



Computation without Representation

$\omega\epsilon\beta$: <http://drorbn.net/o19/>

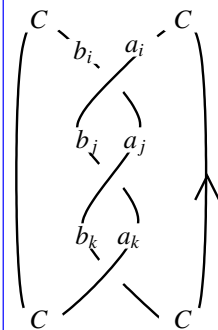
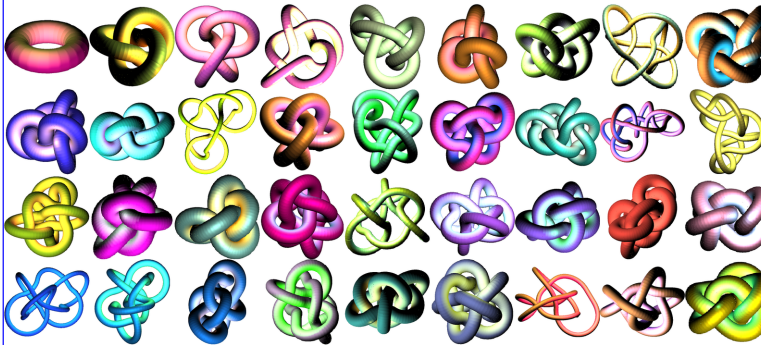
Abstract. A major part of “quantum topology” is the definition and computation of various knot invariants by carrying out computations in quantum groups. Traditionally these computations are carried out “in a representation”, but this is very slow: one has to use tensor powers of these representations, and the dimensions of powers grow exponentially fast.

In my talk, I will describe a direct method for carrying out such computations without having to choose a representation and explain why in many ways the results are better and faster. The two key points we use are a technique for composing infinite-order “perturbed Gaussian” differential operators, and the little-known fact that every semi-simple Lie algebra can be approximated by solvable Lie algebras, where computations are easier.

KiW 43 Abstract ($\omega\epsilon\beta$ /kiw). Whether or not you like the formulas on this page, they describe the strongest truly computable knot invariant we know. (experimental analysis @ $\omega\epsilon\beta$ /kiw)

Knotted Candies

$\omega\epsilon\beta$ /kc



The Yang-Baxter Technique. Given an algebra U (typically $\hat{U}(\mathfrak{g})$ or $\hat{U}_q(\mathfrak{g})$) and elements

$$R = \sum a_i \otimes b_i \in U \otimes U \quad \text{and} \quad C \in U,$$

$$\text{form} \quad Z = \sum_{i,j,k} C a_i b_j a_k C^2 b_i a_j b_k C.$$

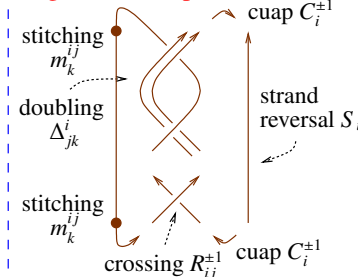
Problem. Extract information from Z .

The Dogma. Use representation theory. In principle finite, but *slow*.

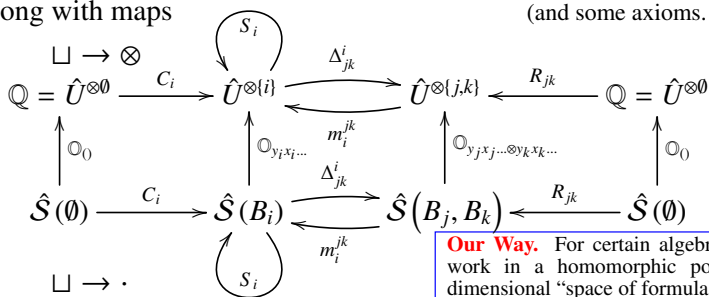
A Knot Theory Portfolio.

- Has operations $\sqcup, m_k^{ij}, \Delta_{jk}^i, S_i$.
- All tangles are generated by $R^{\pm 1}$ and $C^{\pm 1}$ (so “easy” to produce invariants).
- Makes some knot properties (“genus”, “ribbon”) become “definable”.

Tangles and Operations

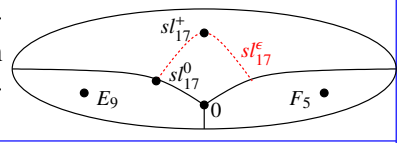


A “Quantum Group” Portfolio consists of a vector space U along with maps (and some axioms...)

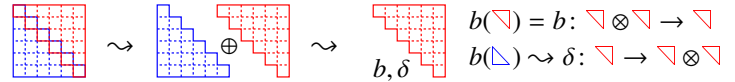


Our Way. For certain algebras, work in a homomorphic polynomial “space of formulas”.

The (fake) moduli of Lie algebras on V , a quadratic variety in $(V^*)^{\otimes 2} \otimes V$ is on the right. We care about $sl_{17}^k := sl_{17}^\epsilon / (\epsilon^{k+1} = 0)$.



