to it and $\phi^D: \langle f_i \rangle \rightarrow \langle g_i \rangle$ maps every face to the sum of the output gaps adjacent to it. So for our D , $\psi_D: f_1 \mapsto g_{34}, f_2 \mapsto g_{31} + g_{14} + g_{24} + g_{33}, f_3 \mapsto g_{32}, f_4 \mapsto g_{11}, f_5 \mapsto g_{13} + g_{21}, f_6 \mapsto g_{23}, f_7 \mapsto g_{12} + g_{22}$ and $\phi^D: f_1 \mapsto g_1, f_2 \mapsto g_2 + g_6, f_3 \mapsto 0, f_4 \mapsto g_3, f_5 \mapsto 0, f_6 \mapsto g_5, f_7 \mapsto g_4$.



Theorem 5. TL and Kas , defined on X and \bar{X} as before, extend to planar algebra morphisms $\{\text{tangles}\} \rightarrow \{S\}$. Restricted to links, $TL = \sigma_{TL}$ and $Kas = \sigma_{Kas}$.



Levine, Tristram, Kashaev

Implementation (sources: <http://drorbn.net/icerm23/ap>). I like it most when the implementation matches the math perfectly. We failed here.

Once[<< KnotTheory`];

Loading KnotTheory` version

of February 2, 2020, 10:53:45.2097.

Read more at <http://katlas.org/wiki/KnotTheory>.

Utilities. The step function, algebraic numbers, canonical forms.

$\theta[x_]$ /; NumericQ[x] := UnitStep[x]

```
 $\omega 2[v\_][p\_]$  := Module[{q = Expand[p], n, c},
  If[q == 0, 0,
    c = Coefficient[q,  $\omega$ , n = Exponent[q,  $\omega$ ]];
     $c v^n + \omega 2[v][q - c (\omega + \omega^{-1})^n]$ ];
```

```
sign[ $\mathcal{E}$ _] := Module[{n, d, v, p, rs, e, k},
  {n, d} = NumeratorDenominator[ $\mathcal{E}$ ];
  {n, d} /=  $\omega^{\text{Exponent}[n, \omega]/2 + \text{Exponent}[n, \omega, \text{Min}]/2}$ ;
  p = Factor[ $\omega 2[v] @ n * \omega 2[v] @ d /. v \rightarrow 4 u^2 - 2$ ];
  rs = Solve[p == 0, u, Reals];
  If[rs == {}, Sign[p /. u -> 0],
    rs = Union@{u /. rs};
    Sign[(-1)e=Exponent[p, u] Coefficient[p, u, e]] + Sum[
      k = 0;
      While[{d = RootReduce[ $\partial_{\{u, ++k\}} p /. u \rightarrow r$ ]} == 0];
      If[EvenQ[k], 0, 2 Sign[d]] *  $\theta[u - r]$ ,
      {r, rs}]]
]
```

SetAttributes[B, Orderless];

```
CF[b_B] := RotateLeft[#, First@Ordering[#] - 1] & /@
DeleteCases[b, {}]
```

```
CF[ $\mathcal{E}$ _] := Module[{ $\gamma$ s = Union@Cases[ $\mathcal{E}$ ,  $\gamma$ _ |  $\bar{\gamma}$ _,  $\infty$ ]},
  Total[CoefficientRules[ $\mathcal{E}$ ,  $\gamma$ s] /.
    (ps_ -> c_) := Factor[c]  $\times$  Times@@ $\gamma$ sps]]
```

CF[{}] = {};

CF[C_List] :=

```
Module[{ $\gamma$ s = Union@Cases[C,  $\gamma$ _,  $\infty$ ],  $\gamma$ },
  CF /@ DeleteCases[0] [
    RowReduce[Table[ $\partial_{\gamma} r$ , {r, C}, { $\gamma$ ,  $\gamma$ s}]] .  $\gamma$ s ]
```

