


Solve the PgUp/PgDn HandoutBrowser bug! 

Dror Bar-Natan: Talks: MIT-1612: (thanks for accepting my invitation!) orβ=http://drorbn.net/MIT-1612/ 

### A Poly-Time Knot Polynomial Via Solvable Approximation

**Abstract.** Rozansky [Ro2] and Overbay [Ov] described a **spectacular** knot polynomial that failed to attract the attention it deserved as the first poly-time-computable knot polynomial since Alexander's [Al, 1928] and (in my opinion) as the second most likely knot polynomial (after Alexander's) to carry topological information. With Roland van der Veen, I will explain how to compute the Rozansky polynomial using some new commutator-calculus techniques and a Lie algebra  $\mathfrak{g}$ , which is at the same time solvable and an approximation of the simple Lie algebra  $sl_2$ .

**Theorem** ([BNG], conjectured [MM], elucidated [Ro1]). Let  $J_d(K)$  be the coloured Jones polynomial of  $K$ , in the  $d$ -dimensional representation of  $sl_2$ . Writing

$$\left. \frac{(q^{1/2} - q^{-1/2})J_d(K)}{q^{d/2} - q^{-d/2}} \right|_{q=e^h} = \sum_{j,m \geq 0} a_{jm}(K) d^j h^m,$$

"below diagonal" coefficients vanish,  $a_{jm}(K) = 0$  if  $j > m$ , and "on diagonal" coefficients give the inverse of the Alexander polynomial:  $(\sum_{m=0}^{\infty} a_{mm}(K) h^m) \cdot A(K)(e^h) = 1$ .

"Above diagonal" we have **Rozansky's Theorem** [Ro3, (1.2)]:

$$J_d(K)(q) = \frac{q^d - q^{-d}}{(q - q^{-1})A(K)(q^d)} \left( 1 + \sum_{k=1}^{\infty} \frac{(q-1)^k R_k(K)(q^d)}{A^{2k}(K)(q^d)} \right).$$

**The Gold Standard** is set by the "T-calculus" Alexander formulas [BNS, BN1]. An  $S$ -component tangle  $T$  has  $\Gamma(T) \in R_S \times M_{S \times S}(R_S) = \left\{ \begin{matrix} \omega & S \\ S & A \end{matrix} \right\}$  with  $R_S := \mathbb{Z}\langle t_a : a \in S \rangle$ :

$$\left( \begin{matrix} a & b \\ b & a \end{matrix} \right) \rightarrow \frac{1}{a} \begin{pmatrix} a & b \\ 0 & 1 - t_a^{-1} \end{pmatrix} \quad T_1 \sqcup T_2 \rightarrow \begin{pmatrix} \omega_1 \omega_2 & S_1 & S_2 \\ S_1 & A_1 & 0 \\ S_2 & 0 & A_2 \end{pmatrix}$$

$$\begin{pmatrix} \omega & a & b & S \\ a & \alpha & \beta & \theta \\ b & \gamma & \delta & \epsilon \\ S & \phi & \psi & \Xi \end{pmatrix} \xrightarrow{m_c^{ab}} \begin{pmatrix} (1-\beta)\omega & c & S \\ c & \gamma + \frac{\alpha\delta}{1-\beta} & \epsilon + \frac{\theta\delta}{1-\beta} \\ S & \phi + \frac{\alpha\psi}{1-\beta} & \Xi + \frac{\psi\theta}{1-\beta} \end{pmatrix}$$

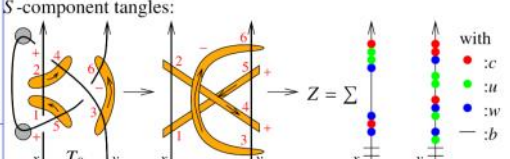
(Roland: "add to  $A$  the product of column  $b$  and row  $a$ , divide by  $(1 - A_{ab})$ , delete column  $b$  and row  $a$ ".)

For long knots,  $\omega$  is Alexander, and that's the fastest Alexander algorithm I know! Dunfield: 1000-crossing fast. (There are also formulas for strand doubling and strand reversal).