Solvable Approximation. In gl_n , half is enough! Indeed $gl_n \oplus \mathfrak{a}_n = \mathcal{D}(\nabla, b, \delta)$:

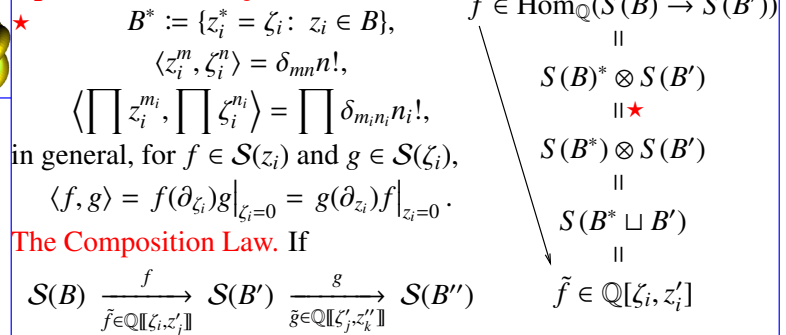


Now define $gl_n^\epsilon := \mathcal{D}(\nabla, b, \epsilon\delta)$. Schematically, this is $[\nabla, \nabla] = \nabla$, $[\Delta, \Delta] = \epsilon\Delta$, and $[\nabla, \Delta] = \Delta + \epsilon\nabla$. The same process works for all semi-simple Lie algebras, and at $\epsilon^{k+1} = 0$ always yields a solvable Lie algebra.

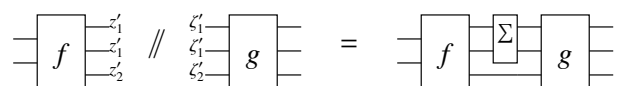
CU and QU. Starting from sl_2 , get $CU_\epsilon = \langle y, a, x, t \rangle / ([t, -] = 0, [a, y] = -y, [a, x] = x, [x, y] = 2\epsilon a - t)$. Quantize using standard tools (I’m sorry) and get $QU_\epsilon = \langle y, a, x, t \rangle / ([t, -] = 0, [a, y] = -y, [a, x] = x, xy - e^{\hbar\epsilon}yx = (1 - T e^{-2\hbar\epsilon a})/\hbar)$.

PBW Bases. The U ’s we care about always have “Poincaré-Birkhoff-Witt” bases; there is some finite set $B = \{y, x, \dots\}$ of “generators” and isomorphisms $\mathbb{O}_{y,x,\dots} : \hat{S}(B) \rightarrow U$ defined by “ordering monomials” to some fixed y, x, \dots order. The quantum group portfolio now becomes a “symmetric algebra” portfolio, or a “power series” portfolio.

Operations are Objects.



then $(\tilde{f}/\tilde{g}) = (\tilde{g} \circ f) = (\tilde{g}|_{z'_j \rightarrow \partial_{z'_j} \tilde{f}})_{z'_j=0} = (\tilde{f}|_{z'_j \rightarrow \partial_{z'_j} \tilde{g}})_{z'_j=0}$



1. The 1-variable identity map $I: S(z) \rightarrow S(z)$ is given by $\tilde{I}_1 = e^{z\zeta}$ and the n -variable one by $\tilde{I}_n = e^{z_1\zeta_1 + \dots + z_n\zeta_n}$.

$$\tilde{I}_1 = \text{[diagram]} + \text{[diagram]} + \frac{1}{2} \text{[diagram]} + \frac{1}{6} \text{[diagram]} + \dots$$

2. The “archetypal multiplication map $m_k^{ij}: S(z_i, z_j) \rightarrow S(z_k)$ ” has $\tilde{m} = e^{-z_k(\zeta_i + \zeta_j)}$.

3. The “archetypal coproduct $\Delta_{jk}^i: S(z_i) \rightarrow S(z_j, z_k)$ ”, given by $z_i \rightarrow z_j + z_k$ or $\Delta z = z \otimes 1 + 1 \otimes z$, has $\tilde{\Delta} = e^{(z_j + z_k)\zeta_i}$.

4. R -matrices tend to have terms of the form $e^{h_{y_1 x_2}} \in \mathcal{U}_q \otimes \mathcal{U}_q$. The “baby R -matrix” is $\tilde{R} = e^{h_{yx}} \in S(y, x)$.

5. The “Weyl form of the canonical commutation relations” states that if $[y, x] = tI$ then $e^{\xi x} e^{\eta y} = e^{\eta y} e^{\xi x} e^{-\eta \xi t}$. So with

$$SW_{xy} \left(S(y, x) \xrightarrow{\mathcal{O}_{yx}} \mathcal{U}(y, x) \right) \text{ we have } \widetilde{SW}_{xy} = e^{\eta y + \xi x - \eta \xi t}.$$

The Real Thing. In the algebra QU_ϵ , over $\mathbb{Q}[[\hbar]]$ using the $yaxt$ order, $T = e^{\hbar t}$, $\bar{T} = T^{-1}$, $\mathcal{A} = e^\alpha$, and $\bar{\mathcal{A}} = \mathcal{A}^{-1}$, we have

$$\tilde{R}_{ij} = e^{\hbar(y_i x_j - t_i a_j)} \left(1 + \epsilon \hbar \left(a_i a_j - \hbar^2 y_i^2 x_j^2 / 4 \right) + O(\epsilon^2) \right)$$

in $\mathcal{S}(B_i, B_j)$, and in $\mathcal{S}(B_1^*, B_2^*, B)$ we have

$$\tilde{m} = e^{(\alpha_1 + \alpha_2) a + \eta_2 \xi_1 (1-T) / \hbar + (\xi_1 \bar{\mathcal{A}}_2 + \xi_2) x + (\eta_1 + \eta_2 \bar{\mathcal{A}}_1) y} \left(1 + \epsilon \lambda + O(\epsilon^2) \right),$$

where $\lambda = \frac{2a\eta_2 \xi_1 T + \eta_2^2 \xi_1^2 (3T^2 - 4T + 1) / 4\hbar - \eta_2 \xi_1^2 (3T - 1) x \bar{\mathcal{A}}_2 / 2 - \eta_2^2 \xi_1 (3T - 1) y \bar{\mathcal{A}}_1 / 2 + \eta_2 \xi_1 x y \hbar \bar{\mathcal{A}}_1 \bar{\mathcal{A}}_2}{}$.