(\mathcal{E} _)* := $\mathcal{E} /. \{\bar{\gamma} \rightarrow \gamma, \gamma \rightarrow \bar{\gamma}, \omega \rightarrow \omega^{-1}, c_Complex \rightarrow c^*\}$;

r_Rule* := {r, r*}

RulesOf[γ_i + rest_] := ($\gamma_i \rightarrow -rest$)*;

```
CF[PQ[C_, q_]] := Module[{nC = CF[C]},
  PQ[nC, CF[q /. Union@@RulesOf /@ nC]] ]
```

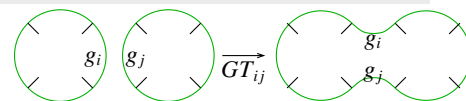
CF[Σ_b [σ _, pq_]] := $\Sigma_{CF[b]}$ [σ , CF[pq]]

Pretty-Printing.

```
Format[ $\Sigma_{b,B}[\sigma_, PQ[C_, q_]]$ ] := Module[{ $\gamma$ s},
   $\gamma$ s =  $\gamma$ # & /@ Join@@b;
  Column[{TraditionalForm@ $\sigma$ ,
    TableForm[Join[
      Prepend[""] /@ Table[TraditionalForm[ $\partial_c r$ ],
        {r, C}, {c,  $\gamma$ s}],
      {Prepend[""] [
        Join@@
          (b /. {L_, m___, r_} =>
            {DisplayForm@RowBox[{"(", L}],
              m, DisplayForm@RowBox[{r, ")"}]}) /-
            i_Integer =>  $\gamma_i$  ]}],
      MapThread[Prepend,
        {Table[TraditionalForm[ $\partial_{r,c} q$ ], {r,  $\gamma$ s*}],
          {c,  $\gamma$ s}],  $\gamma$ s*}]]
    ], TableAlignments -> Center]
  ], Center] ];
```

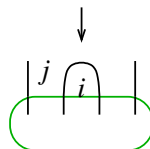
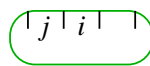
The Face-Centric Core.

```
 $\Sigma_{b1}[\sigma_1, PQ[C1_, q1_]] \oplus \Sigma_{b2}[\sigma_2, PQ[C2_, q2_]] \wedge :=$ 
  CF@ $\Sigma_{\text{Join}[b1, b2]}$ [ $\sigma_1 + \sigma_2, PQ[C1 \cup C2, q1 + q2]$ ];
```



GT for Gap Touch:

```
GTi,j@ $\Sigma_B[\{li___, i, ri___\}, \{lj___, j, rj___\}, bs___]$ [ $\sigma$ _,
  PQ[C_, q_]] :=
  CF@ $\Sigma_B[\{ri, li, j, rj, lj, i\}, bs]$ [ $\sigma$ _, PQ[C  $\cup$  { $\gamma_i - \gamma_j$ }, q]]
```



cor·don (kôr'dn)
n.



1. A line of people, military posts, or ships stationed around an area to enclose or guard it: a *police cordon*.
2. A rope, line, tape, or similar border stretched around an area, usually by the police, indicating that access is restricted.

use ϕ_p to kill its row and column, drop a $\begin{pmatrix} 01 \\ 10 \end{pmatrix}$ summand

use λ to kill θ , let $s += \text{sign}(\lambda)$

append θ to C_{rest} .

$$s \begin{pmatrix} 0 & \phi C_{\text{rest}} \\ \bar{\phi}^T & \lambda \theta \\ \bar{C}_{\text{rest}}^T & \bar{\theta}^T A_{\text{rest}} \end{pmatrix} \rightarrow \begin{cases} \exists p \phi_p \neq 0 & \text{use } \phi_p \text{ to kill its row and column, drop a } \begin{pmatrix} 01 \\ 10 \end{pmatrix} \text{ summand} \\ \phi = 0, \lambda \neq 0 & \text{use } \lambda \text{ to kill } \theta, \text{ let } s += \text{sign}(\lambda) \\ \phi = 0, \lambda = 0 & \text{append } \theta \text{ to } C_{\text{rest}}. \end{cases}$$

```
Cordoni@ $\Sigma_B[\{li___, i, ri___\}, bs___]$ [ $\sigma$ _, PQ[C_, q_]] :=
```