**Theorem** [EK, Ha, En, Se]. There is a "homomorphic expansion"  $Z: \{S\text{-component } (v/b)\text{-tangles}\} \rightarrow \mathcal{A}_S^v :=$

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**Algebras and Invariants.** Given any unital algebra  $A$  (even better if  $A$  is Hopf; typically,  $A \sim \hat{\mathcal{U}}(\mathfrak{g})$ ), appropriate orange  $R \in A \otimes A$ , and appropriate cuaps  $\epsilon \in A$ , get an  $A^{\otimes S}$ -valued invariant of pure  $S$ -component tangles:





**Good News.** In theory, enough to know  $R$ , the cuaps, and stitching/multiplication  $m_k^j: A_j \otimes A_j \rightarrow A_k$ .


**Textbook Solution.** Use representation theory ... works, slowly.

**Today's Solution** (with van der Veen). For some specific  $\mathfrak{g}$ 's, work in a space of "formulas of a specific type" for elements of  $\hat{\mathcal{U}}(\mathfrak{g})^{\otimes S}$ :

$$\{ \text{ordered perturbed Gaussian formulas} \} \rightarrow \hat{\mathcal{U}}(\mathfrak{g})^{\otimes S}$$

van der Veen 

Leopold Kronecker (modified) www.katlas.org 

*bring to front?* 

**1-Smidgen  $sl_2$**  Let  $\mathfrak{g}_1$  be the 4-dimensional Lie algebra  $\mathfrak{g}_1 = \langle b, c, u, w \rangle$  over the ring  $R = \mathbb{Q}[\epsilon]/(\epsilon^2 = 0)$ , with  $b$  central and with  $[w, c] = w, [c, u] = u$ , and  $[u, w] = b - 2\epsilon c$ , with CYBE  $r_{ij} = (b_i - \epsilon c_i)c_j + u_i w_j$  in  $\mathcal{U}(\mathfrak{g}_1)^{\otimes(i,j)}$ . Over  $\mathbb{Q}$ ,  $\mathfrak{g}_1$  is a solvable approximation of  $sl_2$ :  $\mathfrak{g}_1 \supset \langle b, u, w, \epsilon b, \epsilon c, \epsilon u, \epsilon w \rangle \supset \langle b, \epsilon b, \epsilon c, \epsilon u, \epsilon w \rangle \supset 0$ . (note:  $\deg(b, c, u, w, \epsilon) = (1, 0, 1, 0, 1)$ )

**0-Smidgen  $sl_2 \oplus$**  Let  $\mathfrak{g}_0$  be  $\mathfrak{g}_1$  at  $\epsilon = 0$ , or  $\mathbb{Q}\langle b, c, u, w \rangle / ([b, \cdot] = 0, [c, u] = u, [c, w] = -w, [u, w] = b$  with  $r_{ij} = b_i c_j + u_i w_j$ . It is  $\mathfrak{b}^+ \rtimes \mathfrak{b}$  where  $\mathfrak{b}$  is the 2D Lie algebra  $\mathbb{Q}\langle c, w \rangle$  and  $(b, u)$  is the dual basis of  $(c, w)$ . For topology, it is more valuable than  $\mathfrak{g}_1 / sl_2$ , but topology already got by other means almost everything  $\mathfrak{g}_0$  gives.

**How did these arise?**  $sl_2 = \mathfrak{b}^+ \oplus \mathfrak{b}^- / \mathfrak{h} =: sl_2^+ / \mathfrak{h}$ , where  $\mathfrak{b}^+ = \langle c, w \rangle / [w, c] = w$  is a Lie bialgebra with  $\delta: \mathfrak{b}^+ \rightarrow \mathfrak{b}^+ \otimes \mathfrak{b}^+$  by  $\delta: (c, w) \mapsto (0, c \wedge w)$ . Going back,  $sl_2^+ = \mathcal{D}(\mathfrak{b}^+) = (\mathfrak{b}^+)^* \oplus \mathfrak{b}^+ = \langle b, u, c, w \rangle / \dots$ . **Idea.** Replace  $\delta \rightarrow \epsilon \delta$  over  $\mathbb{Q}[\epsilon]/(\epsilon^{k+1} = 0)$ . At  $k = 0$ , get  $\mathfrak{g}_0$ . At  $k = 1$ , get  $[w, c] = w, [w, b'] = -\epsilon w, [c, u] = u, [b', u] = -\epsilon u, [b', c] = 0$ , and  $[u, w] = b' - \epsilon c$ . Now note that  $b' + \epsilon c$  is central, so switch to  $b := b' + \epsilon c$ . This is  $\mathfrak{g}_1$ .