Finally,

$$\tilde{\Delta} = e^{\tau(t_1 + t_2) + \eta(y_1 + T_1 y_2) + \alpha(a_1 + a_2) + \xi(x_1 + x_2)} (1 + O(\epsilon)) \in \mathcal{S}(B^*, B_1, B_2),$$

$$\text{and } \tilde{S} = e^{-\tau t - \alpha a - \eta \xi (1-T) \mathcal{A} / \hbar - T \eta y \mathcal{A} - \xi x \mathcal{A}} (1 + O(\epsilon)) \in \mathcal{S}(B^*, B).$$

The Zipping Issue.

(between unbound and bound lies half-zipped).



Zipping. If $P(\zeta^j, z_i)$ is a polynomial, or whenever otherwise convergent, set $\langle P(\zeta^j, z_i) \rangle_{(\zeta^j)} = P(\partial_{z_j}, z_i) \Big|_{z_i=0}$. (E.g., if $P = \sum a_{nm} \zeta^n z^m$ then $\langle P \rangle_{\zeta} = \sum a_{nm} \partial_z^n z^m \Big|_{z=0} = \sum n! a_{nm}$).

The Zipping / Contraction Theorem. If $P = P(\zeta^j, z_i)$ has a finite ζ -degree and the y 's and the q 's are "small" then

$$\left\langle P e^{c + \eta^j z_i + y_j \zeta^j + q_j^k z_k} \right\rangle_{(\zeta^j)} = \det(\tilde{q}) e^{c + \eta^j \tilde{q}_i^k y_k} \left\langle P \Big|_{z_i \rightarrow \tilde{q}_i^k (z_k + y_k)} \right\rangle_{(\zeta^j)}$$

where \tilde{q} is the inverse matrix of $1 - q$: $(\delta_j^i - q_j^i) \tilde{q}_k^j = \delta_k^i$.

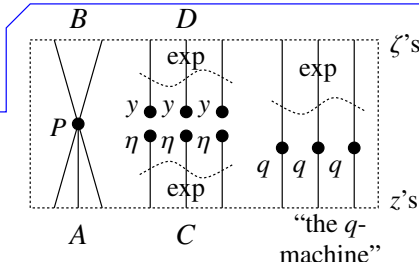
Exponential Reservoirs. The true Hilbert hotel is exp! Remove one x from an "exponential reservoir" of x 's and you are left with the same exponential reservoir:

$$e^x = \left[\dots + \frac{xxxxx}{120} + \dots \right] \xrightarrow{\partial_x} \left[\dots + \frac{xxxxx}{120} + \dots \right] = (e^x)' = e^x,$$

and if you let each element choose left or right, you get twice the same reservoir:

$$e^x \xrightarrow{x \rightarrow x_l + x_r} e^{x_l + x_r} = e^{x_l} e^{x_r}.$$

A Graphical Proof. Glue top to bottom on the right, in all possible ways. Several scenarios occur:



1. Start at A, go through the q -machine $k \geq 0$ times, stop at B. Get $\langle P(\zeta, \sum_{k \geq 0} q^k z) \rangle = \langle P(\zeta, \tilde{q} z) \rangle$.
2. Loop through the q -machine and swallow your own tail. Get $\exp(\sum q^k / k) = \exp(-\log(1 - q)) = \tilde{q}$.
3. ...

By the reservoir splitting principle, these scenarios contribute multiplicatively. \square

Implementation. $(\mathbb{E}[Q, P] \text{ means } e^Q P)$ $\omega\epsilon\beta/\text{Zip}$

```
Zip_gs_List @E [Q, P] :=
Module[{z, zs, c, ys, ns, qt, zrule, grule},
  zs = Table[z^*, {z, zs}];
  c = Q /. Alternatives @@ (zs | z) -> 0;
  ys = Table[partial_z (Q /. Alternatives @@ zs -> 0), {z, zs}];
  ns = Table[partial_z (Q /. Alternatives @@ z -> 0), {z, zs}];
  qt = Inverse@Table[KD_z, z^* - partial_z, z, {z, zs}, {z, zs}];
  zrule = Thread[zs -> qt. (zs + ys)];
  grule = Thread[z -> z + ns.qt];
  Simplify /@
  E[c + ns.qt.ys, Det[qt] Zip_gs[P /. (zrule | grule)]];];
```

Real Zipping is a minor mess, and is done in two phases:

	τa -phase		ξy -phase	
ζ -like variables	τ	a	ξ	y
z -like variables	t	α	x	η

Already at $\epsilon = 0$ we get the best known formulas for the Alexander polynomial!