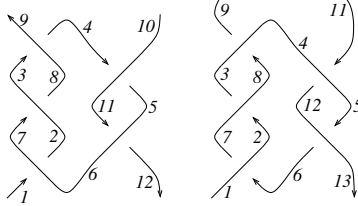
```
Module[{ $\phi = \partial_{\gamma_i} C$ ,  $\lambda = \partial_{\bar{\gamma}_i, \gamma_i} q$ , n $\sigma = \sigma$ , nC, nq, p},
  {p} = FirstPosition[({# != 0} & /@  $\phi$ , True, {0})];
  {nC, nq} = Which[
    p > 0, {C, q} /. ( $\gamma_i \rightarrow -C[[p]] / \phi[[p]]$ ) + ( $\gamma_i \rightarrow \theta$ )*,
     $\lambda \neq 0$ , (n $\sigma += \text{sign}[\lambda]$ ;
      {C, q} /. ( $\gamma_i \rightarrow -(\partial_{\bar{\gamma}_i} q) / \lambda$ ) + ( $\gamma_i \rightarrow \theta$ )*),
     $\lambda == 0$ , {C  $\cup$  { $\partial_{\bar{\gamma}_i} q$ }, q} /. ( $\gamma_i \rightarrow \theta$ )*];
  CF@ $\Sigma_B[\text{Most}@\{ri, li\}, bs]$ [n $\sigma$ ,
    PQ[nC, nq] /. ( $\gamma_{\text{Last}@\{ri, li\}} \rightarrow \gamma_{\text{First}@\{ri, li\}}$ )* ] ]
```


The Conway-Kinoshita-Terasaka Tangles.



$$T1 = PD[\bar{X}_{-6,2,7,-1}, \bar{X}_{-2,8,3,-7}, \bar{X}_{-8,4,9,-3}, X_{-11,6,12,-5}, X_{-4,11,5,-10}];$$

$$T2 = PD[X_{-6,2,7,-1}, X_{-2,8,3,-7}, X_{-8,4,9,-3}, \bar{X}_{-12,6,13,-5}, \bar{X}_{-4,12,5,-11}, \bar{X}_{-10,15,11,-14}, \bar{X}_{-15,10,16,-9}];$$



Column@{TL [T1], Kas [T1]}

$$-2\theta(u - \frac{\sqrt{3}}{2}) + 2\theta(u + \frac{\sqrt{3}}{2}) - 1$$

(Y_{-10})	Y_9	Y_{-1}	Y_{12}
\bar{Y}_{-10}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega^2-\omega+1}$
\bar{Y}_9	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega^2-\omega+1}$	$\frac{2\omega}{\omega^2-\omega+1}$
\bar{Y}_{12}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega^2-\omega+1}$	$\frac{2\omega}{\omega^2-\omega+1}$

(Y_{-10})	Y_9	Y_{-1}	Y_{12}
\bar{Y}_{-10}	$2(u-1)(u+1)(4u^2-3)$	0	$-2(u-1)(u+1)(4u^2-3)$
\bar{Y}_9	0	$\frac{1}{2(4u^2-3)}$	0
\bar{Y}_{-1}	$-2(u-1)(u+1)(4u^2-3)$	0	$2(u-1)(u+1)(4u^2-3)$
\bar{Y}_{12}	0	$-\frac{1}{2(4u^2-3)}$	0

Column@{TL [T2], Kas [T2]}

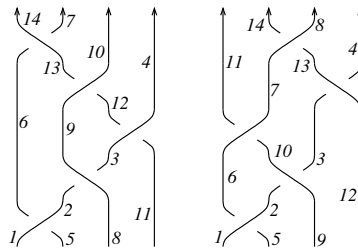
(Y_{-14})	Y_{16}	Y_{-1}	Y_{13}
\bar{Y}_{-14}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{16}	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{13}	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$

(Y_{-14})	Y_{16}	Y_{-1}	Y_{13}
\bar{Y}_{-14}	$\frac{1}{2}(-16u^4+28u^2-13)$	0	$\frac{1}{2}(16u^4-28u^2+13)$
\bar{Y}_{16}	0	$-\frac{2(u-1)(u+1)}{16u^4-28u^2+13}$	0
\bar{Y}_{-1}	$\frac{1}{2}(16u^4-28u^2+13)$	0	$\frac{1}{2}(-16u^4+28u^2-13)$
\bar{Y}_{13}	0	$\frac{2(u-1)(u+1)}{16u^4-28u^2+13}$	0

