

# Shifted Partial Quadratics, their Pushforwards, and Signature Invariants for Tangles

**Abstract.** Following a general discussion of the computation of zombians of unfinished columbaria (with examples), I will tell you about my recent joint work w/ Jessica Liu on what we feel is the “textbook” extension of knot signatures to tangles, which for unknown reasons, is not in any of the textbooks that we know.



Jessica Liu



Columbaria in an East Sydney Cemetery



Jacobian, Hamiltonian, Zombian

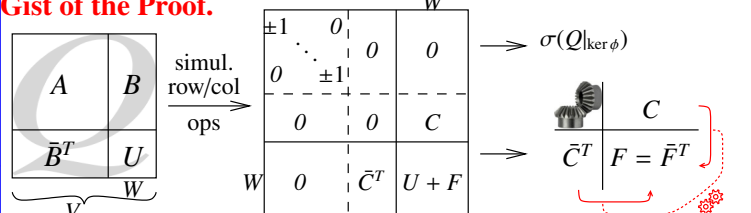
**Kashaev’s Conjecture** [Ka] For links,  $\sigma_{Kas} = 2\sigma_{TL}$ .  
**Liu’s Theorem** [Li].

A **Partial Quadratic (PQ)** on  $V$  is a quadratic  $Q$  defined only on a subspace  $\mathcal{D}_Q \subset V$ . We add PQs with  $\mathcal{D}_{Q_1+Q_2} := \mathcal{D}_{Q_1} \cap \mathcal{D}_{Q_2}$ . Given a linear  $\psi: V \rightarrow W$  and a PQ  $Q$  on  $W$ , there is an obvious pullback  $\psi^*Q$ , a PQ on  $V$ .

**Theorem 1.** Given a linear  $\phi: V \rightarrow W$  and a PQ  $Q$  on  $V$ , there is a unique pushforward PQ  $\phi_*Q$  on  $W$  such that for every PQ  $U$  on  $W$ ,  $\sigma_V(Q + \phi^*U) = \sigma_{\ker\phi}(Q|_{\ker\phi}) + \sigma_W(U + \phi_*Q)$ . (If you must,  $\mathcal{D}(\phi_*Q) = \phi(\text{ann}_Q(\mathcal{D}(Q) \cap \ker\phi))$  and  $(\phi_*Q)(w) = Q(v)$ , where  $v$  is s.t.  $\phi(v) = w$  and  $Q(v, \text{rad } Q|_{\ker\phi}) = 0$ ).

**Prior Art** on signatures for tangles / braids. Gambaudo and Ghys [GG], Cimasoni and Conway [CC], Conway [Co], Merz [Me]. All define signatures of tangles / braids by first closing them to links and then work hard to derive composition properties.

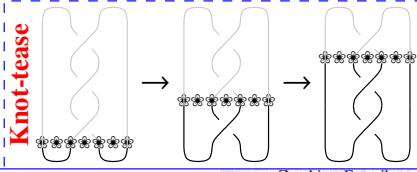
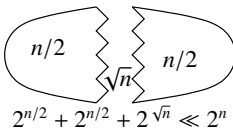
### Gist of the Proof.



... and the quadratic  $F := \phi_*Q$  is well-defined only on  $D := \ker C$ . Exactly what we want, if the Zombian is the signature!

### Why Tangles? • Faster!

- Conceptually clearer proofs of invariance (and of skein relations).
- Often fun and consequential:
  - The Jones Polynomial  $\rightsquigarrow$  The Temperley-Lieb Algebra.
  - Khovanov Homology  $\rightsquigarrow$  “Unfinished complexes”, complexes in a category.
  - The Kontsevich Integral  $\rightsquigarrow$  Associators.
  - HFK  $\rightsquigarrow$  OMG, type  $D$ , type  $A$ ,  $\mathcal{A}_\infty, \dots$



### Computing Zombians of Unfinished Columbaria.

- Must be no slower than for finished ones.
- Future zombies must be able to complete the computation.
- Future zombies must not even know the size of the task that today’s zombies were facing.
- We must be able to extend to ZPUCs, Zombie Processed Unfinished Columbaria!

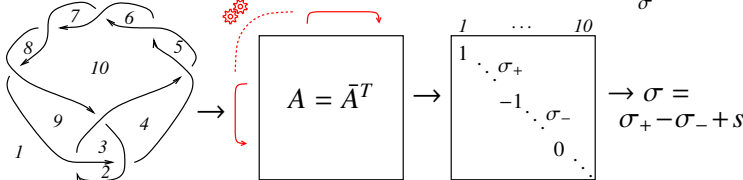


Columbarium near Assen

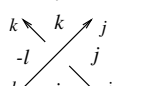
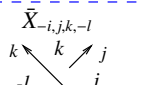
**Example / Exercise.** Compute the determinant of a  $1,000 \times 1,000$  matrix in which 50 entries are not yet given.

**Homework / Research Projects.** • What with ZPUCs? • Use this to get an Alexander tangle invariant.

**Reminders.** {links}  $\rightleftharpoons$  {matrices / quadratic forms}  $\xrightarrow{\text{signature}} \mathbb{Z}$ :

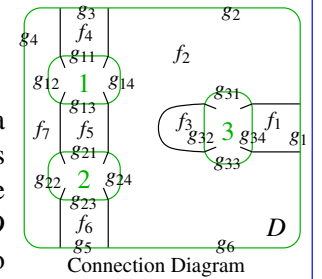


With  $|\omega| = 1$ ,  $t = 1 - \omega$ ,  $r = t + \bar{t}$ ,  $v = \text{Re}(\omega)$ , and  $u = \text{Re}(\omega^{1/2})$ :

$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} -r & -t & 2t & \bar{t} \\ -\bar{t} & t & -r & -\bar{t} \\ 2\bar{t} & t & -r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A += \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$s = 0$		$s = -1$
$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} r & -t & -2\bar{t} & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ -2t & t & r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A -= \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$s = 0$		$s = +1$

**Definition.**  $S \left( \begin{matrix} g_2 \\ g_3 \dots g_1 \end{matrix} \right) := \left\{ \begin{matrix} \text{SPQ } S \\ \text{on } \langle g_i \rangle \end{matrix} \right\}$ .