Generic Docility. A "docile perturbed Gaussian" in the variables $(z_i)_{i \in S}$ over the ring R is an expression of the form

$$e^{q^{ij} z_i z_j} P = e^{q^{ij} z_i z_j} \left(\sum_{k \geq 0} \epsilon^k P_k \right),$$

where all coefficients are in R and where P is a "docile series": $\deg P_k \leq 4k$.

Our Docility. In the case of QU_ϵ , all invariants and operations are of the form $e^{L+Q} P$, where

- L is a quadratic of the form $\sum l_{z\zeta} z \zeta$, where z runs over $\{t_i, \alpha_i\}_{i \in S}$ and ζ over $\{\tau_i, a_i\}_{i \in S}$, with integer coefficients $l_{z\zeta}$.
- Q is a quadratic of the form $\sum q_{z\zeta} z \zeta$, where z runs over $\{x_i, \eta_i\}_{i \in S}$ and ζ over $\{\xi_i, y_i\}_{i \in S}$, with coefficients $q_{z\zeta}$ in the ring R_S of rational functions in $\{T_i, \mathcal{A}_i\}_{i \in S}$.
- P is a docile power series in $\{y_i, a_i, x_i, \eta_i, \xi_i\}_{i \in S}$ with coefficients in R_S , and where $\deg(y_i, a_i, x_i, \eta_i, \xi_i) = (1, 2, 1, 1, 1)$.

Docility Matters! The rank of the space of docile series to ϵ^k is polynomial in the number of variables $|S|$. **!!!!**

- At $\epsilon^2 = 0$ we get the Rozansky-Overbay [Ro1, Ro2, Ro3, Ov] invariant, which is stronger than HOMFLY-PT polynomial and Khovanov homology taken together!
- In general, get "higher diagonals in the Melvin-Morton-Rozansky expansion of the coloured Jones polynomial" [MM, BNG], but why spoil something good?

[BNG] D. Bar-Natan and S. Garoufalidis, *On the Melvin-Morton-Rozansky conjecture*, Invent. Math. **125** (1996) 103–133.

[BV] D. Bar-Natan and R. van der Veen, *A Polynomial Time Knot Polynomial*, arXiv:1708.04853.

[Fa] L. Faddeev, *Modular Double of a Quantum Group*, arXiv:math/9912078.

[GR] S. Garoufalidis and L. Rozansky, *The Loop Expansion of the Kontsevich Integral, the Null-Move, and S-Equivalence*, arXiv:math.GT/0003187.

[MM] P. M. Melvin and H. R. Morton, *The coloured Jones function*, Commun. Math. Phys. **169** (1995) 501–520.

[Ov] A. Overbay, *Perturbative Expansion of the Colored Jones Polynomial*, University of North Carolina PhD thesis, $\omega\epsilon\beta/\text{Ov}$.

[Qu] C. Quesne, *Jackson's q-Exponential as the Exponential of a Series*, arXiv:math-ph/0305003.

[Ro1] L. Rozansky, *A contribution of the trivial flat connection to the Jones polynomial and Witten's invariant of 3d manifolds, I*, Comm. Math. Phys. **175-2** (1996) 275–296, arXiv:hep-th/9401061.

[Ro2] L. Rozansky, *The Universal R-Matrix, Braid Representation and the Melvin-Morton Expansion of the Colored Jones Polynomial*, Adv. Math. **134-1** (1998) 1–31, arXiv:q-alg/9604005.

[Ro3] L. Rozansky, *A Universal U(1)-RCC Invariant of Links and Rationality Conjecture*, arXiv:math/0201139.

[Za] D. Zagier, *The Dilogarithm Function*, in Cartier, Moussa, Julia, and Vanhove (eds) *Frontiers in Number Theory, Physics, and Geometry II*. Springer, Berlin, Heidelberg, and $\omega\epsilon\beta/\text{Za}$.

References.



"God created the knots, all else in topology is the work of mortals."
Leopold Kronecker (modified)

The Algebras H and H^* . Let $q = e^{\hbar\epsilon\gamma}$ and set $H = \langle a, x \rangle / ([a, x] = \gamma x)$ with

$$A = e^{-\hbar\epsilon a}, \quad xA = qAx, \quad S_H(a, A, x) = (-a, A^{-1}, -A^{-1}x),$$

$$\Delta_H(a, A, x) = (a_1 + a_2, A_1A_2, x_1 + A_1x_2)$$

and dual $H^* = \langle b, y \rangle / ([b, y] = -\epsilon y)$ with

$$B = e^{-\hbar y b}, \quad By = qyB, \quad S_{H^*}(b, B, y) = (-b, B^{-1}, -yB^{-1}),$$

$$\Delta_{H^*}(b, B, y) = (b_1 + b_2, B_1B_2, y_1B_2 + y_2).$$

Pairing by $(a, x)^* = (b, y) (\Rightarrow \langle B, A \rangle = q)$ making $\langle y^j b^i, a^j x^k \rangle = \delta_{ij} \delta_{kl} j! [k]_q!$ so $R = \sum \frac{y^k b^i \otimes a^j x^k}{j! [k]_q!}$.