**Ordering Symbols.**  $\mathcal{O}$  (*poly* | *specs*) plants the variables of *poly* in  $\mathcal{S}(\mathfrak{g})$  on several tensor copies of  $\mathcal{U}(\mathfrak{g})$  according to *specs*. E.g.,  $\mathcal{O}(c_1^3 u_1 c_2 e^{u_3} w_3^2 | x: w_3 c_1, y: u_1 u_3 c_2) = w^9 c^3 \otimes u e^u c \in \mathcal{U}(\mathfrak{g}_1) \otimes \mathcal{U}(\mathfrak{g}_1)$ . This enables the description of elements of  $\mathcal{U}(\mathfrak{g})^{\otimes S}$  using commutative polynomials / power series.

**0-Smidgen Invariants.**  $r = Id \in \mathfrak{b}^- \otimes \mathfrak{b}^+$  solves the CYBE  $[r_{12}, r_{13}] + [r_{12}, r_{23}] + [r_{13}, r_{23}] = 0$  in  $\mathcal{U}(\mathfrak{g}_0)^{\otimes 3}$  and, by luck,

$$\begin{array}{c} \begin{array}{|c|} \hline \times \\ \hline \end{array} \begin{array}{|c|} \hline \rightarrow \\ \hline \end{array} \begin{array}{|c|} \hline \uparrow \\ \hline \end{array} \begin{array}{|c|} \hline \downarrow \\ \hline \end{array} \\ \hline \end{array} = R_{ij} = e^{c_j} u_i = e^{b_i c_j + u_i w_j} \in \mathcal{U}(\mathfrak{g}_0) \otimes \mathfrak{g}_0$$

solves YB/R3.

**Lemma.**  $R_{ij} = e^{b_i c_j + u_i w_j} = \mathcal{O}\left(\exp\left(b_i c_j + \frac{e^{b_i-1}}{b_i} u_i w_j\right) | i: u_i, j: c_j w_j\right)$

**Example.**  $Z(T_0) = \sum_{m,n} \frac{b_i^{m-1} (e^{b_i-1})^n}{m! n!} u^m c^n e^{m u w^n}$

$$\mathcal{O}\left(\exp\left(b_5 c_1 + \frac{e^{b_5-1}}{b_5} u_5 w_1 + b_2 c_4 + \frac{e^{b_2-1}}{b_2} u_2 w_4 - b_3 c_6 + \frac{e^{b_3-1}}{b_3} u_3 w_6\right) | \begin{array}{l} \text{"ucw form"} \\ x: c_1 w_1 u_2, y: u_3 c_4 w_4 u_5 c_6 w_6 \end{array}\right) = \mathcal{O}\left(\zeta | x: u_i c_j w_k, y: u_i c_j w_k\right)$$

**Goal.** Write  $\zeta$  as a Gaussian:  $\omega e^{L+Q}$  where  $L$  bilinear in  $b_i$  and  $c_j$  with integer coefficients,  $Q$  a balanced quadratic in  $u_i$  and  $w_j$  with coefficients in  $R_S := \mathbb{Q}(b_i, e^{b_i})$ , and  $\omega \in R_S$ .

- The Big  $\mathfrak{g}_0$  Lemma.** Under  $[c, u] = u, [c, w] = -w$ , and  $[u, w] = b$ :
- $N^{uc} := \mathcal{O}(e^{\gamma c + \beta u} | uc) \equiv \mathcal{O}(e^{\gamma c + \epsilon^{-\gamma} \beta u} | cu)$  (means  $e^{\beta u} e^{\gamma c} = e^{\gamma c} e^{-\gamma \beta u}$ )
  - $N^{wc} := \mathcal{O}(e^{\gamma c + \alpha w} | wc) \equiv \mathcal{O}(e^{\gamma c + e^{\alpha} w} | cw)$  ... in the  $(\alpha x + b)$  group
  - $\mathcal{O}(e^{\alpha w + \beta u} | wu) = \mathcal{O}(e^{-b\alpha\beta + \alpha w + \beta u} | uw)$  (the Weyl relations)
  - $\mathcal{O}(e^{\delta u w} | wu) e^{\beta u} = e^{\gamma \beta u} \mathcal{O}(e^{\delta u w} | wu)$ , with  $\gamma = (1 + b\delta)^{-1}$
  - (a. expand and crunch. b. use  $w = b\hat{x}, u = \partial_x$ . c. use "scatter and glow".)
  - $\mathcal{O}(e^{\delta u w} | wu) = \mathcal{O}(v e^{\gamma \delta u w} | uw)$  (same techniques)
  - $N^{wu} := \mathcal{O}(e^{\beta u + \alpha w + \delta u w} | wu) \equiv \mathcal{O}(v e^{-b\alpha\beta + \alpha w + \beta u + \gamma \delta u w} | uw)$
  - $N_k^{c_j} := \mathcal{O}(\zeta | c_j c_j) \equiv \mathcal{O}(\zeta / (c_i, c_j \rightarrow c_k) | c_k)$

