

Rewrite the bi-algebras box following the GWU-1612 talking points

*I wish I had started the talk w/ that story!!!*

Dror Bar-Natan: Talks: MIT-1612: (thanks for accepting my invitation!)

oeβ:=<http://drorbn.net/MIT-1612/>

Work in Progress! Fluid! Help Needed!

### A Poly-Time Knot Polynomial Via Solvable Approximation

**Abstract.** Rozansky [Ro2] and Overbay [Ov] described a **spectacular** knot polynomial that failed to attract the attention it deserved as the first poly-time-computable knot polynomial since Alexander's [Al, 1928] and (in my opinion) as the second most likely knot polynomial (after Alexander's) to carry topological information. With Roland van der Veen, I will explain how to compute the Rozansky polynomial using some new commutator-calculus techniques and a Lie algebra  $\mathfrak{g}_1$  which is at the same time solvable and an approximation of the simple Lie algebra  $\mathfrak{sl}_2$ .

**Theorem** ([BNG], conjectured [MM], elucidated [Ro1]). Let  $J_d(K)$  be the coloured Jones polynomial of  $K$ , in the  $d$ -dimensional representation of  $\mathfrak{sl}_2$ . Writing

$$\left. \frac{(q^{1/2} - q^{-1/2})J_d(K)}{q^{d/2} - q^{-d/2}} \right|_{q=e^h} = \sum_{j,m \geq 0} a_{jm}(K) d^j h^m,$$

"below diagonal" coefficients vanish,  $a_{jm}(K) = 0$  if  $j > m$ , and "on diagonal" coefficients give the inverse of the Alexander polynomial:

$$\left( \sum_{m=0}^{\infty} a_{mm}(K) h^m \right) \cdot A(K)(e^h) = 1.$$

"Above diagonal" we have **Rozansky's Theorem** [Ro3, (1.2)]:

$$J_d(K)(q) = \frac{q^d - q^{-d}}{(q - q^{-1})A(K)(q^d)} \left( 1 + \sum_{k=1}^{\infty} \frac{(q-1)^k R_k(K)(q^d)}{A^{2k}(K)(q^d)} \right).$$

$U \in \mathcal{T}_n \xrightarrow{\tau} \mathcal{T}_{2n} \xrightarrow{\kappa} \mathcal{A}_{2n} \xrightarrow{z} \mathcal{A}_1$

$I \in \mathcal{A}_n \xrightarrow{\tau} \mathcal{A}_{2n} \xrightarrow{\kappa} \mathcal{A}_1$

with  $\mathcal{R} := \kappa(\tau^{-1}(1))$

ribbon  $K \in \mathcal{T}_1$   $z(K) \in \mathcal{A}_1$

Faster is better, leaner is meaner!

$$A^+ = -t^8 + 2t^7 - t^6 - 2t^4 + 5t^3 - 2t^2 - 7t + 13$$

$$\rho_1^+ = 5t^{15} - 18t^{14} + 33t^{13} - 32t^{12} + 2t^{11} + 42t^{10} - 62t^9 - 8t^8 + 166t^7 - 242t^6 + 108t^5 + 132t^4 - 226t^3 + 148t^2 - 11t - 36$$

**The Gold Standard** is set by the "Γ-calculus" Alexander formulas [BNS, BN1]. An  $S$ -component tangle  $T$  has

$$\Gamma(T) \in R_S \times M_{S \times S}(R_S) = \left\{ \begin{array}{c|c} \omega & S \\ \hline S & A \end{array} \right\} \text{ with } R_S := \mathbb{Z}\langle t_a : a \in S \rangle:$$

$$\left( \begin{array}{c|c} a & b \\ \hline a & b \end{array} \right) \rightarrow \begin{array}{c|c} a & b \\ \hline a & b \end{array} \quad T_1 \sqcup T_2 \rightarrow \begin{array}{c|c} \omega_1 \omega_2 & S_1 \ S_2 \\ \hline S_1 \ S_2 & A_1 \ A_2 \end{array}$$

$$\begin{array}{c|c} \omega & a \ b \ S \\ \hline a & \alpha \ \beta \ \theta \\ b & \gamma \ \delta \ \epsilon \\ S & \phi \ \psi \ \Xi \end{array} \xrightarrow{m_c^{ab}} \begin{array}{c|c} (1-\beta)\omega & c \ S \\ \hline c & \gamma + \frac{\alpha\delta}{1-\beta} \ \epsilon + \frac{\theta\delta}{1-\beta} \\ S & \phi + \frac{\alpha\delta}{1-\beta} \ \Xi + \frac{\theta\delta}{1-\beta} \end{array}$$