The Algebra QU . Using the Drinfel'd double procedure, $QU_{\gamma, \epsilon} := H^{*cop} \otimes H$ with $(\phi f)(\psi g) = \langle \psi_1 S^{-1} f_3 \rangle \langle \psi_3, f_1 \rangle (\phi \psi_2)(f_2 g)$ and

$$S(y, b, a, x) = (-B^{-1}y, -b, -a, -A^{-1}x),$$

$$\Delta(y, b, a, x) = (y_1 + y_2 B_1, b_1 + b_2, a_1 + a_2, x_1 + A_1 x_2).$$

Note also that $t := \epsilon a - \gamma b$ is central and can replace b , and set $QU = QU_\epsilon = QU_{1, \epsilon}$.

The 2D Lie Algebra. One may show* that if $[a, x] = \gamma x$ then $e^{\xi x} e^{aa} = e^{aa} e^{\gamma a \xi x}$. Ergo with

$$SW_{ax} \left(\begin{array}{c} \curvearrowright \\ \mathcal{S}(a, x) \end{array} \begin{array}{c} \xrightarrow{\mathcal{O}_{ax}} \\ \mathcal{U}(a, x) \\ \xleftarrow{\mathcal{O}_{xa}} \end{array} \right)$$

we have $\widetilde{SW}_{ax} = e^{aa + e^{-\gamma a} \xi x}$.

* Indeed $xa = (a - \gamma)x$ thus $xa^n = (a - \gamma)^n x$ thus $x e^{aa} = e^{a(a-\gamma)} x = e^{-\gamma a} e^{aa} x$ thus $x^n e^{aa} = e^{aa} (e^{-\gamma a})^n x^n$ thus $e^{\xi x} e^{aa} = e^{aa} e^{\gamma a \xi x}$.

Faddeev's Formula (In as much as we can tell, first appeared without proof in Faddeev [Fa], rediscovered and proven in Quesne [Qu], and again with easier proof, in Zagier [Za]). With $[n]_q := \frac{q^n - 1}{q - 1}$, with $[n]_q! := [1]_q [2]_q \cdots [n]_q$ and with $e_q^x := \sum_{n \geq 0} \frac{x^n}{[n]_q!}$, we have

$$\log e_q^x = \sum_{k \geq 1} \frac{(1-q)^k x^k}{k(1-q^k)} = x + \frac{(1-q)^2 x^2}{2(1-q^2)} + \dots$$

Proof. We have that $e_q^x = \frac{e^{qx} - e^{-qx}}{qx - x}$ ("the q -derivative of e_q^x is itself"), and hence $e_q^{qx} = (1 + (1-q)x)e_q^x$, and

$$\log e_q^{qx} = \log(1 + (1-q)x) + \log e_q^x.$$

Writing $\log e_q^x = \sum_{k \geq 1} a_k x^k$ and comparing powers of x , we get $q^k a_k = -(1-q)^k/k + a_k$, or $a_k = \frac{(1-q)^k}{k(1-q^k)}$. □

A Full Implementation.

$\omega \epsilon \beta / \text{Full}$

Utilities

```
CF[sd_SeriesData] := MapAt[CF, sd, 3];
CF[_] := ExpandDenominator@ExpandNumerator@Together[
  Expand[_] /. e^-v -> e^v /. e^x -> e^CF[x];
```

```
Kδ /: Kδ_{i,j} := If[i===j, 1, 0];
E /: E[L1_, Q1_, P1_] ≡ E[L2_, Q2_, P2_] :=
  CF[L1 == L2] ∧ CF[Q1 == Q2] ∧ CF[Normal[P1 - P2] == 0];
E /: E[L1_, Q1_, P1_] E[L2_, Q2_, P2_] :=
  E[L1 + L2, Q1 + Q2, P1 * P2];
E[L_, Q_, P_]_{$k} := E[L, Q, Series[Normal@P, {ε, 0, $k}]];
```