**Theorem 4.**  $\{S(\text{cyclic sets})\}$  is a planar algebra, with compositions  $S(D)((S_i)) := \phi_*^D(\psi_D^*(\bigoplus_i S_i))$ , where  $\psi_D: \langle f_i \rangle \rightarrow \langle g_{\alpha i} \rangle$  maps every face of  $D$  to the sum of the input gaps adjacent to it and  $\phi^D: \langle f_i \rangle \rightarrow \langle g_i \rangle$  maps every face to the sum of the output gaps adjacent to it. So for our  $D$ ,  $\psi_D: f_1 \mapsto g_{34}, f_2 \mapsto g_{31} + g_{14} + g_{24} + g_{33}, f_3 \mapsto g_{32}, f_4 \mapsto g_{11}, f_5 \mapsto g_{13} + g_{21}, f_6 \mapsto g_{23}, f_7 \mapsto g_{12} + g_{22}$  and  $\phi^D: f_1 \mapsto g_1, f_2 \mapsto g_2 + g_6, f_3 \mapsto 0, f_4 \mapsto g_3, f_5 \mapsto 0, f_6 \mapsto g_5, f_7 \mapsto g_4$ .



**Theorem 5.** TL and Kas, defined on X and  $\bar{X}$  as before, extend to planar algebra morphisms {tangles}  $\rightarrow \{S\}$ . Restricted to links,  $TL = \sigma_{TL}$  and  $Kas = \sigma_{Kas}$ .



Levine, Tristram, Kashaev

**Implementation** (sources: <http://drorbn.net/icerm23/ap>). I like it most when the implementation matches the math perfectly. We failed here.

Once[<< KnotTheory`];

Loading KnotTheory` version

of February 2, 2020, 10:53:45.2097.

Read more at <http://katlas.org/wiki/KnotTheory>.

**Utilities.** The step function, algebraic numbers, canonical forms.

$\theta[x\_]$  /; NumericQ[x] := UnitStep[x]

```
 $\omega 2[v\_][p\_]$  := Module[{q = Expand[p], n, c},
  If[q == 0, 0,
    c = Coefficient[q,  $\omega$ , n = Exponent[q,  $\omega$ ]];
     $c v^n + \omega 2[v][q - c (\omega + \omega^{-1})^n]$ ];
```

```
sign[ $\mathcal{E}$ _] := Module[{n, d, v, p, rs, e, k},
  {n, d} = NumeratorDenominator[ $\mathcal{E}$ ];
  {n, d} /=  $\omega^{\text{Exponent}[n, \omega]/2 + \text{Exponent}[n, \omega, \text{Min}]/2}$ ;
  p = Factor[ $\omega 2[v] @ n * \omega 2[v] @ d /. v \rightarrow 4 u^2 - 2$ ];
  rs = Solve[p == 0, u, Reals];
  If[rs == {}, Sign[p /. u -> 0],
    rs = Union@{u /. rs};
    Sign[(-1)e=Exponent[p, u] Coefficient[p, u, e]] + Sum[
      k = 0;
      While[{d = RootReduce[ $\partial_{\{u, ++k\}} p /. u \rightarrow r$ ]} == 0];
      If[EvenQ[k], 0, 2 Sign[d]] *  $\theta[u - r]$ ,
      {r, rs}]]
]
```

SetAttributes[B, Orderless];

CF[b\_B] := RotateLeft[#, First@Ordering[#] - 1] & /@ DeleteCases[b, {}]

```
CF[ $\mathcal{E}$ _] := Module[{ $\gamma$ s = Union@Cases[ $\mathcal{E}$ ,  $\gamma$ _ |  $\bar{\gamma}$ _,  $\infty$ ]},
  Total[CoefficientRules[ $\mathcal{E}$ ,  $\gamma$ s] /.
    (ps_ -> c_) := Factor[c]  $\times$  Times@@ $\gamma$ sps]]
```

CF[{}] = {};

CF[C\_List] :=

```
Module[{ $\gamma$ s = Union@Cases[C,  $\gamma$ _,  $\infty$ ],  $\gamma$ },
  CF /@ DeleteCases[0] [
    RowReduce[Table[ $\partial_{\gamma} r$ , {r, C}, { $\gamma$ ,  $\gamma$ s}]] .  $\gamma$ s ]
```

( $\mathcal{E}$ \_)\* :=  $\mathcal{E} /. \{\bar{\gamma} \rightarrow \gamma, \gamma \rightarrow \bar{\gamma}, \omega \rightarrow \omega^{-1}, c\_Complex \rightarrow c^*\}$ ;

r\_Rule\* := {r, r\*}

RulesOf[ $\gamma_i$  + rest\_] := ( $\gamma_i \rightarrow -rest$ )\*;

```
CF[PQ[C_, q_]] := Module[{nC = CF[C]},
  PQ[nC, CF[q /. Union@@RulesOf /@ nC]] ]
```

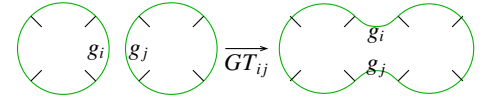
CF[ $\Sigma_b$ [ $\sigma$ \_, pq\_]] :=  $\Sigma_{CF[b]}$ [ $\sigma$ , CF[pq]]

## Pretty-Printing.

```
Format[ $\Sigma_{b,B}[\sigma_, PQ[C_, q_]]$ ] := Module[{ $\gamma$ s},
   $\gamma$ s =  $\gamma$ # & /@ Join@@b;
  Column[{TraditionalForm@ $\sigma$ ,
    TableForm[Join[
      Prepend[""] /@ Table[TraditionalForm[ $\partial_c r$ ],
        {r, C}, {c,  $\gamma$ s}],
      {Prepend[""] [
        Join@@
          (b /. {L_, m___, r_} =>
            {DisplayForm@RowBox[{"(", L}],
              m, DisplayForm@RowBox[{r, ")"}]}) /.
            i_Integer =>  $\gamma_i$  ]}],
      MapThread[Prepend,
        {Table[TraditionalForm[ $\partial_{r,c} q$ ], {r,  $\gamma$ s*}],
          {c,  $\gamma$ s}],  $\gamma$ s*}]]],
    TableAlignments -> Center]
  ], Center] ];
```

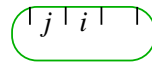
## The Face-Centric Core.