**Sneaky.**  $\alpha$  may contain (other)  $u$ 's,  $\beta$  may contain (other)  $w$ 's.

**Strand Stitching.**  $m_k^{ij}$  is defined as the composition

$$u_i c_j \overline{w_i} \overline{u_j} c_j w_j \xrightarrow{N_k^{w_j}} u_i c_j \overline{u_x} \overline{w_x} c_j w_j \xrightarrow{N_k^{u_x} // N_k^{w_j}} \overline{u_i} \overline{u_x} c_x c_x \overline{w_x} w_j \xrightarrow{i, j, x \rightarrow k} u_k c_k w_k$$


**1-Smidgen Invariants.** Much is the same: **The Big  $\mathfrak{g}_1$  Lemma.** Parts 1 and 6 are the same, yet  $\mathcal{O}(e^{\alpha w + \beta u + \delta u w} | wu) = \mathcal{O}(v(1 + \epsilon v \Lambda) e^{v(-b\alpha\beta + \alpha w + \beta u + \delta u w)} | cuw)$ . Here  $\Lambda$  is for  $\Lambda\delta\gamma\alpha\delta$ , "a principle of order and knowledge", a balanced quartic in  $\alpha, \beta, c, u$ , and  $w$ :

$$\begin{aligned} \Lambda = & -bv(v^2\alpha^2\beta^2 + 4\delta v\alpha\beta + 2\delta^2)/2 - \delta v^3(3b\delta + 2)\beta^2 u^2/2 \\ & - b\delta^4 v^3 u^2 w^2/2 - \delta^2 v^3(2b\delta + 1)\beta u^2 w \\ & - v^2(2b\delta + 1)(v\alpha\beta + 2\delta)\beta u - 2b\delta^2 v^2(v\alpha\beta + \delta)uw \\ & + \delta v^3(b\delta + 2)\alpha^2 w^2/2 + 2(v\alpha\beta + \delta)c + 2\delta v\beta c u + 2\delta^2 v c u w \\ & + 2\delta v\alpha c w + \delta^2 v^3 a u w^2 + v^2(v\alpha\beta + 2\delta)aw. \end{aligned}$$

**Proof.** A lengthy computation. **Problem.** We now need to normal-order perturbed Gaussians! **Solution.** Borrow some tactics from QFT:

$$\mathcal{O}(\epsilon P(c, u) e^{\gamma c + \beta u} | uc) = \mathcal{O}(\epsilon P(\partial_\gamma, \partial_\beta) e^{\gamma c + \beta u} | uc) = \mathcal{O}(\epsilon P(\partial_\gamma, \partial_\beta) e^{\gamma c + \epsilon^{-\gamma} \beta u} | cu)$$

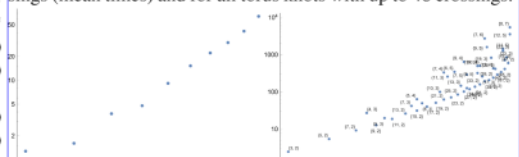
and likewise  $\mathcal{O}(\epsilon P(u, w) e^{\alpha w + \beta u + \delta u w} | wu) = \mathcal{O}(\epsilon P(\partial_\beta, \partial_\alpha) v e^{v(-b\alpha\beta + \alpha w + \beta u + \delta u w)} | cuw)$

**Note.** Strand stitching requires a tiny extra step. **Finally,** the values of the generators  $\zeta, \bar{\zeta}, \vec{n}$ , and  $\vec{u}$  are set by solving many equations, non-uniquely.

