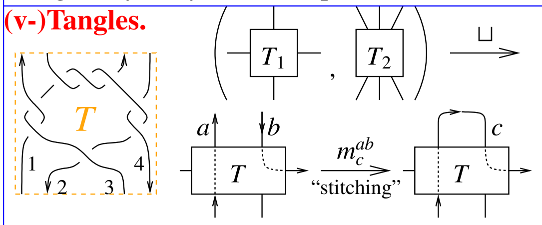
Dror Bar-Natan: Talks: Macquarie-191016: ωεβ:=http://drorbn.net/mac19/

Algebraic Knot Theory

Abstract. This will be a very "light" talk: I will explain why about 13 years ago, in order to have a say on some problems in knot theory, I've set out to find tangle invariants with some nice compositional properties. In other talks in Sydney (ωεβ/talks) I have explained / will explain how such invariants were found - though they are yet to be explored and utilized.

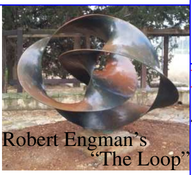
(v-)Tangles.



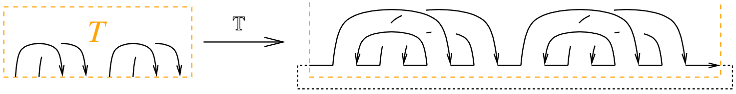
(meta-associativity: $m^{ab}_x /\!\!/ m^{xc}_y = m^{bc}_x /\!\!/ m^{ax}_y$) (tangles are generated by ⤯ and ⤰)

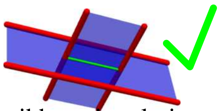
Strand doubling: $\xrightarrow{a}\ \xrightarrow{\Delta^a_{bc}}\ \begin{matrix}\xrightarrow{b}\\ \xrightarrow{c}\end{matrix}$ Strand reversal: $\xrightarrow{a}\ \xrightarrow{S_a}\ \xleftarrow{a}$

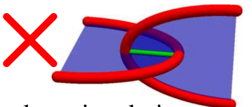
Genus. Every knot is the boundary of an orientable "Seifert Surface" (ωεβ/SS), and the least of their genera is the "genus" of the knot.


Robert Engman's "The Loop"

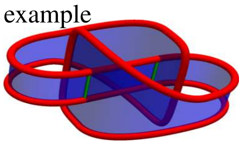
Claim. The knots of genus ≤ 2 are precisely the images of 4-component tangles via




a ribbon singularity


a clasp singularity

example

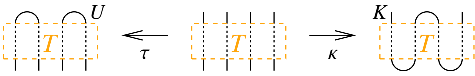


A Bit about Ribbon Knots. A "ribbon knot" is a knot that can be presented as the boundary of a disk that has "ribbon singularities", but no "clasp singularities". A "slice knot" is a knot in $S^3 = \partial B^4$ which is the boundary of a non-singular disk in B^4. Every ribbon knots is clearly slice, yet,

Conjecture. Some slice knots are not ribbon.

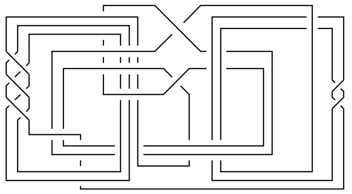
Fox-Milnor. The Alexander polynomial of a ribbon knot is always of the form $A(t) = f(t)f(1/t)$. (also for slice)

Theorem. K is ribbon iff it is κT for a tangle T for which τT is the untangle U.




Gompf, Scharlemann, Thompson [GST]

$$\begin{tikzcd}
 & U \in \mathcal{T}_n & & 1 \in \mathcal{A}_n & \\
\mathcal{T}_{2n} \arrow[ur, "\tau"] \arrow[dr, "\kappa"'] \arrow[rr, "z"] & & \mathcal{A}_{2n} \arrow[ur, "\tau"] \arrow[dr, "\kappa"'] & & \text{with } \mathcal{R} := \kappa(\tau^{-1}(1)) \\
 & \text{ribbon } K \in \mathcal{T}_1 & & z(K) \in \mathcal{R} \subseteq \mathcal{A}_1 &
\end{tikzcd}$$



Faster is better, leaner is meaner!