$\Sigma_{b1}[\sigma_1, PQ[C1_, q1_]] \oplus \Sigma_{b2}[\sigma_2, PQ[C2_, q2_]] \wedge :=$   
 $CF @ \Sigma_{\text{Join}[b1, b2]}[\sigma_1 + \sigma_2, PQ[C1 \cup C2, q1 + q2]];$



GT for Gap Touch:

```
GTi,j@ $\Sigma_B[\{li___, i, ri___\}, \{lj___, j, rj___\}, bs___]$ [ $\sigma$ _,
  PQ[C_, q_]] :=
  CF@ $\Sigma_B[\{ri, li, j, rj, lj, i\}, bs]$ [ $\sigma$ _, PQ[C  $\cup$  { $\gamma_i - \gamma_j$ }, q]]
```



cordon (kôr'dn)



n.

1. A line of people, military posts, or ships stationed around an area to enclose or guard it: *a police cordon*.
2. A rope, line, tape, or similar border stretched around an area, usually by the police, indicating that access is restricted.

$$s \begin{pmatrix} 0 & \phi C_{\text{rest}} \\ \bar{\phi}^T & \lambda \theta \\ \bar{C}_{\text{rest}}^T & \bar{\theta}^T A_{\text{rest}} \end{pmatrix} \rightarrow \begin{cases} \exists p \phi_p \neq 0 & \text{use } \phi_p \text{ to kill its row and} \\ \phi = 0, \lambda \neq 0 & \text{column, drop a } \begin{pmatrix} 01 \\ 10 \end{pmatrix} \text{ summand} \\ \phi = 0, \lambda = 0 & \text{use } \lambda \text{ to kill } \theta, \text{ let } s += \text{sign}(\lambda) \\ & \text{append } \theta \text{ to } C_{\text{rest}}. \end{cases}$$

Cordon<sub>i</sub>@ $\Sigma_B[\{li___, i, ri___\}, bs___]$ [ $\sigma$ \_, PQ[C\_, q\_]] :=

```
Module[{ $\phi = \partial_{\gamma_i} C$ ,  $\lambda = \partial_{\bar{\gamma}_i, \gamma_i} q$ , n $\sigma = \sigma$ , nC, nq, p},
  {p} = FirstPosition[ (# != 0) & /@  $\phi$ , True, {0}];
  {nC, nq} = Which[
    p > 0, {C, q} /. ( $\gamma_i \rightarrow -C[[p]] / \phi[[p]]$ )* /. ( $\gamma_i \rightarrow \theta$ )*,
     $\lambda \neq 0$ , (n $\sigma += \text{sign}[\lambda]$ ;
      {C, q} /. ( $\gamma_i \rightarrow -(\partial_{\bar{\gamma}_i} q) / \lambda$ )* /. ( $\gamma_i \rightarrow \theta$ )*),
     $\lambda == 0$ , {C  $\cup$  { $\partial_{\bar{\gamma}_i} q$ }, q} /. ( $\gamma_i \rightarrow \theta$ )*];
  CF@ $\Sigma_B[\text{Most}@\{ri, li\}, bs]$ [n $\sigma$ ,
    PQ[nC, nq] /. ( $\gamma_{\text{Last}@\{ri, li\}} \rightarrow \gamma_{\text{First}@\{ri, li\}}$ )* ] ]
```

**Strand Operations.** c for contract, mc for magnetic contract:

$$C_{i,j}@t : \Sigma_B[\{li\_ , i, ri\_ \}, \{ \_ , j, \_ \}, \_ ] [ \_ ] := t // GT_{j, First\{ri, li\}} // Cordon_j$$

$$C_{i,j}@t : \Sigma_B[\{ \_ , i, j, \_ \}, \_ ] [ \_ ] := Cordon_j @ t$$

$$C_{i,j}@t : \Sigma_B[\{j, \_ , i, \_ \}, \_ ] [ \_ ] := Cordon_j @ t$$

$$C_{i,j}@t : \Sigma_B[\{ \_ , j, i, \_ \}, \_ ] [ \_ ] := Cordon_i @ t$$

$$C_{i,j}@t : \Sigma_B[\{i, \_ , j, \_ \}, \_ ] [ \_ ] := Cordon_i @ t$$

$$mc[\mathcal{E}_] := \mathcal{E} //$$

$$t : \Sigma_B[\{ \_ , i, \_ \}, \{ \_ , j, \_ \}, \_ ] [ \_ ] | \Sigma_B[\{ \_ , i, j, \_ \}, \_ ] [ \_ ] | \Sigma_B[\{j, \_ , i, \_ \}, \_ ] [ \_ ] / ; i + j == 0 \Rightarrow C_{i,j}@t$$

**The Crossings** (and empty strands).

$$Kas@P_{i,j} := CF@ \Sigma_B[\{i,j\}] [\theta, PQ[\{ \}, \theta]] ;$$

$$TL@P_{i,j} := CF@ \Sigma_B[\{i,j\}] [\theta, PQ[\{ \}, \theta]]$$

$$Kas[x : X[i, j, k, l]] :=$$

$$Kas@If[PositiveQ[x], X_{-i,j,k,-l}, \bar{X}_{-j,k,l,-i}] ;$$

$$Kas[(x : X | \bar{X})_{fs\_}] := Module[\{v = 2u^2 - 1, p, \gamma s, m\},$$

$$\gamma s = \gamma_{\#} \& /@ \{fs\}; p = (x == X) ;$$

$$m = If[p, \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}, -\begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}] ;$$

$$CF@ \Sigma_B[\{fs\}] [If[p, -1, 1], PQ[\{ \}, \gamma s^* . m . \gamma s]]$$

$$TL[x : X[i, j, k, l]] :=$$

$$TL@If[PositiveQ[x], X_{-i,j,k,-l}, \bar{X}_{-j,k,l,-i}] ;$$

$$TL[(x : X | \bar{X})_{fs\_}] := Module[\{t = 1 - \omega, r, \gamma s, m\},$$

$$r = t + t^*; \gamma s = \gamma_{\#} \& /@ \{fs\};$$

$$m = If[x == X,$$

$$\begin{pmatrix} -r & -t & 2t & t^* \\ -t^* & \theta & t^* & \theta \\ 2t^* & t & -r & -t^* \\ t & \theta & -t & \theta \end{pmatrix}, \begin{pmatrix} r & -t & -2t^* & t^* \\ -t^* & \theta & t^* & \theta \\ -2t & t & r & -t^* \\ t & \theta & -t & \theta \end{pmatrix}] ;$$