Zip and Bind

```
{t*, b*, y*, a*, x*, z*} = {τ, β, η, α, ξ, ζ};
{τ*, β*, η*, α*, ξ*, ζ*} = {t, b, y, a, x, z};
(u_{-i})_* := (u*)_i;
```

```
collect[sd_SeriesData, _] :=
  MapAt[collect[#, _] &, sd, 3];
collect[_] := Collect[_];
Zip[_][P_] := P; Zip[_][P_] :=
  (collect[P // Zip[_]] /. f_{-} -> ∂_{ε*, d} f) /. ε* -> 0
QZip[_List@E[L_, Q_, P_] :=
```

```
Module[{ξ, z, zs, c, ys, ηs, qt, zrule, grule},
  zs = Table[ξ*, {ξ, ζs}];
  c = CF[Q /. Alternatives @@ (ξs ∪ zs) -> 0];
  ys = CF@Table[∂_ε (Q /. Alternatives @@ zs -> 0), {ξ, ζs}];
  ηs = CF@Table[∂_z (Q /. Alternatives @@ ξs -> 0), {z, zs}];
  qt = CF@Inverse@Table[Kδ_{z, ξ*} - ∂_{z, ε} Q, {ξ, ζs}, {z, zs}];
  zrule = Thread[zs -> CF[qt.(zs + ys)]];
  grule = Thread[ξs -> ξs + ηs.qt];
  CF /@ E[L, c + ηs.qt.ys,
    Det[qt] Zip[_][P /. (zrule ∪ grule)]]];
```

```
U21 = {B_{-}^{p_{-}} -> e^{-p_{-} \hbar y b_i}, B_{-}^{p_{-}} -> e^{-p_{-} \hbar \gamma b}, T_{-}^{p_{-}} -> e^{p_{-} \hbar t_i},
  T_{-}^{p_{-}} -> e^{p_{-} \hbar t}, \mathcal{A}_{-}^{p_{-}} -> e^{p_{-} \gamma \alpha_i}, \mathcal{A}_{-}^{p_{-}} -> e^{p_{-} \gamma \alpha}};
L2U = {e^{c_{-} \cdot b_i + d_{-}} -> B_{-}^{-c/(h \gamma)} e^d, e^{c_{-} \cdot b + d_{-}} -> B^{-c/(h \gamma)} e^d,
  e^{c_{-} \cdot t_i + d_{-}} -> T_{-}^{c/h} e^d, e^{c_{-} \cdot t + d_{-}} -> T^{c/h} e^d,
  e^{c_{-} \cdot \alpha_i + d_{-}} -> \mathcal{A}_{-}^{c/\gamma} e^d, e^{c_{-} \cdot \alpha + d_{-}} -> \mathcal{A}^{c/\gamma} e^d,
  e^{\epsilon_{-}} -> e^{Expand[_]}};
```

```
LZip[_List@E[L_, Q_, P_] :=
  Module[{ξ, z, zs, c, ys, ηs, lt, zrule, L1, L2, Q1, Q2},
  zs = Table[ξ*, {ξ, ζs}];
  c = L /. Alternatives @@ (ξs ∪ zs) -> 0;
  ys = Table[∂_ε (L /. Alternatives @@ zs -> 0), {ξ, ζs}];
  ηs = Table[∂_z (L /. Alternatives @@ ξs -> 0), {z, zs}];
  lt = Inverse@Table[Kδ_{z, ξ*} - ∂_{z, ε} L, {ξ, ζs}, {z, zs}];
  zrule = Thread[zs -> lt.(zs + ys)];
  L2 = (L1 = c + ηs.zs /. zrule) /. Alternatives @@ zs -> 0;
  Q2 = (Q1 = Q /. U21 /. zrule) /. Alternatives @@ zs -> 0;
  CF /@ E[L2, Q2, Det[lt] e^{-L2-Q2}
    Zip[_][e^{L1+Q1} (P /. U21 /. zrule)]] // . L2U];
```

```
B_{-}[L_, R_] := LR;
B_{is}_{-}[L_{-E}, R_{-E}] := Module[{n}, Times[
  L /. Table[{v: b | B | t | T | a | x | y}_i -> v_{nei}, {i, {is}}],
  R /. Table[{v: β | τ | α | \mathcal{A} | ξ | η}_i -> v_{nei}, {i, {is}}]
] // LZipJoin@Table[{β_{nei}, τ_{nei}, a_{nei}}, {i, {is}}] //
  QZipJoin@Table[{ξ_{nei}, y_{nei}}, {i, {is}}];
B_{is}_{-}[L_, R_] := B_{is}[L, R];
```

E morphisms with domain and range.

```
B_{is}_{-}List[E_{d1 -> r1}[L1_, Q1_, P1_], E_{d2 -> r2}[L2_, Q2_, P2_] :=
  E[(d1 ∪ Complement[d2, is]) -> (r2 ∪ Complement[r1, is])] @@
  B_{is}[E[L1, Q1, P1], E[L2, Q2, P2]];
E_{d1 -> r1}[L1_, Q1_, P1_] // E_{d2 -> r2}[L2_, Q2_, P2_] :=
  B_{r1} ∩_{d2} [E_{d1 -> r1}[L1, Q1, P1], E_{d2 -> r2}[L2, Q2, P2]];
E_{d1 -> r1}[L1_, Q1_, P1_] ≡ E_{d2 -> r2}[L2_, Q2_, P2_] ^:=
  (d1 == d2) ∧ (r1 == r2) ∧ (E[L1, Q1, P1] ≡ E[L2, Q2, P2]);
E_{d1 -> r1}[L1_, Q1_, P1_] E_{d2 -> r2}[L2_, Q2_, P2_] ^:=
  E[(d1 ∪ d2) -> (r1 ∪ r2)] @@ (E[L1, Q1, P1] E[L2, Q2, P2]);
E_{d -> r}_{-}[L_, Q_, P_]_{$k} := E_{d -> r} @@ E[L, Q, P]_{$k};
E_{-}[_{-}]_{i_{-}} := {ε}_{i};
```

"Define" code

```
SetAttributes[Define, HoldAll];
Define[def_, defs_] := (Define[def]; Define[defs]);
```



```

Define[op_is_ = ε_] :=
Module[{SD, ii, jj, kk, isp, nis, nisp, sis},
Block[{i, j, k},
ReleaseHold[Hold[
SD[opnisp, $k_Integer, Block[{i, j, k}, opisp, $k = ε;
opnis, $k]]];
SD[opisp, op{is}, $k]; SD[opsis_, op{sis}];
] /. {SD → SetDelayed,
isp → {is} /. {i → ii_, j → jj_, k → kk_},
nis → {is} /. {i → ii_, j → jj_, k → kk_},
nisp → {is} /. {i → ii_, j → jj_, k → kk_}
}]]]