The Gold Standard is set by the "Γ-calculus" Alexander formulas [BNS, BN]. An S-component tangle T has $\Gamma(T) \in R_S \times M_{S\times S}(R_S) = \left\{ \begin{array}{c|c} \omega & S \\ \hline S & A \end{array} \right\}$ with $R_S := \mathbb{Z}(\{T_a : a \in S\})$:

$$\left({}_a\!\nearrow\!\!\!\nwarrow_b,\ {}_b\!\nwarrow\!\!\!\nearrow_a\right) \to \begin{array}{c|cc} 1 & a & b \\ \hline a & 1 & 1-T_a^{\pm1} \\ b & 0 & T_a^{\pm1} \end{array} \qquad T_1 \sqcup T_2 \to \begin{array}{c|cc} \omega_1\omega_2 & S_1 & S_2 \\ \hline S_1 & A_1 & 0 \\ S_2 & 0 & A_2 \end{array}$$

$$\begin{array}{c|ccc} \omega & a & b & S \\ \hline a & \alpha & \beta & \theta \\ b & \gamma & \delta & \epsilon \\ S & \phi & \psi & \Xi \end{array} \xrightarrow[T_a, T_b \to T_c]{m^{ab}_c} \left(\begin{array}{c|cc} (1-\beta)\omega & c & S \\ \hline c & \gamma + \frac{\alpha\delta}{1-\beta} & \epsilon + \frac{\delta\theta}{1-\beta} \\ S & \phi + \frac{\alpha\psi}{1-\beta} & \Xi + \frac{\psi\theta}{1-\beta} \end{array} \right)$$

For long knots, ω is Alexander, and that's the fastest Alexander algorithm I know! Dunfield: 1000-crossing fast.

Strand Doubling and Reversal.

$$\begin{array}{c|cc} \omega & a & S \\ \hline a & \alpha & \theta \\ S & \phi & \Xi \end{array} \xrightarrow[\substack{\mu := T_a - 1 \\ \nu := \alpha - \sigma_a \\ T_a \mapsto T_b T_c}]{q\Delta^a_{bc}} \left(\begin{array}{c|ccc} \omega & b & c & S \\ \hline b & (\sigma_a - \alpha T_a - \nu T_c)/\mu & (T_b - 1)T_c\nu/\mu & (T_b - 1)T_c\theta/\mu \\ c & (T_c - 1)\nu/\mu & (\alpha - \sigma_a T_a - \nu T_c)/\mu & (T_c - 1)\theta/\mu \\ S & \phi & \phi & \Xi \end{array} \right)$$

$$\Big\downarrow dS^a \ \ T_a \to T_a^{-1}$$

$$\left(\begin{array}{c|cc} \alpha\omega/\sigma_a & a & S \\ \hline a & 1/\alpha & \theta/\alpha \\ S & -\phi/\alpha & (\alpha\Xi - \phi\theta)/\alpha \end{array} \right)$$

Where σ assigns to every $a \in S$ a Laurent monomial σ_a in $\{t_b\}_{b\in S}$ subject to $\sigma\left({}_a\!\nearrow\!\!\!\nwarrow_b,\ {}_b\!\nwarrow\!\!\!\nearrow_a\right) = (a \to 1, b \to t_a^{\pm1})$, $\sigma(T_1 \sqcup T_2) = \sigma(T_1) \sqcup \sigma(T_2)$, and $\sigma /\!\!/ m^{ab}_c = (\sigma \setminus \{a,b\}) \cup (c \to \sigma_a\sigma_b)|_{t_a,t_b\to t_c}$.

Vo's Thesis [Vo]. A proof of the Fox-Milnor theorem for ribbon knots using this technology (and more).

Implementation key idea: ωεβ/AlexDemo

$$(\omega, A = (\alpha_{ab})) \leftrightarrow (\omega, \lambda = \textstyle\sum \alpha_{ab} t_a h_b)$$

```
Γ /: Γ[ω1_, λ1_] Γ[ω2_, λ2_] := Γ[ω1 * ω2, λ1 + λ2];
m_{a_ b_→c_}[Γ[ω_, λ_]] := Module[{α, β, γ, δ, θ, ϵ, ϕ, ψ, Ξ, μ},
   (α β θ)   (∂_{t_a,h_a}λ  ∂_{t_a,h_b}λ  ∂_{t_a}λ)
   (γ δ ϵ) = (∂_{t_b,h_a}λ  ∂_{t_b,h_b}λ  ∂_{t_b}λ) /. (t | h)_{a|b} → 0;
   (ϕ ψ Ξ)   (  ∂_{h_a}λ      ∂_{h_b}λ      λ   )
   Γ[(μ = 1 - β) ω, {t_c, 1}.(γ + α δ / μ   ϵ + δ θ / μ).{h_c, 1}]
                             (ϕ + α ψ / μ   Ξ + ψ θ / μ)
     /. {T_a → T_c, T_b → T_c} // ΓCollect];
Rp_{a_ b_} := Γ[1, {t_a, t_b}.(1  1 - T_a).{h_a, h_b}];
                              (0    T_a  )
Rm_{a_ b_} := Rp_{ab} /. T_a → 1 / T_a;
```

```
ΓCollect[Γ[ω_, λ_]] := Γ[Simplify[ω],
   Collect[λ, h_, Collect[#, t_, Factor] &]];
Format[Γ[ω_, λ_]] := Module[{S, M},
   S = Union@Cases[Γ[ω, λ], (h | t)_{a_} :→ a, ∞];
   M = Outer[Factor[∂_{h_{#1} t_{#2}} λ] &, S, S];
   M = Prepend[M, t_# & /@ S] // Transpose;
   M = Prepend[M, Prepend[h_# & /@ S, ω]];
   M // MatrixForm];
```