$$CF@ \Sigma_B[\{fs\}] [\theta, PQ[\{ \}, \gamma s^* . m . \gamma s]]$$

**Evaluation on Tangles and Knots.**

$$Kas[K_] := Fold[mc[\#1 \oplus \#2] \&, \Sigma_B[] [\theta, PQ[\{ \}, \theta]], List@@ (Kas /@ PD@K)] ;$$

$$KasSig[K_] := Expand[Kas[K][[1]] / 2]$$

$$TL[K_] :=$$

$$Fold[mc[\#1 \oplus \#2] \&, \Sigma_B[] [\theta, PQ[\{ \}, \theta]], List@@ (TL /@ PD@K)] / .$$

$$\theta[c\_ + u] / ; Abs[c] \ge 1 \Rightarrow \theta[c] ;$$

$$TLSig[K_] := TL[K][[1]]$$

**Reidemeister 3.**

$$R3L = PD[X_{-2,5,4,-1}, X_{-3,7,6,-5}]$$

$$X_{-6,9,8,-4} ;$$

$$R3R = PD[X_{-3,5,4,-2}, X_{-4,6,8,-1}]$$

$$X_{-5,7,9,-6} ;$$

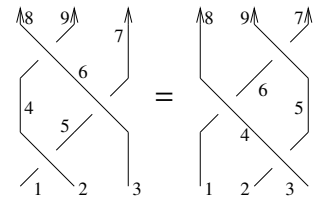
$$\{TL@R3L == TL@R3R, Kas@R3L == Kas@R3R\}$$

$$\{True, True\}$$

**Kas@R3L**

$$2\theta(u - \frac{1}{2}) - 2\theta(u + \frac{1}{2}) - 2$$

	$\gamma_3$	$\gamma_7$	$\gamma_9$	$\gamma_8$	$\gamma_{-1}$	$\gamma_{-2}$
$\bar{\gamma}_3$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$
$\bar{\gamma}_7$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2(2u^2-1)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$
$\bar{\gamma}_9$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$
$\bar{\gamma}_8$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$
$\bar{\gamma}_{-1}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2(2u^2-1)}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$
$\bar{\gamma}_{-2}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$-\frac{2u}{(2u-1)(2u+1)}$	$-\frac{1}{(2u-1)(2u+1)}$	$\frac{u(4u^2-3)}{(2u-1)(2u+1)}$	$\frac{2u^2(4u^2-3)}{(2u-1)(2u+1)}$



**Reidemeister 2.**

$$TL@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}]$$

$$\begin{matrix} & \theta & & & \\ & 1 & \theta & -1 & \theta \\ (\gamma_{-2} & & \gamma_6 & \gamma_5 & \gamma_{-1}) \\ \bar{\gamma}_{-2} & \theta & \theta & \theta & \theta \\ \bar{\gamma}_6 & \theta & \theta & \theta & \theta \\ \bar{\gamma}_5 & \theta & \theta & \theta & \theta \\ \bar{\gamma}_{-1} & \theta & \theta & \theta & \theta \end{matrix}$$

$$\{TL@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}] == GT_{5,-2}@TL@PD[P_{-1,5}, P_{-2,6}], Kas@PD[X_{-2,4,3,-1}, \bar{X}_{-4,6,5,-3}] == GT_{5,-2}@Kas@PD[P_{-1,5}, P_{-2,6}]\}$$

$$\{True, True\}$$

**Reidemeister 1.**

$$\{TL@PD[X_{-3,3,2,-1}] == TL@P_{-1,2},$$

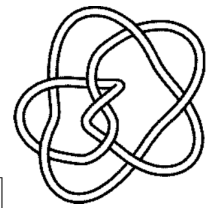
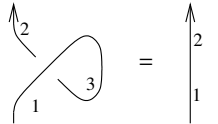
$$Kas@PD[X_{-3,3,2,-1}] == Kas@P_{-1,2}\}$$

$$\{True, True\}$$

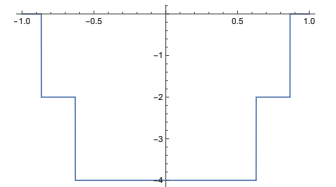
**A Knot.**

$$f = TLSig[Knot[8, 5]]$$

$$2\theta\left[-\frac{\sqrt{3}}{2} + u\right] - 2\theta\left[\frac{\sqrt{3}}{2} + u\right] - 2\theta\left[u - \left(\text{clockwise} - 0.630\dots\right)\right] + 2\theta\left[u - \left(\text{counterclockwise} 0.630\dots\right)\right]$$



$$Plot[f, \{u, -1, 1\}]$$

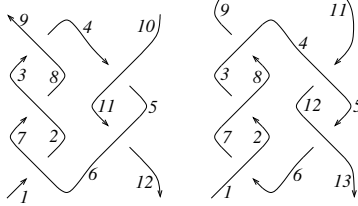


# The Conway-Kinoshita-Terasaka Tangles.



$$T1 = PD[\bar{X}_{-6,2,7,-1}, \bar{X}_{-2,8,3,-7}, \bar{X}_{-8,4,9,-3}, X_{-11,6,12,-5}, X_{-4,11,5,-10}];$$

$$T2 = PD[X_{-6,2,7,-1}, X_{-2,8,3,-7}, X_{-8,4,9,-3}, \bar{X}_{-12,6,13,-5}, \bar{X}_{-4,12,5,-11}, \bar{X}_{-10,15,11,-14}, \bar{X}_{-15,10,16,-9}];$$



## Column@{TL [T1], Kas [T1]}

$$-2\theta(u - \frac{\sqrt{3}}{2}) + 2\theta(u + \frac{\sqrt{3}}{2}) - 1$$

$\bar{Y}_{-10}$	$Y_9$	$Y_{-1}$	$Y_{12}$
$\frac{\omega-1}{\omega}$	$1-\omega$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$