(Roland: "add to  $A$  the product of column  $b$  and row  $a$ , divide by  $(1 - A_{ab})$ , delete column  $b$  and row  $a$ ".)

For long knots,  $\omega$  is Alexander, and that's the fastest Alexander algorithm I know! Dunfield: 1000-crossing fast.

(There are also formulas for strand doubling and strand reversal).

**Theorem** [EK, Ha, En, Se]. There is a "homomorphic expansion"

$$Z: \left\{ \begin{array}{l} S\text{-component} \\ (v/b)\text{-tangles} \end{array} \right\} \rightarrow \mathcal{A}_S^v := \left\{ \begin{array}{l} \text{AS: } \begin{array}{c} \text{---} \\ \diagup \ \diagdown \\ \text{---} \end{array} + \begin{array}{c} \text{---} \\ \diagdown \ \diagup \\ \text{---} \end{array} = 0 \\ \text{STU: } \begin{array}{c} \text{---} \\ \diagup \ \diagdown \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ \diagdown \ \diagup \\ \text{---} \end{array} - \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \\ \text{HX: } \begin{array}{c} \text{---} \\ \diagup \ \diagdown \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ \diagdown \ \diagup \\ \text{---} \end{array} - \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \end{array} \right.$$

**Genus.**

a ribbon singularity a clasp singularity

**A bit about ribbon knots.** A "ribbon knot" is a knot that can be presented as the boundary of a disk that has "ribbon singularities", but no "clasp singularities". A "slice knot" is a knot in  $S^3 = \partial B^4$  which is the boundary of a non-singular disk in  $B^4$ . Every ribbon knots is clearly slice, yet,

**Conjecture.** Some slice knots are not ribbon.

**Fox-Milnor.** The Alexander polynomial of a ribbon knot is always of the form  $A(t) = f(t)f(1/t)$ . (also for slice)

**Algebras and Invariants.** Given any unital algebra  $A$  (even better if  $A$  is Hopf; typically,  $A \sim \mathcal{U}(\mathfrak{g})$ ), appropriate orange  $R \in A \otimes A$ , and appropriate cuaps  $\in A$ , get an  $A^{\otimes S}$ -valued invariant of pure  $S$ -component tangles:

with

- :  $c$
- :  $u$
- :  $w$
- :  $b$

**Good News.** In theory, enough to know  $R$ , the cuaps, and stitching/multiplication  $m_k^{ij}: A_i \otimes A_j \rightarrow A_k$ .

**Problem.** Extract information out of  $Z$ .

**Textbook Solution.** Use representation theory ... works, slowly.

**Today's Solution** (with van der Veen). For some specific  $\mathfrak{g}$ 's, work in a space of "formulas of a specific type" for elements of  $\mathcal{U}(\mathfrak{g})^{\otimes S}$ :

$$\left\{ \begin{array}{l} \text{ordered perturbed} \\ \text{Gaussian formulas} \end{array} \right\} \rightarrow \mathcal{U}(\mathfrak{g})^{\otimes S}$$

van der Veen

"God created the knots, all else in topology is the work of mortals."

Leopold Kronecker (modified) [www.katlas.org](http://www.katlas.org)