Examples with non-trivial co-dimension.

$$B1 = PD[X_{-5,2,6,-1}, \bar{X}_{-8,3,9,-2}, X_{-11,4,12,-3}, X_{-12,10,13,-9}, \bar{X}_{-13,7,14,-6}];$$

$$B2 = PD[X_{-5,2,6,-1}, \bar{X}_{-9,3,10,-2}, X_{-10,7,11,-6}, \bar{X}_{-12,4,13,-3}, X_{-13,8,14,-7}];$$



Column@{TL [B1], Kas [B1]}

(Y_{-11})	Y_4	Y_{10}	Y_7	Y_{14}	Y_{-1}	Y_{-5}	Y_{-8}
\bar{Y}_{-11}	0	0	0	0	0	0	0
\bar{Y}_4	0	0	0	$\frac{\omega-1}{\omega^2}$	0	$-\frac{\omega-1}{\omega^2}$	0
\bar{Y}_{10}	0	0	0	$-\frac{\omega-1}{\omega^2}$	0	$\frac{\omega-1}{\omega^2}$	0
\bar{Y}_7	0	0	0	$\frac{(\omega-1)^2}{\omega^2}$	0	$-\frac{(\omega-1)^2}{\omega^2}$	0
\bar{Y}_{14}	0	0	0	$-\frac{\omega-1}{\omega^2}$	0	$\frac{\omega-1}{\omega^2}$	0
\bar{Y}_{-1}	0	0	0	0	$1-\omega$	0	0
\bar{Y}_{-5}	0	$(\omega-1)\omega$	$1-\omega$	$-(\omega-1)^2$	$1-\omega$	$\frac{(\omega-1)^2}{\omega}$	0
\bar{Y}_{-8}	0	0	0	0	0	0	0

(Y_{-11})	Y_4	Y_{10}	Y_7	Y_{14}	Y_{-1}	Y_{-5}	Y_{-8}
\bar{Y}_{-11}	0	0	0	0	0	0	0
\bar{Y}_4	0	0	-1	$-u$	0	u	1
\bar{Y}_{10}	0	0	$-u$	$1-2u^2$	0	$2u^2-1$	0
\bar{Y}_7	0	-1	$-u$	$2u^2-3$	$-u$	0	1
\bar{Y}_{14}	0	$-u$	$1-2u^2$	$-u$	-1	$-2(u-1)(u+1)$	0
\bar{Y}_{-1}	0	0	-1	$-u$	0	u	1
\bar{Y}_{-5}	0	u	$2u^2-1$	0	$-2(u-1)(u+1)$	u	$4u^2-3$
\bar{Y}_{-8}	0	1	u	1	u	0	$1-2u^2$

Column@{TL [B2], Kas [B2]}

(Y_{-12})	Y_4	Y_8	Y_{14}	Y_{11}	Y_{-1}	Y_{-5}	Y_{-9}
\bar{Y}_{-12}	$\frac{\omega-1}{\omega}$	$\omega-1$	$-2(\omega-1)$	$\frac{2(\omega-1)^2}{\omega}$	0	$-\frac{2(\omega-1)^2}{\omega}$	$-\frac{(\omega-1)(2\omega-1)}{\omega}$
\bar{Y}_4	0	0	0	0	0	0	0
\bar{Y}_8	$\frac{2(\omega-1)}{\omega}$	$1-\omega$	0	$-\frac{(\omega-1)(2\omega-1)}{\omega}$	0	$\frac{2(\omega-1)}{\omega}$	$\frac{2(\omega-1)(\omega-1)}{\omega}$
\bar{Y}_{14}	0	0	0	0	0	0	0
\bar{Y}_{11}	$-2(\omega-1)\omega$	0	$2(\omega-1)\omega$	$-(\omega-1)(2\omega-1)$	0	$-\frac{2(\omega-1)^2}{\omega}$	$-\frac{2(\omega-1)(\omega-1)}{\omega}$
\bar{Y}_{-1}	0	0	0	0	0	$1-\omega$	0
\bar{Y}_{-5}	$2(\omega-1)\omega$	0	$-2(\omega-1)\omega$	$2(\omega-1)\omega$	0	$-\frac{2(\omega-1)^2}{\omega}$	$-\frac{2(\omega-1)(\omega-1)}{\omega}$
\bar{Y}_{-9}	$-\frac{(\omega-1)(2\omega-1)}{\omega}$	0	$\frac{2(\omega-1)(2\omega-1)}{\omega}$	$-\frac{2(\omega-1)(2\omega-1)}{\omega}$	0	$\frac{2(\omega-1)^2}{\omega}$	$-\frac{2(\omega-1)(\omega-1)}{\omega}$