$\bar{Y}_{-10}$	$Y_9$	$Y_{-1}$	$Y_{12}$
$2(u-1)(u+1)(4u^2-3)$	$\frac{1}{2(4u^2-3)}$	$-2(u-1)(u+1)(4u^2-3)$	$\frac{1}{2(4u^2-3)}$
$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$
$-2(u-1)(u+1)(4u^2-3)$	$\frac{1}{2(4u^2-3)}$	$2(u-1)(u+1)(4u^2-3)$	$\frac{1}{2(4u^2-3)}$
$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$	$\frac{1}{2(4u^2-3)}$

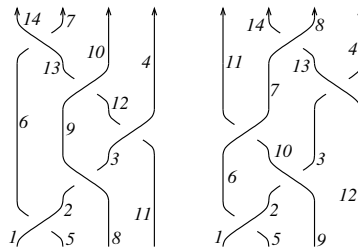
## Column@{TL [T2], Kas [T2]}

$\bar{Y}_{-14}$	$Y_{16}$	$Y_{-1}$	$Y_{13}$
$\frac{\omega-1}{\omega}$	$1-\omega$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
$\frac{\omega-1}{\omega}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$
$\frac{\omega-1}{\omega}$	$\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2(\omega-1)^2\omega}{\omega^4-3\omega^3+5\omega^2-3\omega+1}$

## Examples with non-trivial co-dimension.

$$B1 = PD[X_{-5,2,6,-1}, \bar{X}_{-8,3,9,-2}, X_{-11,4,12,-3}, X_{-12,10,13,-9}, \bar{X}_{-13,7,14,-6}];$$

$$B2 = PD[X_{-5,2,6,-1}, \bar{X}_{-9,3,10,-2}, X_{-10,7,11,-6}, \bar{X}_{-12,4,13,-3}, X_{-13,8,14,-7}];$$



## Column@{TL [B1], Kas [B1]}

$\bar{Y}_{-11}$	$Y_4$	$Y_{10}$	$Y_7$	$Y_{14}$	$Y_{-1}$	$Y_{-5}$	$Y_{-8}$
$\frac{\omega-1}{\omega}$	$1-\omega$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$

## Column@{TL [B2], Kas [B2]}

$\bar{Y}_{-12}$	$Y_4$	$Y_8$	$Y_{14}$	$Y_{11}$	$Y_{-1}$	$Y_{-5}$	$Y_{-9}$
$\frac{\omega-1}{\omega}$	$1-\omega$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$
$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$

$\bar{Y}_{-12}$	$Y_4$	$Y_8$	$Y_{14}$	$Y_{11}$	$Y_{-1}$	$Y_{-5}$	$Y_{-9}$
$\frac{1}{\omega}$	$1-\omega$	$\frac{1}{\omega}$	$\frac{1}{\omega}$	$\frac{1}{\omega}$	$\frac{1}{\omega}$	$\frac{1}{\omega}$	$\frac{1}{\omega}$
$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$-\frac{2\omega}{\omega^2-\omega+1}$
$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$	$\frac{1}{\omega}$	$\frac{2\omega}{\omega^2-\omega+1}$

$\begin{pmatrix} A & B \\ C & U \end{pmatrix} \xrightarrow{\det(A)} \begin{pmatrix} I & A^{-1}B \\ C & U \end{pmatrix} \xrightarrow{1} \begin{pmatrix} I & A^{-1}B \\ 0 & U - CA^{-1}B \end{pmatrix}$ . Roughly,  $\det(A)$  is "det on ker",  $-CA^{-1}B$  is "a pushforward of  $\begin{pmatrix} A & B \\ C & U \end{pmatrix}$ ".  
so  $\det \begin{pmatrix} A & B \\ C & U \end{pmatrix} = \det(A) \det(U - CA^{-1}B)$ . (what if  $\mathbb{A}A^{-1}$ ?)

**Questions.** 1. Does this have a topological meaning? 2. Is there a version of the Kashaev Conjecture for tangles? 3. Find all solutions of R123 in our "algebra". 4. Braids and the Burau representation. 5. Recover the work in "Prior Art". 6. Are there any concordance properties? 7. What is the "SPQ group"? 8. The jumping points of signatures are the roots of the Alexander polynomial. Does this generalize to tangles? 9. Which of the three Cordon cases is the most common? 10. Are there interesting examples of tangles for which rels is non-trivial? 11. Is the  $pq$  part determined by  $\Gamma$ -calculus? 12. Is the  $pq$  part determined by finite type invariants? 13. Does it work with closed components / links? 14. Strand-doubling formulas? 15. A multivariable version? 16. Mutation invariance? 17. Ribbon knots? 18. Are there "face-virtual knots"? 19. Does the pushforward story extend to ranks? To formal Gaussian measures? To super Gaussian measures?

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**Acknowledgement.** This work was partially supported by NSERC grant RGPIN-2018-04350 and by the Chu Family Foundation (NYC).

**Some Rigor.**

(Exercises hints and partial solutions at end)

**Exercise 1.** Show that if two SPQ's  $S_1$  and  $S_2$  on  $V$  satisfy  $\sigma(S_1 + U) = \sigma(S_2 + U)$  for every quadratic  $U$  on  $V$ , then they have the same shifts and the same domains.

**Exercise 2.** Show that if two full quadratics  $Q_1$  and  $Q_2$  satisfy  $\sigma(Q_1 + U) = \sigma(Q_2 + U)$  for every  $U$ , then  $Q_1 = Q_2$ .

**Proof of Theorem 1'.** Fix  $W$  and consider triples  $(V, S, \phi: V \rightarrow W)$  where  $S = (s, D, Q)$  is an SPQ on  $V$ . Say that two triples are "push-equivalent",  $(V_1, S_1, \phi_1) \sim (V_2, S_2, \phi_2)$  if for every quadratic  $U$  on  $W$ ,

$$\sigma_{V_1}(S_1 + \phi_1^* U) = \sigma_{V_2}(S_2 + \phi_2^* U).$$

Given our  $(V, S, \phi)$ , we need to show:

1. There is an SPQ  $S'$  on  $W$  such that  $(V, S, \phi) \sim (W, S', I)$ .
2. If  $(W, S', I) \sim (W, S'', I)$  then  $S' = S''$ .