Meta-Associativity **Runs.**

```
ζ = Γ[ω, {t1, t2, t3, tS}.(α11 α12 α13 θ1).{h1, h2, h3, hS}];
                           (α21 α22 α23 θ2)
                           (α31 α32 α33 θ3)
                           (ϕ1  ϕ2  ϕ3  Ξ )
(ζ // m12→1 // m13→1) == (ζ // m23→2 // m12→1)
```

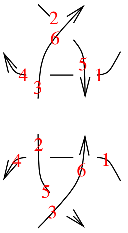
```
True
```

R3

```
{Rm51 Rm62 Rp34 // m14→1 // m25→2 // m36→3,
  Rp61 Rm24 Rm35 // m14→1 // m25→2 // m36→3}
```

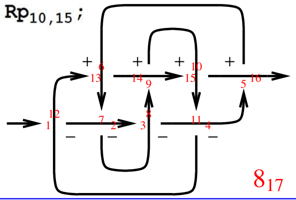
```
  ( 1     h1        h2        h3 )     ( 1     h1        h2        h3 )
  ( t1    T3/T2     0         0  )     ( t1    T3/T2     0         0  )
{ ( t2    (-1+T2)/T2  1/T3    0  )  ,  ( t2    (-1+T2)/T2  1/T3    0  ) }
  ( t3  -(-1+T3)/T2 (-1+T3)/T3 1 )     ( t3  -(-1+T3)/T2 (-1+T3)/T3 1 )
```

. . . divide and conquer!



```
z = Rm12,1 Rm27 Rm83 Rm4,11 Rp16,5 Rp6,13 Rp14,9 Rp10,15;
Do[z = z // m1k→1, {k, 2, 16}];
z
```

```
( 11 - 1/T1^3 + 4/T1^2 - 8/T1 - 8 T1 + 4 T1^2 - T1^3   h1 )
(                     t1                               1  )
```



Fact. Γ is better viewed as an invariant of a certain class of 2D knotted objects in $\mathbb{R}^4$ [BND, BN].

Fact. Γ is the "0-loop" part of an invariant that generalizes to "n-loops" (1D tangles only, see further talks and future publications with van der Veen).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Speculation. Stepping stones to categorification?

Ask me about geography vs. identity!

References.

[BN] D. Bar-Natan, *Balloons and Hoops and their Universal Finite Type Invariant, BF Theory, and an Ultimate Alexander Invariant,* ωεβ/KBH, arXiv:1308.1721.

[BND] D. Bar-Natan and Z. Dancso, *Finite Type Invariants of W-Knotted Objects I: w-Knots and the Alexander Polynomial,* Alg. and Geom. Top. **16-2** (2016) 1063–1133, arXiv:1405.1956, ωεβ/WKO1.

[BNS] D. Bar-Natan and S. Selmani, *Meta-Monoids, Meta-Bicrossed Products, and the Alexander Polynomial,* J. of Knot Theory and its Ramifications **22-10** (2013), arXiv:1302.5689.

[GST] R. E. Gompf, M. Scharlemann, and A. Thompson, *Fibered Knots and Potential Counterexamples to the Property 2R and Slice-Ribbon Conjectures,* Geom. and Top. **14** (2010) 2305–2347, arXiv:1103.1601.

[Vo] H. Vo, *Alexander Invariants of Tangles via Expansions,* University of Toronto Ph.D. thesis, ωεβ/Vo.