**Pragmatic Simplifications.** Get rid of  $\zeta = (e^b - 1)/b$  factors by rescaling  $u \rightarrow \bar{u} = \zeta u$ . Complement this with  $\beta \rightarrow \bar{\beta} = \zeta^{-1} \beta$ ,  $\delta \rightarrow \bar{\delta} = \zeta^{-1} \delta$ ,  $\epsilon \rightarrow \bar{\epsilon} = \zeta^{-1} \epsilon$ . Simplify further by naming  $e^b \rightarrow t$ ; e.g.,  $v \rightarrow \bar{v} = (1 + (t-1)\delta)^{-1}$ . Get confused by renaming  $(\bar{u}, \bar{\beta}, \bar{\delta}, \bar{v}) \rightarrow (u, \beta, \delta, v)$ , and more confused by working with  $\mu = v^{-1}$  and  $\mathbb{E}(\omega, L, Q, P) := \omega^{-1}(1 + \epsilon \omega^{-d} P) e^{L + \omega^{-1} Q}$ , where  $\omega \in R := \mathbb{Q}(t_k)$ ,  $L = \sum l_{ij} b_i c_j$  with  $l_{ij} \in \mathbb{Z}$ ,  $Q = \sum q_{ij} u_i w_j$  with  $q_{ij} \in R$ , and  $P$  is a balanced quartic polynomial in  $c_i, u_i$ , and  $w_j$  with coefficients in  $R$ . Magically, all coefficients are now Laurent polynomials in the  $t_k$ 's.

**Rough complexity estimate,** after  $t_k \rightarrow t$ :  $n$ : xing number;  $w$ : width, maybe  $\frac{B}{A} \sum_{d=0}^4 \frac{w^{4-d} w^d n^2}{E F G} = n^3 w^4 \in [n^5, n^7]$ ;  $\sim \sqrt{n}$ .  $A$ : go over stitchings in order.  $B$ : multiplication ops per  $N^{u_i w_j}$ .  $d$ : deg of  $u_i, w_j$  in  $P$ .  $E$ : #terms of deg  $d$  in  $P$ .  $F$ : ops per term.  $G$ : cost per polynomial multiplication op.

**Experimental Analysis** ( $\omega\epsilon\beta/\text{Exp}$ ). Log-log plot of computation time (sec) vs. crossing number, for all knots with up to 12 crossings (mean times) and for all torus knots with up to 48 crossings:



**Conjecture** (checked on the same collections). Given a knot  $K$  with Alexander polynomial  $A$ , there is a polynomial  $\rho_1$  such that  $P = A^2 \frac{(t-1)^3 \rho_1 + t^2(2vw + (1-t)(1-2c))AA'}{(1-t)t}$ .

Furthermore,  $A$  and  $\rho_1$  are symmetric under  $t \rightarrow t^{-1}$ , so let  $A^+$  and  $\rho_1^+$  be their "positive parts", so e.g.,  $\rho_1(t) = \rho_1^+(t) + \rho_1^+(t^{-1}) - \rho_1^+(0)$ . **Power.** On the 250 knots with at most 10 crossings, the pair  $(A, \rho_1)$  attains 250 distinct values, while (Khovanov, HOMFLY-PT) attains only 249 distinct values. To 11 crossings the numbers are (802, 788, 772) and to 12 they are (2978, 2883, 2786).

improved

improved

ncu

expm?

ucw

changed

include a ref?

ucw

revised

**Genus.** Up to 12 crossings, always  $\text{deg } \rho_1^+ \leq 2g - 1$ , where  $g$  is the 3-genus of  $K$  (equality for 2530 knots). This gives a lower bound on  $g$  in terms of  $\rho_1^+$  (conjectural, but undoubtedly true). This bound is often weaker than the Alexander bound, yet for 10 of the 12-crossing Alexander failures it does give the right answer.

**Demo Programs for 0-Co.**

ωεβ/Demo

```
R0_{i,j} := E[b_i c_j + b_i^{-1} (e^{b_i} - 1) u_i w_j];
R0_{i,j} := E[-b_i c_j + b_i^{-1} (e^{-b_i} - 1) u_i w_j];
```

The R-matrices

**Utilities**

```
CF[ω_ E[Q_]] := Simplify[ω E[Simplify[Q]];
E /: E[Q1] E[Q2] := CF@E[Q1 + Q2];
ω1_ E[Q1] := ω2_ E[Q2] := Simplify[ω1 = ω2 ∧ Q1 = Q2];
```