Property 2 is easy (Exercises 1, 2). Property 1 follows from the following three claims, each of which is easy.

**Claim 1.** If  $v \in \ker \phi \cap D(S)$ , and  $\lambda := Q(v, v) \neq 0$ , then  $(V, S, \phi) \sim$

$$(V/\langle v \rangle, (s + \text{sign}(\lambda), D(S)/\langle v \rangle, Q - \lambda^{-1} Q(-, v) \otimes Q(v, -)), \phi/\langle v \rangle).$$

So wlog  $Q|_{\ker \phi} = 0$  (meaning,  $Q|_{\ker \phi \otimes \ker \phi} = 0$ ).

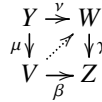
**Claim 2.** If  $Q|_{\ker \phi} = 0$  and  $v \in \ker \phi \cap D(S)$ , let  $V' = \ker Q(v, -)$  and then  $(V, S, \phi) \sim (V', S|_{V'}, \phi|_{V'})$  so wlog  $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$ .

**Claim 3.** If  $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$  then  $S = \phi^* S'$  for some SPQ  $S'$  on  $\text{im } \phi$  and then  $(V, S, \phi) \sim (W, S', I)$ .

**Proof of Theorem 2.** The functoriality of pullbacks needs no proof.

Now assume  $V_0 \xrightarrow{\alpha} V_1 \xrightarrow{\beta} V_2$  and that  $S$  is an SPQ on  $V_0$ . Then for every SPQ  $U$  on  $V_2$  we have, using reciprocity three times, that  $\sigma(\beta_* \alpha_* S + U) = \sigma(\alpha_* S + \beta^* U) = \sigma(S + \alpha^* \beta^* U) = \sigma(S + (\beta \alpha)^* U) = \sigma((\beta \alpha)_* S + U)$ . Hence  $\beta_* \alpha_* S = (\beta \alpha)_* S$ .

**Definition.** A commutative square as on the right is called *admissible* if  $\gamma^* \beta_* = \nu_* \mu^*$ .



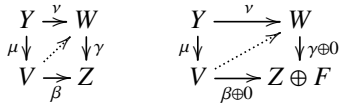
**Lemma 1.** If  $V = W = Y = Z$  and  $\beta = \gamma = \mu = \nu = I$ , the square is admissible.

**Lemma 2.** The following are equivalent:

1. A square as above is admissible.
2. The *Pairing Condition* holds. Namely, if  $S_1$  is an SPQ on  $V$  (write  $S_1 \vdash V$ ) and  $S_2 \vdash W$ , then  $\sigma(\mu^* S_1 + \nu^* S_2) = \sigma(\beta_* S_1 + \gamma_* S_2)$ .
3. The square is mirror admissible:  $\beta^* \gamma_* = \mu_* \nu^*$ .

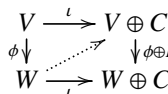
**Proof.** Using Exercises 1 and 2 below, and then using reciprocity on both sides, we have  $\forall S_1 \gamma^* \beta_* S_1 = \nu_* \mu^* S_1 \Leftrightarrow \forall S_1 \forall S_2 \sigma(\gamma^* \beta_* S_1 + S_2) = \sigma(\nu_* \mu^* S_1 + S_2) \Leftrightarrow \forall S_1 \forall S_2 \sigma(\beta_* S_1 + \gamma_* S_2) = \sigma(\mu^* S_1 + \nu^* S_2)$ , and thus 1  $\Leftrightarrow$  2. But the condition in 2 is symmetric under  $\beta \leftrightarrow \gamma, \mu \leftrightarrow \nu$ , so also 2  $\Leftrightarrow$  3.

**Lemma 3.** If the first diagram below is admissible, then so is the second.



**Lemma 4.** A pushforward by an inclusion is the do nothing operation (though note that the pushforward via an inclusion of a fully defined quadratic retains its domain of definition, which now may become partial).

**Lemma 5.** For any linear  $\phi: V \rightarrow W$ , the diagram on the right is admissible, where  $\iota$  denotes the inclusion maps.



**Proof.** Follows easily from Lemma 4.

**Definition.** If  $S$  is an SPQ with domain  $D$  and quadratic  $Q$ , the radical of  $S$  is the radical of  $Q$  considered as a fully-defined quadratic on  $D$ . Namely,  $\text{rad } S := \{u \in D: \forall v \in D, Q(u, v) = 0\}$ .

**Lemma 6.** Always,  $\phi(\text{rad } S) \subset \text{rad } \phi_* S$ .

**Proof.** Pick  $w \in \phi(\text{rad } S)$  and repeat the proof of Theorem 1' but now considering quadruples  $(V, S, \phi, \nu)$ , where  $(V, S, \phi)$  are as before and  $\nu \in \text{rad } S$  satisfies  $\phi(\nu) = w$ . Clearly our initial triple  $(V, S, \phi)$  can be extended to such a quadruple, and it is easy to repeat the steps of the proof of Theorem 1' extending everything to such quadruples.

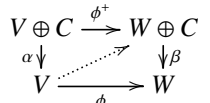
We have to acknowledge that our proof of Lemma 6 is ugly. We wish we had a cleaner one.

**Exercise 3.** Show that if two SPQ's  $S_1$  and  $S_2$  on  $V \oplus A$  satisfy  $A \subset \text{rad } S_i$  and  $\sigma(S_1 + \pi^* U) = \sigma(S_2 + \pi^* U)$  for every quadratic  $U$  on  $V$ , where  $\pi: V \oplus A \rightarrow V$  is the projection, then  $S_1 = S_2$ .

**Exercise 4.** Show that if  $\phi: V \rightarrow W$  is surjective and  $Q$  is a quadratic on  $W$ , then  $\sigma(Q) = \sigma(\phi^* Q)$ .

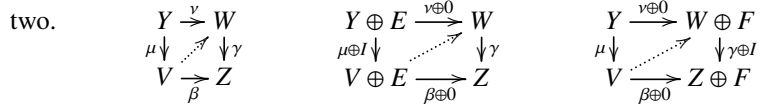
**Exercise 5.** Show that always,  $\phi_* \phi^* S = S|_{\text{im } \phi}$ .