**1-Smidgen  $sl_2$**  Let  $\mathfrak{g}_1$  be the 4-dimensional Lie algebra  $\mathfrak{g}_1 = \langle b, c, u, w \rangle$  over the ring  $R = \mathbb{Q}[\epsilon]/(\epsilon^2 = 0)$ , with  $b$  central and with  $[w, c] = w$ ,  $[c, u] = u$ , and  $[u, w] = b - 2\epsilon c$ , with CYBE  $r_{ij} = (b_i - \epsilon c_i)c_j + u_i w_j$  in  $\mathcal{U}(\mathfrak{g}_1)^{\otimes 2}$ . Over  $\mathbb{Q}$ ,  $\mathfrak{g}_1$  is a **solvable approximation of  $sl_2$** :  $\mathfrak{g}_1 \supset \langle b, u, w, \epsilon b, \epsilon c, \epsilon u, \epsilon w \rangle \supset \langle b, \epsilon b, \epsilon c, \epsilon u, \epsilon w \rangle \supset 0$ . (note:  $\deg(b, c, u, w, \epsilon) = (1, 0, 1, 0, 1)$ )

**0-Smidgen  $sl_2$**   $\otimes$ . Let  $\mathfrak{g}_0$  be  $\mathfrak{g}_1$  at  $\epsilon = 0$ , or  $\mathbb{Q}\langle b, c, u, w \rangle / ([b, \cdot] = 0, [c, u] = u, [c, w] = -w, [u, w] = b$  with  $r_{ij} = b_i c_j + u_i w_j$ . It is  $\mathfrak{b}^* \rtimes \mathfrak{b}$  where  $\mathfrak{b}$  is the 2D Lie algebra  $\mathbb{Q}\langle c, w \rangle$  and  $(b, u)$  is the dual basis of  $(c, w)$ . For topology, it is more valuable than  $\mathfrak{g}_1 / sl_2$ , but topology already got by other means almost everything  $\mathfrak{g}_0$  gives.

**How did these arise?**  $sl_2 = \mathfrak{b}^+ \oplus \mathfrak{b}^- / \mathfrak{h} =: sl_2^+ / \mathfrak{h}$ , where  $\mathfrak{b}^+ = \langle c, w \rangle / [w, c] = w$  is a Lie bialgebra with  $\delta: \mathfrak{b}^+ \rightarrow \mathfrak{b}^+ \otimes \mathfrak{b}^+$  by  $\delta: (c, w) \mapsto (0, c \wedge w)$ . Going back,  $sl_2^+ = \mathcal{D}(\mathfrak{b}^+) = (\mathfrak{b}^+)^* \oplus \mathfrak{b}^+ = \langle b, u, c, w \rangle / \dots$ . **Idea.** Replace  $\delta \rightarrow \epsilon \delta$  over  $\mathbb{Q}[\epsilon]/(\epsilon^{k+1} = 0)$ . At  $k = 0$ , get  $\mathfrak{g}_0$ . At  $k = 1$ , get  $[w, c] = w$ ,  $[w, b'] = -\epsilon w$ ,  $[c, u] = u$ ,  $[b', u] = -\epsilon u$ ,  $[b', c] = 0$ , and  $[u, w] = b' - \epsilon c$ . Now note that  $b' + \epsilon c$  is central, so switch to  $b := b' + \epsilon c$ . This is  $\mathfrak{g}_1$ .

**Ordering Symbols.**  $\otimes$  (*poly* | *specs*) plants the variables of *poly* in  $\mathcal{S}(\otimes \mathfrak{g})$  on several tensor copies of  $\mathcal{U}(\mathfrak{g})$  according to *specs*. E.g.,  $\otimes (c_1^3 u_1 c_2 e^{u_3} w_3 | x: w_3 c_1, y: u_1 u_3 c_2) = w^3 c^3 \otimes u e^u c \in \mathcal{U}(\mathfrak{g})_x \otimes \mathcal{U}(\mathfrak{g})_y$ . This enables the description of elements of  $\hat{\mathcal{U}}(\mathfrak{g})^{\otimes S}$  using commutative polynomials / power series.

**0-Smidgen Invariants.**  $r = Id \in \mathfrak{b}^- \otimes \mathfrak{b}^+$  solves the CYBE  $[r_{12}, r_{13}] + [r_{12}, r_{23}] + [r_{13}, r_{23}] = 0$  in  $\mathcal{U}(\mathfrak{g}_0)^{\otimes 3}$  and, by luck,

$$\begin{array}{c} \nearrow \\ i \end{array} \begin{array}{c} \nearrow \\ j \end{array} = \begin{array}{c} \boxed{+} \\ \nearrow \\ i \end{array} \begin{array}{c} \boxed{+} \\ \nearrow \\ j \end{array} = R_{ij} = e^{r_{ij}} = e^{b_i c_j + u_i w_j} \in \mathcal{U}(\mathfrak{g}_{0,i} \oplus \mathfrak{g}_{0,j})$$

solves YB/R3.