**Normal Ordering Operators**

```
N_{(x:w)u_i c_j -> h_ [ω_ E[Q_]] := CF[
ω E[e^y α x_h + γ c_h + (Q / c_j | x_i → θ)] / . {y → ∂_{c_j} Q, α → ∂_{x_i} Q};
N_{u_i c_j -> h_ [ω_ E[Q_]] := CF[
v ω E[-b_i v α β + v β u_h + v δ u_h + v α w_h + (Q / w_i | u_j → θ)] / .
v → (1 + b_i δ)^{-1} / .
{α → ∂_{u_i} Q / u_j → θ, β → ∂_{u_j} Q / w_i → θ, δ → ∂_{w_i} u_j Q};
```

Stitching

```
m_{i,j} -> h_ [Z_ := Module[{x, z},
CF[Z // N_{u_i v_j -> x} // N_{c_i u_k -> x} // N_{w_k c_j -> x} / . Z_{-i|j|x} → z_h]]
```

Some calculations for T0

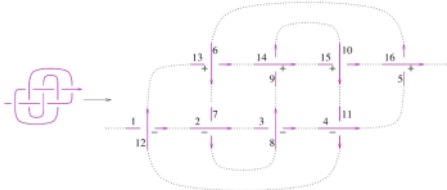
```
T0 = R0_{5,1} R0_{2,4} R0_{3,6}
E[b_5 c_1 + b_2 c_4 - b_3 c_6 + (-1 e^{b_5}) u_5 w_1 + (-1 e^{b_2}) u_2 w_4 + (-1 e^{-b_3}) u_3 w_6]
T0 // m_{1,2+1} // m_{3,4+3} // m_{5,3+3} // m_{3,6+3}
1 - (-1 e^{b_1}) (-1 e^{b_3}) E[b_3 c_1 + b_1 c_3 - b_3 c_3 +
e^{b_3} (-1 e^{b_2}) (-1 e^{b_3}) u_1 w_1 - e^{b_1} (-1 e^{b_3}) u_3 w_3 -
(-1 e^{b_1} e^{b_3} e^{b_1} b_3) b_1 - (-1 (-1 e^{b_1}) (-1 e^{b_3})) b_3 -
e^{-b_3} (-1 e^{b_3}) u_3 w_3 - e^{-b_3} (-1 e^{b_1}) (-e^{b_3} b_3 u_1 + e^{b_1} (-1 e^{b_3}) b_1 u_3) w_3]
```

Verifying meta-associativity

```
Q0 = E[Sum[f_i c_i, {i, 3}] + Sum[f_{i,j} u_i w_j, {i, 3}, {j, 3}]]
E[c_1 f_1 + c_2 f_2 + c_3 f_3 + u_1 w_1 f_{1,1} + u_1 w_2 f_{1,2} + u_1 w_3 f_{1,3} + u_2 w_1 f_{2,1} +
u_2 w_2 f_{2,2} + u_2 w_3 f_{2,3} + u_3 w_1 f_{3,1} + u_3 w_2 f_{3,2} + u_3 w_3 f_{3,3}]
(Q0 // m_{1,2+1} // m_{3,3+1}) = (Q0 // m_{2,3+2} // m_{1,2+1})
True
```

Testing R3

```
t1 = R0_{8,12} R0_{3,4} R0_{5,6} // m_{3,5-x} // m_{1,6-y} // m_{2,4-z}
E[b_x c_y + b_x c_z + b_y c_z + e^{b_x} (-1 e^{b_y}) u_y w_z + (-1 e^{b_x}) u_x (w_y w_z)]
t1 = (R0_{8,12} R0_{3,4} R0_{5,6} // m_{3,3-x} // m_{2,5-y} // m_{4,6-z})
True
```



817

```
z1 = R0_{12,1} R0_{2,7} R0_{8,3} R0_{4,11} R0_{16,5} R0_{6,13} R0_{14,9} R0_{10,15};
Do[z1 = (z1 // m_{1,n-1}) / . b_ -> b, {n, 2, 16}];
{CF@z1, KnotData[{8, 17}, "AlexanderPolynomial"]}[t]
[-(e^{3 b} e^{2 b} - 11 e^{3 b} e^{4 b} e^{5 b} e^{6 b}) / (1 - 4 e^{b} e^{2 b} - 11 e^{3 b} e^{4 b} e^{5 b} e^{6 b}), 11 - 1/3 + 4/2 - 8/t + 4 t^2 - t^3]
```