**Lemma 7.** For any linear  $\phi: V \rightarrow W$ , the diagram on the right is admissible, where  $\phi^+ := \phi \oplus I$  and  $\alpha$  and  $\beta$  denote the projection maps.

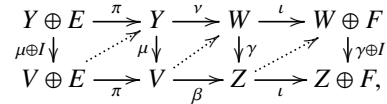


**Proof.** Let  $S$  be an SPQ on  $V$ . Clearly  $C \subset \beta^* \phi_* S$ . Also,  $C \subset \text{rad } \alpha^* S$  so by Lemma 6,  $C = \phi^+(C) \subset \phi^+(\text{rad } \alpha^* S) \subset \text{rad } \phi_* \alpha^* S$ . Hence using Exercise 3, it is enough to show that  $\sigma(\phi_* \alpha^* S + \beta^* U) = \sigma(\beta^* \phi_* S + \beta^* U)$  for every  $U$  on  $W$ . Indeed,  $\sigma(\phi_* \alpha^* S + \beta^* U) \stackrel{(1)}{=} \sigma(\beta_* \phi_* \alpha^* S + U) \stackrel{(2)}{=} \sigma(\phi_* \alpha_* \alpha^* S + U) \stackrel{(3)}{=} \sigma(\phi_* S + U) \stackrel{(4)}{=} \sigma(\beta^*(\phi_* S + U)) \stackrel{(5)}{=} \sigma(\beta^* \phi_* S + \beta^* U)$ , using (1) reciprocity, (2) the commutativity of the diagram and the functoriality of pushing, (3) Exercise 5, (4) Exercise 4, and (5) the additivity of pullbacks.

**Lemma 8.** If the first diagram below is admissible, then so are the other two.



**Proof.** In the diagram



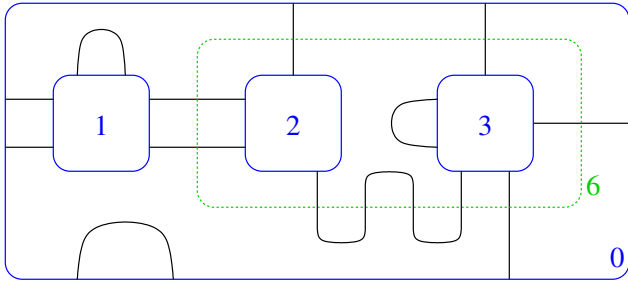
with  $\pi$  marking projections and  $\iota$  inclusions, the left square is admissible by Lemma 7, the middle square by assumption, and the right square by Lemma 5. Along with the functoriality of pushforwards this shows the admissibility of both the left and the right  $1 \times 2$  subrectangles, and these are the diagrams we wanted.

**Proof of Theorem 3.** Decompose  $Z = A \oplus E \oplus F \rightarrow A \oplus C \oplus F$  where  $A = \text{im } \beta \cap \text{im } \gamma$ ,  $\text{im } \beta = A \oplus B$ , and  $\text{im } \gamma = A \oplus C$ . Write  $V \simeq A \oplus B \oplus E$  with  $\beta = I$  on  $A \oplus B$  yet  $\beta = 0$  on  $E$ , and write  $W \simeq A \oplus C \oplus F$  with  $\gamma = I$  on  $A \oplus C$  yet  $\gamma = 0$  on  $F$ . Then  $Y = V \oplus_Z W \simeq A \oplus E \oplus F$  and our square is as shown on the right, with all maps equal to  $I$  on like-named summands and equal to 0 on non-like-named summands. But this diagram is admissible: build it up using Lemma 1 for the  $A$ 's, and then Lemma 8 for  $E$  and  $C$ , and then again Lemma 8 along with the mirror property of Lemma 2 for  $B$  and  $F$ , and then Lemma 3 for  $D$ .

To prove Theorem 4, given three<sup>1</sup> SPQ's  $S_1, S_2$ , and  $S_3$ , we need to show that planar-multiplying them in two steps, first using a planar connection diagram  $D_I$  ( $I$  for Inner) to yield  $S_6 = S(D_I)(S_2, S_3)$  and then using a second planar connection diagram  $D_O$  ( $O$  for Outer) to yield  $S(D_O)(S_1, S_6)$ , gives the same answer as multiplying them all at once using the composition planar connection diagram  $D_B = D_O \circ_6 D_I$  ( $B$  for Big) to yield  $S(D_B)(S_1, S_2, S_3)$ .<sup>2</sup> An example should help:

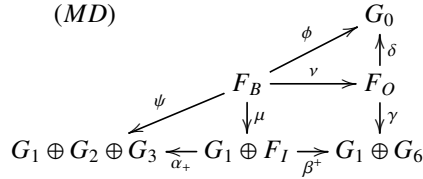
<sup>1</sup>Truly, we need the same for any number of input SPQ's that are divided into two groups, "multiply in the first step" and "multiply in the second step". But there's no added difficulty here, only an added notational complexity.

<sup>2</sup>Aren't we sassy? We picked "6" for the name of the product of "2" and "3".



In this example, if you ignore the dotted green line (marked “6”), you see the planar connection diagram  $D_B$ , which has three inputs (1,2,3) and a single output, the cycle 0. If you only look inside the green line, you see  $D_I$ , with inputs 2 and 3 and an output cycle 6. If you ignore the inside of 6 you see  $D_O$ , with inputs 1 and 6 and output cycle 0.

Let  $F_B$  (Big Faces) denote the vector space whose basis are the faces of  $D_B$ , let  $F_I$  (Inner Faces) be the space of faces of  $D_I$ , and let  $F_O$  (Outer Faces) be the space of faces of  $D_O$ . Let  $G_1, G_2, G_3, G_6$ , and  $G_0$  be the spaces of gaps (edges) along the cycles 1,2,3,6, and 0, respectively. Let  $\psi := \psi_{D_B}$  and  $\phi := \phi^{D_B}$  be the maps defining  $\mathcal{S}(D_B)$  and let  $\gamma := \psi_{D_O}$  and  $\delta := \phi^{D_O}$  be the maps defining  $\mathcal{S}(D_O)$ . Further, let  $\alpha := \psi_{D_I} : F_I \rightarrow G_2 \oplus G_3$  and  $\beta := \phi^{D_I} : F_I \rightarrow G_6$  be the maps defining  $\mathcal{S}(D_I)$ , and let  $\alpha_+ := I \oplus \alpha$  and  $\beta^+ := I \oplus \beta$  be the extensions of  $\alpha$  and  $\beta$  by an identity on an extra factor of  $G_1$ , so that  $\beta^+ \alpha_+^* = I_{G_1} \oplus \mathcal{S}(D_I)$ . Let  $\mu$  map any big face to the sum of  $G_1$  gaps around it, plus the sum of the inner faces it contains. Let  $\nu$  map any big face to the sum of the outer faces it contains. It is easy to see that the master diagram (MD) shown on the right, made of all of these spaces and maps, is commutative.