(Y_{-12})	Y_4	Y_8	Y_{14}	Y_{11}	Y_{-1}	Y_{-5}	Y_{-9}
\bar{Y}_{-12}	0	0	0	0	0	0	0
\bar{Y}_4	0	$\frac{2(\omega-1)(2\omega-1)}{4\omega^2(4\omega^2-3)}$	$\frac{2(\omega-1)}{2\omega}$	0	0	$-\frac{2(\omega-1)(2\omega-1)}{4\omega^2(4\omega^2-3)}$	$-\frac{1}{4\omega^2(4\omega^2-3)}$
\bar{Y}_8	0	$-\frac{2(\omega-1)}{2\omega}$	$-2(u-1)(u+1)$	0	0	0	0
\bar{Y}_{14}	0	$\frac{2(\omega-1)}{4\omega^2(4\omega^2-3)}$	$\frac{2(\omega-1)(2\omega-1)}{2\omega}$	0	0	0	0
\bar{Y}_{11}	0	0	0	0	0	0	0
\bar{Y}_{-1}	0	$-\frac{2(\omega-1)(2\omega-1)}{4\omega^2(4\omega^2-3)}$	$-\frac{1}{2\omega}$	0	0	0	0
\bar{Y}_{-5}	0	$\frac{1}{2\omega(4\omega^2-3)}$	0	0	0	0	0
\bar{Y}_{-9}	0	$\frac{8\omega^4-6\omega^2-1}{4\omega^2(4\omega^2-3)}$	$\frac{1}{2\omega}$	0	0	$\frac{8\omega^4-6\omega^2-1}{4\omega^2(4\omega^2-3)}$	$\frac{32\omega^6-64\omega^4-38\omega^2-1}{4\omega^2(4\omega^2-3)}$

$\begin{pmatrix} A & B \\ C & U \end{pmatrix} \xrightarrow{\det(A)} \begin{pmatrix} I & A^{-1}B \\ C & U \end{pmatrix} \xrightarrow{1} \begin{pmatrix} I & A^{-1}B \\ 0 & U - CA^{-1}B \end{pmatrix}$. Roughly, $\det(A)$ is "det on ker", $-CA^{-1}B$ is "a pushforward of $\begin{pmatrix} A & B \\ C & U \end{pmatrix}$ ".
so $\det \begin{pmatrix} A & B \\ C & U \end{pmatrix} = \det(A) \det(U - CA^{-1}B)$. (what if $\mathbb{A}A^{-1}$?)

Questions. 1. Does this have a topological meaning? 2. Is there a version of the Kashaev Conjecture for tangles? 3. Find all solutions of R123 in our "algebra". 4. Braids and the Burau representation. 5. Recover the work in "Prior Art". 6. Are there any concordance properties? 7. What is the "SPQ group"? 8. The jumping points of signatures are the roots of the Alexander polynomial. Does this generalize to tangles? 9. Which of the three Cordon cases is the most common? 10. Are there interesting examples of tangles for which rels is non-trivial? 11. Is the pq part determined by Γ -calculus? 12. Is the pq part determined by finite type invariants? 13. Does it work with closed components / links? 14. Strand-doubling formulas? 15. A multivariable version? 16. Mutation invariance? 17. Ribbon knots? 18. Are there "face-virtual knots"? 19. Does the pushforward story extend to ranks? To formal Gaussian measures? To super Gaussian measures?