> "God created the knots, all else in topology is the work of mortals."
>
> Leopold Kronecker (modified)

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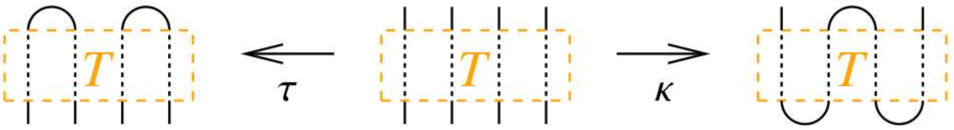
Proof of the Tangle Characterization of Ribbon Knots

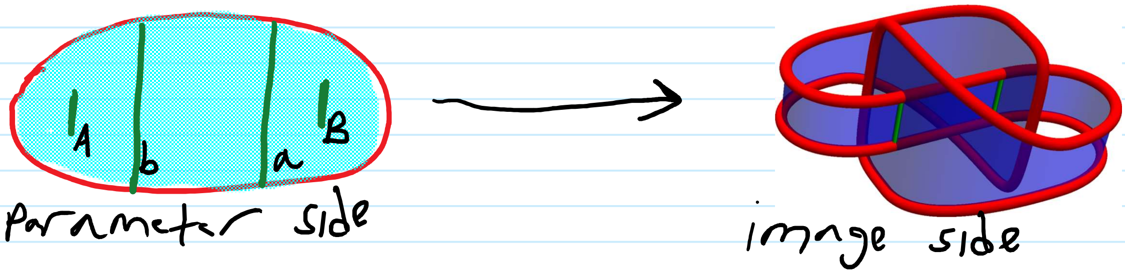


(tangle has 2n strands, here n=2)

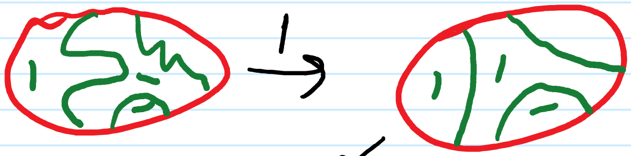
Theorem. A knot *K* is ribbon iff there exists a tangle *T* whose τ closure is the untangle and whose κ closure is *K*.

Proof. The backward ⟸ implication is easy:

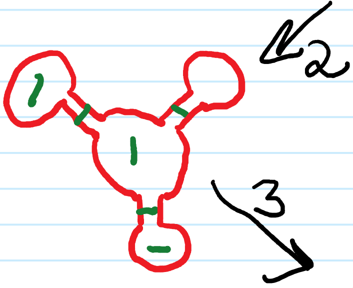
For the forward implication, follow the following 5 steps:



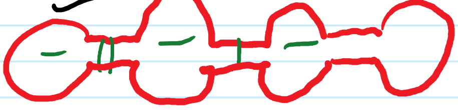
Step I: In-situ cosmetics.
At end: D is a tree of chord-and-arc polygons.

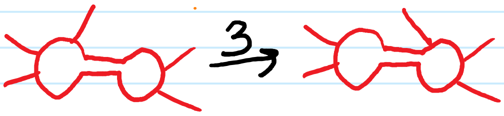


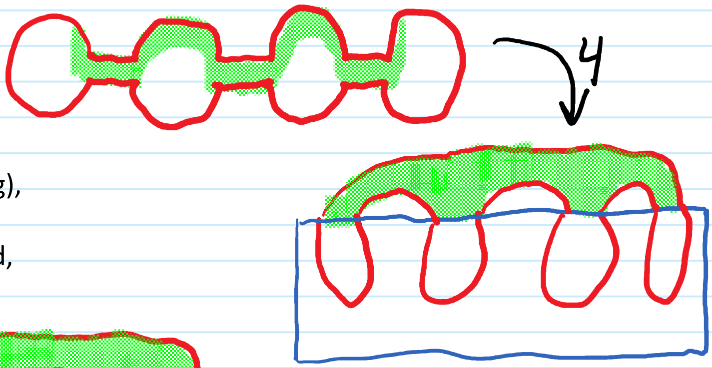
Step 2: Near-situ cosmetics.
At end: D is tree-band-sum of n unknotted disks.



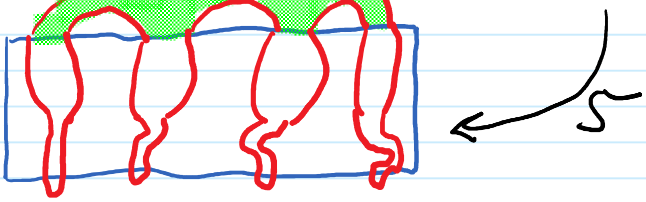
Step 3: Slides.
At end: D is a linear-band-sum of n unknotted disks.







Step 4: Exposure!
The green domain is contractible - so it can be shrank, moved at will (with the blue membrane following along), and expanded back again.
At end: D has (n-1) exposed bridges which when turned, make D a union of n unknotted disks.



Step 5: Pulling bottom handles avoiding the obstacles.
At end: Theorem is proven.