```

The Fundamental Tensors

```

Define[am_i,j→k = E_{i,j}→{k}[(α_i + α_j) a_k, (e^{-γ α_j} ξ_i + ξ_j) x_k, 1]_{$k},
bm_i,j→k = E_{i,j}→{k}[(β_i + β_j) b_k, (η_i + η_j) y_k, e^{(ε - ε β_i - 1) η_j y_k}]_{$k}

```

```

Define[Ri_j =
E_{()→{i,j}}[ħ a_j b_i, ħ x_j y_i, e^{ħ} \sum_{k=2}^{k+1} \frac{(1 - e^{γ ε ħ})^k (ħ y_i x_j)^k}{k (1 - e^{k γ ε ħ})}]_{$k}

```

```

Define[Ri_j = E_{()→{i,j}}[-ħ a_j b_i, -ħ x_j y_i / B_i,
1 + If[$k == 0, 0, (R_{i,j}, $k-1)_{$k} [3] -
((R_{i,j}, 0)_{$k} R_{1,2} (R_{3,4}, $k-1)_{$k} // (bm_{i,1→i} am_{j,2→j} //
(bm_{i,3→i} am_{j,4→j})) [3] ]],

```

```

Pi_j = E_{i,j}→{}[β_i α_j / ħ, η_i ξ_j / ħ,
1 + If[$k == 0, 0, (P_{i,j}, $k-1)_{$k} [3] -
(R_{1,2} // ((P_{i,j}, 0)_{$k} (P_{i,2}, $k-1)_{$k})) [3] ]]]

```

```

Define[as_j = Ri_j ~ B_i ~ Pi_j,
as_i = E_{i}→{i}[-a_i α_i, -x_i A_i ξ_i,
1 + If[$k == 0, 0, (as_{i}, $k-1)_{$k} [3] -
((as_{i}, 0)_{$k} ~ B_i ~ as_{i} ~ B_i ~ (as_{i}, $k-1)_{$k}) [3] ]]]

```

```

Define[bs_i = Ri_1 ~ B_1 ~ as_1 ~ B_1 ~ Pi_1,
bs_i = Ri_1 ~ B_1 ~ as_1 ~ B_1 ~ Pi_1,
aΔ_i→j,k = (R_{i,j} R_{2,k}) // bm_{1,2→3} // P_{3,i},
bΔ_i→j,k = (R_{j,1} R_{k,2}) // am_{1,2→3} // P_{i,3}

```

```

Define[
dm_i,j→k =
(E_{i,j}→{i,j} [β_i b_i + α_j a_j, η_i y_i + ξ_j x_j, 1]
(aΔ_{i→1,2} // aΔ_{2→2,3} // as_3) (bΔ_{j→-1,-2} // bΔ_{-2→-2,-3}) //
(P_{-1,3} P_{-3,1} am_{2,j→k} bm_{i,-2→k}),
dSi = E_{i}→{1,2} [β_i b_1 + α_i a_2, η_i y_1 + ξ_i x_2, 1] // (bs_i as_2) //
dm_{2,1→i},
dΔ_i→j,k = (bΔ_{i→3,1} aΔ_{i→2,4}) // (dm_{3,4→k} dm_{1,2→j})

```

```

Define[Ci = E_{i}→{i} [0, 0, B_i^{1/2} e^{-ħ ε a_i/2}]_{$k},
Ci = E_{i}→{i} [0, 0, B_i^{-1/2} e^{ħ ε a_i/2}]_{$k},
Kink_i = (R_{1,3} C_2) // dm_{1,2→1} // dm_{1,3→i},
Kink_i = (R_{1,3} C_2) // dm_{1,2→1} // dm_{1,3→i}

```

```

Define[
b2t_i = E_{i}→{i} [α_i a_i - β_i t_i / γ, ξ_i x_i + η_i y_i, e^{ε β_i a_i / γ}]_{$k},
t2b_i = E_{i}→{i} [α_i a_i - t_i γ b_i, ξ_i x_i + η_i y_i, e^{ε t_i a_i}]_{$k}
Define[kRi_j = Ri_j // (b2t_i b2t_j) /. {t_i|j → t,
kRi_j = Ri_j // (b2t_i b2t_j) /. {t_i|j → t, Ti|j → T},
km_i,j→k = (t2b_i t2b_j) // dm_i,j→k //
b2t_k /. {t_k → t, T_k → T, t_i|j → 0},
kCi = Ci // b2t_i /. Ti → T, kCi = Ci // b2t_i /. Ti → T,
kKink_i = Kink_i // b2t_i /. {t_i → t, Ti → T},
kKink_i = Kink_i // b2t_i /. {t_i → t, Ti → T}

```

The Trefoil

```

$k = 2; Z = kR_{1,5} kR_{6,2} kR_{3,7} kC_4 kKink_8 kKink_9 kKink_{10};
Do[Z = Z ~ B_{1,r} ~ km_{1,r→1}, {r, 2, 10}];
Simplify /@ Z /. v_{-1} → v
E_{()→{1}} [0, 0, \frac{T}{1-T+T^2} + \frac{1}{(1-T+T^2)^3} T ħ (2 a (-1+T-T^3+T^4) +
T (-1+2T-3T^2+2T^3) γ - 2 (1+T^3) x y γ ħ) ε +
\frac{1}{2 (1-T+T^2)^5} T ħ^2 (4 a^2 (1-T+T^2)^2 (1+T-6T^2+T^3+T^4) +
4 a (1-T+T^2) γ (T (2-5T+8T^2-7T^3-2T^4+2T^5) -
2 (-1-2T+5T^2-4T^3+T^4+2T^5) x y ħ) +
γ^2 (T (1-2T+4T^2-2T^3+6T^5-11T^6+4T^7) +
4 (-1+2T+T^3+T^4+2T^6-T^7) x y ħ +
6 (1-T+T^2)^2 (1+3T+T^2) x^2 y^2 ħ^2)] ε^2 + O[ε]^3