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Some Rigor.

(Exercises hints and partial solutions at end)

Exercise 1. Show that if two SPQ's S_1 and S_2 on V satisfy $\sigma(S_1 + U) = \sigma(S_2 + U)$ for every quadratic U on V , then they have the same shifts and the same domains.

Exercise 2. Show that if two full quadratics Q_1 and Q_2 satisfy $\sigma(Q_1 + U) = \sigma(Q_2 + U)$ for every U , then $Q_1 = Q_2$.

Proof of Theorem 1'. Fix W and consider triples $(V_1, S_1, \phi_1: V \rightarrow W)$ where $S_1 = (s_1, D_1, Q_1)$ is an SPQ on V_1 . Say that two such triples are "push-equivalent", $(V_1, S_1, \phi_1) \sim (V_2, S_2, \phi_2)$ if for every quadratic U on W ,

$$\sigma_{V_1}(S_1 + \phi_1^*U) = \sigma_{V_2}(S_2 + \phi_2^*U).$$

Given our (V, S, ϕ) , we need to show:

1. There is an SPQ S' on W such that $(V, S, \phi) \sim (W, S', I)$.
2. If $(W, S', I) \sim (W, S'', I)$ then $S' = S''$.

Property 2 is easy (Exercises 1, 2). Property 1 follows from the following three claims, each of which is easy.

Claim 1. If $v \in \ker \phi \cap D(S)$, and $\lambda := Q(v, v) \neq 0$, then $(V, S, \phi) \sim (V/\langle v \rangle, (s + \text{sign}(\lambda), D(S)/\langle v \rangle, Q - \lambda^{-1}Q(-, v) \otimes Q(v, -)), \phi/\langle v \rangle)$.

So wlog $Q|_{\ker \phi} = 0$ (meaning, $Q|_{\ker \phi \otimes \ker \phi} = 0$). \square

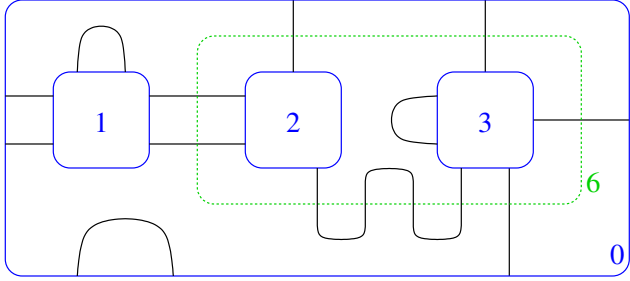
Claim 2. If $Q|_{\ker \phi} = 0$ and $v \in \ker \phi \cap D(S)$, let $V' = \ker Q(v, -)$ and then $(V, S, \phi) \sim (V', S|_{V'}, \phi|_{V'})$ so wlog $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$. \square

Claim 3. If $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$ then $S = \phi^*S'$ for some SPQ S' on $\text{im } \phi$ and then $(V, S, \phi) \sim (W, S', I)$. $\square \square$

Proof of Theorem 2. The functoriality of pullbacks needs no proof. Now assume $V_0 \xrightarrow{\alpha} V_1 \xrightarrow{\beta} V_2$ and that S is an SPQ on V_0 . Then for every SPQ U on V_2 we have, using reciprocity three times, that $\sigma(\beta_*\alpha_*S + U) = \sigma(\alpha_*S + \beta^*U) = \sigma(S + \alpha^*\beta^*U) = \sigma(S + (\beta\alpha)^*U) = \sigma((\beta\alpha)_*S + U)$. Hence $\beta_*\alpha_*S = (\beta\alpha)_*S$. \square

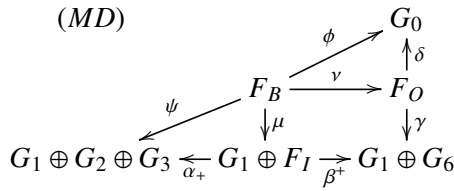
Proof of Theorem 3.