**Claim.** The bottom right square of (MD) is an equalizer square, namely  $F_B \simeq EQ(\beta^+, \gamma)$ . Hence  $\nu_* \mu^* = \gamma^* \beta^+$ .

**Proof.** A big face (an element of  $F_B$ ) is a sum of outer faces  $f_o$  and a sum of inner faces  $f_i$ , and it has a boundary  $g_1$  on input cycle 1, such that the boundary of the outer pieces  $f_o$  is equal to the boundary of the inner pieces  $f_i$  plus  $g_1$ . That matches perfectly with the definition of the equalizer:  $EQ(\beta^+, \gamma) = \{(g_1, f_i, f_o) : \beta^+(g_1, f_i) = \gamma(f_o)\} = \{(g_1, f_i, f_o) : \gamma(f_o) = (g_1, \beta(f_i))\}$ .  $\square$

**Proof of Theorem 4.** With notation as above, with the example above (which is general enough), and with the claim above, and also using functoriality, we have  $\mathcal{S}(D_B) = \phi_* \psi^* = \delta_* \nu_* \mu^* \alpha_+^* = \delta_* \gamma^* \beta^+ \alpha_+^* = \mathcal{S}(D_O) \circ (I_{G_1} \oplus \mathcal{S}(D_I))$ , as required.  $\square$

**Proof of Theorem 5.** We need to verify the Reidemeister moves and that was done in the computational section, and the statement about the restriction to links, which is easy: simply assemble an  $n$ -crossing knot using an  $n$ -input planar connection diagram, and the formulas clearly match.  $\square$

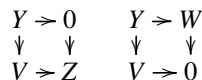
**Further Homework.**

**Exercise 6.** By taking  $U = 0$  in the reciprocity statement, prove that always  $\sigma(\phi_* S) = \sigma(S)$ . But that seems wrong, if  $\phi = 0$ . What saves the day?

**Exercise 7.** By taking  $S = 0$  in the reciprocity statement, prove that always  $\sigma(\phi^* U) = \sigma(U)$ . But wait, this is nonsense! What went wrong?

**Exercise 8.** Given  $\phi : V \rightarrow W$  and a subspace  $D \subset V$ , show that there is a unique subspace  $\phi_* D \subset W$  such that for every quadratic  $Q$  on  $W$ ,  $\sigma(\phi^* Q|_D) = \sigma(Q|_{\phi_* D})$ .

**Exercise 9.** When are diagrams as on the right equalizer diagrams? What then do we learn from Theorem 3?



**Exercise 10.** There are 11 types of irreducible commutative squares:  $1 \rightrightarrows 0, 0 \rightrightarrows 1, 0 \rightrightarrows 0, 0 \rightrightarrows 0, 1 \rightrightarrows 1, 0 \rightrightarrows 1, 0 \rightrightarrows 1, 0 \rightrightarrows 0, 0 \rightrightarrows 0, 0 \rightrightarrows 0, 1 \rightrightarrows 0, 0 \rightrightarrows 1, 0 \rightrightarrows 0, 0 \rightrightarrows 1, 0 \rightrightarrows 1, 1 \rightrightarrows 1$ . Show that pushing commutes with pulling for all but four of them. Compare with the statement of Theorem 3.

**Exercise 11.** Prove that a square is admissible iff it is an equalizer square, with an additional direct summand  $A$  added to the  $Y$  term, and with the maps  $\mu$  and  $\nu$  extended by 0 on  $A$ .

**Exercise 12.** Prove that the direct sum of two admissible squares is admissible. *Warning:* Harder than it seems! Not all quadratics on  $V_1 \oplus V_2$  are direct sums of quadratics on  $V_1$  and on  $V_2$ .

**Exercise 13.** Given a quadratic  $Q$  on a space  $V$ , let  $\pi$  be the projection  $V \rightarrow V/\text{rad}(Q)$  and show that  $\pi_* Q = Q/\text{rad}(Q)$ , with the obvious definition for the latter.

**Exercise 14.** Show that for any partial quadratic  $Q$  on a space  $W$  there exists a space  $A$  and a fully-defined quadratic  $F$  on  $W \oplus A$  such that  $\pi_* F = Q$ , where  $\pi : W \oplus A \rightarrow W$  is the projection (these are not unique). Furthermore, if  $\phi : V \rightarrow W$ , then  $\phi^* Q = \pi_* \phi_+^* F$ , where  $\phi_+ = \phi \oplus I : V \oplus A \rightarrow W \oplus A$  and  $\pi$  also denotes the projection  $V \oplus A \rightarrow V$ .

**Solutions / Hints.**

- Hint for 1.** On a vector in the domain of one but not the other, the signature will rise or lower the signature.
- Hint for 2.**  $Q_{\text{DW}}$  is idempotent in  $\mathbb{Q}$ .
- Hint for 3.** It's enough to test that against  $\mathbb{U}$  with  $\mathbb{U} = (\mathbb{U})$ .
- Hint for 4.** The "shift" part of  $\mathcal{S}$  is  $(\mathcal{S})$ .
- Hint for 7.**  $\phi$  is the pairing  $\mathcal{S}(\phi)$  and  $\phi$  is the pairing  $\mathcal{S}(\phi)$ .
- Hint for 10.** The exceptions are  $\begin{matrix} 10 & 00 & 10 \\ 01 & 01 & 00 \end{matrix}$  and  $\begin{matrix} 11 & 01 \end{matrix}$ .
- Hint for 12.** Use Exercise 11.