```

diagram	n_k^r Alexander's ω^+	genus / ribbon	diagram	n_k^r Alexander's ω^+	genus / ribbon	diagram	n_k^r Alexander's ω^+	genus / ribbon
	Today's ρ_1^+	unknotting # / amphi?		Today's ρ_1^+	unknotting # / amphi?		Today's ρ_1^+	unknotting # / amphi?
	0_1^a 1	0 / ✓		3_1^a t - 1	1 / ✗		4_1^a 3 - t	1 / ✗
	0	0 / ✓		t	1 / ✗		0	1 / ✓
	5_1^a t^2 - t + 1	2 / ✗		5_2^a 2t - 3	1 / ✗		6_1^a 5 - 2t	1 / ✓
	$2t^3 + 3t$	2 / ✗		$5t - 4$	1 / ✗		t - 4	1 / ✗
	6_2^a -t^2 + 3t - 3	2 / ✗		6_3^a t^2 - 3t + 5	2 / ✗		7_1^a t^3 - t^2 + t - 1	3 / ✗
	$t^3 - 4t^2 + 4t - 4$	1 / ✗		0	1 / ✓		$3t^3 + 5t^3 + 6t$	3 / ✗
	7_2^a 3t - 5	1 / ✗		7_3^a 2t^2 - 3t + 3	2 / ✗		7_4^a 4t - 7	1 / ✗
	14t - 16	1 / ✗		$-9t^3 + 8t^2 - 16t + 12$	2 / ✗		$32 - 24t$	2 / ✗
	7_5^a 2t^2 - 4t + 5	2 / ✗		7_6^a -t^2 + 5t - 7	2 / ✗		7_7^a t^2 - 5t + 9	2 / ✗
	$9t^3 - 16t^2 + 29t - 28$	2 / ✗		$t^3 - 8t^2 + 19t - 20$	1 / ✗		8 - 3t	1 / ✗
	8_1^a 7 - 3t	1 / ✗		8_2^a -t^3 + 3t^2 - 3t + 3	3 / ✗		8_3^a 9 - 4t	1 / ✗
	5t - 16	1 / ✗		$2t^3 - 8t^4 + 10t^3 - 12t^2 + 13t - 12$	2 / ✗		0	2 / ✓
	8_4^a -2t^2 + 5t - 5	2 / ✗		8_5^a -t^3 + 3t^2 - 4t + 5	3 / ✗		8_6^a -2t^2 + 6t - 7	2 / ✗
	$3t^3 - 8t^2 + 6t - 4$	2 / ✗		$-2t^3 + 8t^4 - 13t^3 + 20t^2 - 22t + 24$	2 / ✗		$5t^3 - 20t^2 + 28t - 32$	2 / ✗
	8_7^a t^3 - 3t^2 + 5t - 5	3 / ✗		8_8^a 2t^2 - 6t + 9	2 / ✓		8_9^a -t^3 + 3t^2 - 5t + 7	3 / ✓
	$-t^5 + 4t^4 - 10t^3 + 12t^2 - 13t + 12$	1 / ✗		$-t^3 + 4t^2 - 12t + 16$	2 / ✗		0	1 / ✓
	8_{10}^a t^3 - 3t^2 + 6t - 7	3 / ✗		8_{11}^a -2t^2 + 7t - 9	2 / ✗		8_{12}^a t^2 - 7t + 13	2 / ✗
	$-t^5 + 4t^4 - 11t^3 + 16t^2 - 21t + 20$	2 / ✗		$5t^3 - 24t^2 + 39t - 44$	1 / ✗		0	2 / ✓
	8_{13}^a 2t^2 - 7t + 11	2 / ✗		8_{14}^a -2t^2 + 8t - 11	2 / ✗		8_{15}^a 3t^2 - 8t + 11	2 / ✗
	$-t^3 + 4t^2 - 14t + 20$	1 / ✗		$5t^3 - 28t^2 + 57t - 68$	1 / ✗		$21t^3 - 64t^2 + 120t - 140$	2 / ✗
	8_{16}^a t^3 - 4t^2 + 8t - 9	3 / ✗		8_{17}^a -t^3 + 4t^2 - 8t + 11	3 / ✗		8_{18}^a -t^3 + 5t^2 - 10t + 13	3 / ✗
	$t^5 - 6t^4 + 17t^3 - 28t^2 + 35t - 36$	2 / ✗		0	1 / ✓		0	2 / ✓
	8_{19}^a t^3 - t^2 + 1	3 / ✗		8_{20}^a t^2 - 2t + 3	2 / ✓		8_{21}^a -t^2 + 4t - 5	2 / ✗
	$-3t^5 - 4t^2 - 3t$	3 / ✗		4t - 4	1 / ✗		$t^3 - 8t^2 + 16t - 20$	1 / ✗