**Demo Programs for 1-Co.**

ωεβ/Demo

```
Δ[h_] := ((t_h - 1) (2 (α β + δ μ)^2 - α^2 β^2) - 4 v_h c_h w_h δ^2 μ^2 -
δ (1 + μ) (w_h^2 α^2 + v_h^2 β^2) - v_h^2 w_h^2 δ^3 (1 + 3 μ) -
2 (α β + 2 δ μ + v_h w_h δ^2 (1 + 2 μ) + 2 c_h δ μ^2) (w_h α + v_h β) -
4 (c_h μ^2 + v_h w_h δ (1 + μ) (α β + δ μ) (1 + t_h) / 4; The Δ_ωγ
```

```
R_{i,j}^+ := E[1, Log[t_i] c_j, v_i w_j, v_i c_i w_j + c_i c_j + v_i^2 w_j^2] / 4;
R_{i,j}^- := E[1, -Log[t_i] c_j, -t_i^{-1} v_i w_j,
-c_i c_j + t_i^{-1} v_i c_j w_j - t_i^{-2} v_i^2 w_j^2 / 4];
(ur_{i,j} := E[t_i^{1/2}, θ, θ, c_i t_i^{-2}]; nr_{i,j} := E[t_i^{1/2}, θ, θ, -c_i t_i^{-2}];
```

Differential Polynomials

```
DP_{x→0, y→0, z→0} [f_] := (* means P[∂_α, ∂_β] [f] *)
Total[CoefficientRules[P, {x, y}]] / .
({m_-, n_} → c_ -> c D[f, {α, m}, {β, n}]]
CF[E_ := Expand /@ Together /@ E;
E /: E[ω1_ L1_ Q1_ P1_] E[ω2_ L2_ Q2_ P2_] :=
CF@E[ω1 ω2, L1 + L2, ω2 Q1 + ω1 Q2, ω2^2 P1 + ω1^2 P2];
```

Normal Ordering Operators

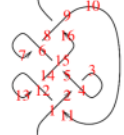
```
N_{c_j (x:w)u_i c_j -> h_ [E[ω_ L_ Q_ P_]] := With[{q = e^y β x_h + γ c_h}, CF[
E[ω y c_h + (L / c_j → θ), ω e^y β x_h + (Q / x_i → θ),
e^{-q} DP_{c_j -> 0, y_i -> 0} [P] [e^q]] / . {y → ∂_{c_j} L, β → ω^{-1} ∂_{y_i} Q}]];
N_{u_i v_j -> h_ [E[ω_ L_ Q_ P_]] :=
With[{q = ((1 - t_h) α β + β v_h + δ v_h w_h + α w_h) / μ}, CF[
E[μ ω y L, μ ω q + μ (Q / w_i | v_j → θ),
μ^4 e^{-q} DP_{w_i -> 0, v_j -> 0} [P] [e^q] + ω^4 Δ[h]] / . μ → 1 + (t_h - 1) δ / .
{α → ω^{-1} (∂_{u_i} Q / v_j → θ), β → ω^{-1} (∂_{v_j} Q / w_i → θ),
δ → ω^{-2} ∂_{w_i} v_j Q}]];
```

Stitching

```
m_{i,j} -> h_ [Z_ := Module[{x, z},
CF[Z // N_{u_i v_j -> x} // N_{c_i u_k -> x} // N_{w_k c_j -> x} / . Z_{-i|j|x} → z_h]]
```

```
z2 = R_{1,11}^+ R_{4,2}^+ nr_{3,15}^+ R_{6,8}^+ ur_{7,16}^+ nr_{10,12,14}^+ ur_{13}^+;
(Do[z2 = z2 // m_{1,k-1}, {k, 2, 16}];
z2 = z2 / . a_{-1} → a)
```

The 0-Framed Trefoil



```
E[-1 + 1/t + t, θ, θ,
16 + 2/c - 1/3 - 6/c^2 + 4/t^2 + 10/c^2 - 10/t - 8/c - 18t + 8ct +
14t^2 - 10ct^2 - 7t^3 + 6ct^3 + 2t^4 - 2ct^4 + 2vw -
2xw + 4yw - 6xw/t^2 + 2xw/t - 6tvw + 4t^2vw - 2t^3vw]
```