To prove Theorem 4, given three¹ SPQ's S_1 , S_2 , and S_3 , we need to show that planar-multiplying them in two steps, first using a planar connection diagram D_I (I for Inner) to yield $S_6 = \mathcal{S}(D_I)(S_2, S_3)$ and then using a second planar connection diagram D_O (O for Outer) to yield $\mathcal{S}(D_O)(S_1, S_6)$, gives the same answer as multiplying them all at once using the composition planar connection diagram $D_B = D_O \circ_6 D_I$ (B for Big) to yield $\mathcal{S}(D_B)(S_1, S_2, S_3)$.² An example should help:



In this example, if you ignore the dotted green line (marked “6”), you see the planar connection diagram D_B , which has three inputs (1,2,3) and a single output, the cycle 0. If you only look inside the green line, you see D_I , with inputs 2 and 3 and an output cycle 6. If you ignore the inside of 6 you see D_O , with inputs 1 and 6 and output cycle 0.

Let F_B (Big Faces) denote the vector space whose basis are the faces of D_B , let F_I (Inner Faces) be the space of faces of D_I , and let F_O



(Outer Faces) be the space of faces of D_O . Let G_1, G_2, G_3, G_6 , and G_0 be the spaces of gaps (edges) along the cycles 1,2,3,6, and 0, respectively. Let $\psi := \psi_{D_B}$ and $\phi := \phi^{D_B}$ be the maps defining $\mathcal{S}(D_B)$ and let $\gamma := \psi_{D_O}$ and $\delta := \phi^{D_O}$ be the maps defining $\mathcal{S}(D_O)$. Further, let $\alpha := \psi_{D_I}: F_I \rightarrow G_2 \oplus G_3$ and $\beta := \phi^{D_I}: F_I \rightarrow G_6$ be the maps defining $\mathcal{S}(D_I)$, and let $\alpha_+ := I \oplus \alpha$ and $\beta^+ := I \oplus \beta$ be the extensions of α and β by an identity on an extra factor of G_1 , so that $\beta^+ \alpha_+^* = I_{G_1} \oplus \mathcal{S}(D_I)$. Let μ map any big face to the sum of G_1 gaps around it, plus the sum of the inner faces it contains. Let ν map any big face to the sum of the outer faces it contains. It is easy to see that the master diagram (MD) shown on the right, made of all of these spaces and maps, is commutative.

Claim. The bottom right square of (MD) is an equalizer square, namely $F_B \simeq EQ(\beta^+, \gamma)$. Hence $\nu_* \mu^* = \gamma^* \beta^*$.

Proof. A big face (an element of F_B) is a sum of outer faces f_o and a sum of inner faces f_i , and it has a boundary g_1 on input cycle 1, such that the boundary of the outer pieces f_o is equal to the boundary of the inner pieces f_i plus g_1 . That matches perfectly with the definition of the equalizer: $EQ(\beta^+, \gamma) = \{(g_1, f_i, f_o) : \beta^+(g_1, f_i) = \gamma(f_o)\} = \{(g_1, f_i, f_o) : \gamma(f_o) = (g_1, \beta(f_i))\}$. \square

Proof of Theorem 4. With notation as above, with the example above (which is general enough), and with the claim above, and

also using functoriality, we have $\mathcal{S}(D_B) = \phi_* \psi^* = \delta_* \nu_* \mu^* \alpha_+^* = \delta_* \gamma^* \beta^+ \alpha_+^* = \mathcal{S}(D_O) \circ (I_{G_1} \oplus \mathcal{S}(D_I))$, as required. \square

Proof of Theorem 5. We need to verify the Reidemeister moves and that was done in the computational section, and the statement about the restriction to links, which is easy: simply assemble an n -crossing knot using an n -input planar connection diagram, and the formulas clearly match. \square

Further Homework.

Exercise 3. Show that if two SPQ's S_1 and S_2 on $V \oplus A$ satisfy $A \subset \text{rad } S_i$ and $\sigma(S_1 + \pi^* U) = \sigma(S_2 + \pi^* U)$ for every quadratic U on V , where $\pi: V \oplus A \rightarrow V$ is the projection, then $S_1 = S_2$.

Exercise 4. Show that if $\phi: V \rightarrow W$ is surjective and Q is a quadratic on W , then $\sigma(Q) = \sigma(\phi^* Q)$.

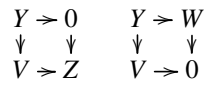
Exercise 5. Show that always, $\phi_* \phi^* S = S|_{\text{im } \phi}$.

Exercise 6. By taking $U = 0$ in the reciprocity statement, prove that always $\sigma(\phi_* S) = \sigma(S)$. But that seems wrong, if $\phi = 0$. What saves the day?

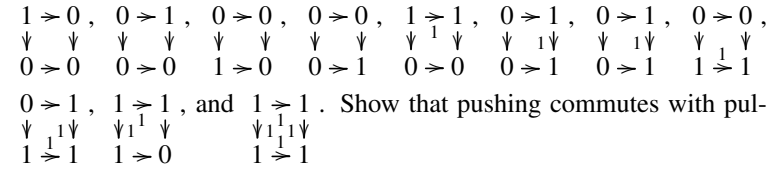
Exercise 7. By taking $S = 0$ in the reciprocity statement, prove that always $\sigma(\phi^* U) = \sigma(U)$. But wait, this is nonsense! What went wrong?

Exercise 8. Given $\phi: V \rightarrow W$ and a subspace $D \subset V$, show that there is a unique subspace $\phi_* D \subset W$ such that for every quadratic Q on W , $\sigma(\phi^* Q|_D) = \sigma(Q|_{\phi_* D})$.

Exercise 9. When are diagrams as on the right equalizer diagrams? What then do we learn from Theorem 3?



Exercise 10. There are 11 types of irreducible commutative squares:



ling for all but four of them. Compare with the statement of Theorem 3.

Exercise 11. Prove that a square is admissible iff it is an equalizer square, with an additional direct summand A added to the Y term, and with the maps μ and ν extended by 0 on A .

Exercise 12. Prove that the direct sum of two admissible squares is admissible. *Warning:* Harder than it seems! Not all quadratics on $V_1 \oplus V_2$ are direct sums of quadratics on V_1 and on V_2 .

Exercise 13. Given a quadratic Q on a space V , let π be the projection $V \rightarrow V/\text{rad}(Q)$ and show that $\pi_* Q = Q/\text{rad}(Q)$, with the obvious definition for the latter.

Exercise 14. Show that for any partial quadratic Q on a space W there exists a space A and a fully-defined quadratic F on $W \oplus A$ such that $\pi_* F = Q$, where $\pi: W \oplus A \rightarrow W$ is the projection (these are not unique). Furthermore, if $\phi: V \rightarrow W$, then $\phi^* Q = \pi_* \phi_+^* F$, where $\phi_+ = \phi \oplus I: V \oplus A \rightarrow W \oplus A$ and π also denotes the projection $V \oplus A \rightarrow V$.

Solutions / Hints.

Hint for 1. On a vector in the domain of one but not the other, take an outrageous value for U , that will raise or lower the signature.

Hint for 2. WLOG, Q_1 is diagonal and $Q_1 = 0$.

Hint for 5. It's enough to test that against U with $\mathcal{D}(U) = \text{im } \phi$.

Hint for 6. The “shift” part of $0_* S$ is $\sigma(S)$.

Hint for 7. $\phi_* S$ isn't 0, it's the *partial* quadratic “0 on $\text{im } \phi$ ” (and indeed, $\sigma(\phi^* U) = \sigma(U)$ if ϕ is surjective).

Hint for 10. The exceptions are $\begin{smallmatrix} 01 \\ 00 \end{smallmatrix}$, $\begin{smallmatrix} 00 \\ 10 \end{smallmatrix}$, $\begin{smallmatrix} 01 \\ 11 \end{smallmatrix}$, and $\begin{smallmatrix} 11 \\ 10 \end{smallmatrix}$.

Hint for 12. Use Exercise 11.

¹Truly, we need the same for any number of input SPQ's that are divided into two groups, “multiply in the first step” and “multiply in the second step”. But there's no added difficulty here, only an added notational complexity.

²Aren't we sassy? We picked “6” for the name of the product of “2” and “3”.