**Questions and To Do List.** • Clean up and write up. • Implement well, compute for everything in sight. • Why are our quantities polynomials rather than just rational functions? • Bounds on their degrees? • Their integrality ( $\mathbb{Z}$ ) properties? • Can everything be re-stated using integrals ( $\int$ )? • Find the 2-variable version (for knots). How complex is it? • What about links / closed components? • Fully digest the "expansion" theorem; include cups. • Explore the (non-)dependence on  $R$ . • Is there a canonical  $R$ ? • What does "group like" mean? • Strand removal? Strand doubling? Strand reversal? • Say something about knot genus. • Find the EK/AT/KV "vertex". • Use as a playground to study associators/braidors. • Restate in topological language. • Study the associated ( $v$ -)braid representations. • Study mirror images and the  $b^+ \leftrightarrow b^-$  involution. • Study ribbon knots. • Make precise the relationship with  $\Gamma$ -calculus and Alexander. • Relate to the coloured Jones polynomial. • Relate with "ordinary"  $q$ -algebra.

optimization as in 2016-09-... optimization - s.n.6 and beyond.

- $k$ -smidgen  $sl_n$ , etc.
- Are there “solvable” CYBE algebras not arising from semi-simple algebras?
- Categorify and appease the Gods.

Help Needed!

balance columns?

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diagram	$n_k^c$ Alexander's $A_+$ Today's / Rozansky's $\rho_k^+$	genus / ribbon unknotting number / amphicheiral	diagram	$n_k^c$ Alexander's $A_+$ Today's / Rozansky's $\rho_k^+$	genus / ribbon unknotting number / amphicheiral
	$0_1^c$ 1 0	0 / ✓ 0 / ✓		$3_1^c$ $t - 1$ $t$	1 / ✗ 1 / ✗
	$4_1^c$ $3 - t$ 0	1 / ✗ 1 / ✓		$5_1^c$ $t^2 - t + 1$ $2t^3 + 3t$	2 / ✗ 2 / ✗
	$5_2^c$ $2t - 3$ $5t - 4$	1 / ✗ 1 / ✗		$6_1^c$ $5 - 2t$ $t - 4$	1 / ✓ 1 / ✗
	$6_2^c$ $-t^2 + 3t - 3$ $t^3 - 4t^2 + 4t - 4$	2 / ✗ 1 / ✗		$6_3^c$ $t^2 - 3t + 5$ 0	2 / ✗ 1 / ✓
	$7_1^c$ $t^3 - t^2 + t - 1$ $3t^3 + 5t^2 + 6t$	3 / ✗ 3 / ✗		$7_2^c$ $3t - 5$ $14t - 16$	1 / ✗ 1 / ✗
	$7_3^c$ $2t^2 - 3t + 3$ $-9t^3 + 8t^2 - 16t + 12$	2 / ✗ 2 / ✗		$7_4^c$ $4t - 7$ $32 - 24t$	1 / ✗ 2 / ✗
	$7_5^c$ $2t^2 - 4t + 5$ $9t^3 - 16t^2 + 29t - 28$	2 / ✗ 2 / ✗		$7_6^c$ $-t^2 + 5t - 7$ $t^3 - 8t^2 + 19t - 20$	2 / ✗ 1 / ✗
	$7_7^c$ $t^2 - 5t + 9$ $8 - 3t$	2 / ✗ 1 / ✗		$8_1^c$ $7 - 3t$ $5t - 16$	1 / ✗ 1 / ✗
	$8_1^c$ $-t^3 + 3t^2 - 3t + 3$ $2t^3 - 8t^2 + 10t^3 - 12t^2 + 13t - 12$	3 / ✗ 2 / ✗		$8_2^c$ $9 - 4t$ 0	1 / ✗ 2 / ✓
	$8_3^c$ $-2t^2 + 5t - 5$ $3t^3 - 8t^2 + 6t - 4$	2 / ✗ 2 / ✗		$8_3^c$ $-t^3 + 3t^2 - 4t + 5$ $-2t^5 + 8t^4 - 13t^3 + 20t^2 - 22t + 24$	3 / ✗ 2 / ✗
	$8_4^c$ $-2t^2 + 6t - 7$ $5t^3 - 20t^2 + 28t - 32$	2 / ✗ 2 / ✗		$8_4^c$ $t^3 - 3t^2 + 5t - 5$ $-t^5 + 4t^4 - 10t^3 + 12t^2 - 13t + 12$	3 / ✗ 1 / ✗