**Lemma.**  $R_{ij} = e^{b_i c_j + u_i w_j} = \otimes (\exp(b_i c_j + \frac{e^{b_i} - 1}{b_i} u_i w_j) | i: u_i, j: c_j w_j)$

**Example.**  $Z(T_0) = \sum_{m,n} \frac{b^{m-n} (e^{b_i} - 1)^n}{m! n!} u^m \otimes c^m w^n$

$$\otimes \left( \exp \left( b_5 c_1 + \frac{e^{b_5} - 1}{b_5} u_5 w_1 + b_2 c_4 + \frac{e^{b_2} - 1}{b_2} u_2 w_4 - b_3 c_6 + \frac{e^{b_3} - 1}{b_3} u_3 w_6 \right) \right)$$

"ucw form"

$$x: c_1 w_1 u_2, y: u_3 c_4 w_4 u_5 c_6 w_6 = \otimes (\zeta | x: u_x c_x w_x, y: u_y c_y w_y)$$

**Goal.** Write  $\zeta$  as a Gaussian:  $\omega e^{L+Q}$  where  $L$  bilinear in  $b_i$  and  $c_i$  with integer coefficients,  $Q$  a balanced quadratic in  $u_i$  and  $w_i$  with coefficients in  $R_S := \mathbb{Q}(b_i, e^{b_i})$ , and  $\omega \in R_S$ .

**The Big  $\mathfrak{g}_0$  Lemma.** Under  $[c, u] = u$ ,  $[c, w] = -w$ , and  $[u, w] = b$ :

- $N^{cu} := \otimes (e^{\gamma c + \beta u} | uc) \stackrel{\cong}{=} \otimes (e^{\gamma c + e^{\beta} \beta u} | cu)$  (means  $e^{\beta u} e^{\gamma c} = e^{\gamma c} e^{\beta u}$ )
- $N^{wc} := \otimes (e^{\gamma c + \alpha w} | wc) \stackrel{\cong}{=} \otimes (e^{\gamma c + e^{\alpha} \alpha w} | cw)$  ... in the  $\{ax + b\}$  group)
- $\otimes (e^{\alpha w + \beta u} | wu) = \otimes (e^{-b\alpha\beta + \alpha w + \beta u} | wu)$  (the Weyl relations)
- $\otimes (e^{\delta u v} | wu) e^{\beta u} = e^{\beta u} \otimes (e^{\delta u v} | wu)$ , with  $\gamma = (1 + b\delta)^{-1}$
- (a. expand and crunch. b. use  $w = b\hat{x}, u = \partial_x$ . c. use "scatter and glow".)
- $\otimes (e^{\delta u v} | wu) = \otimes (v e^{\delta u v} | wu)$  (same techniques)
- $N^{wu} := \otimes (e^{\beta u + \alpha w + \delta u v} | wu) \stackrel{\cong}{=} \otimes (v e^{-b\gamma\alpha\beta + \gamma\alpha w + \gamma\beta u + \gamma\delta u v} | wu)$
- $N_k^{c_i c_j} := \otimes (\zeta | c_i c_j) \stackrel{\cong}{=} \otimes (\zeta / (c_i, c_j \rightarrow c_k) | c_k)$

**Sneaky.**  $\alpha$  may contain (other)  $u$ 's,  $\beta$  may contain (other)  $w$ 's.

**Strand Stitching,**  $m_{ij}^k$ , is defined as the composition

$$u_i c_i \overline{w_j} u_j c_j w_j \xrightarrow{N_x^{u_i u_j}} u_i \overline{c_i} u_x \overline{w_x} c_j w_j \xrightarrow{N_x^{c_i c_x} // N_x^{u_x c_j}} \overline{u_i} u_x \overline{c_x} c_x \overline{w_x} w_j \xrightarrow{i, j, x \rightarrow k} u_k c_k w_k$$

On to 1-smidgen invariants, where much is the same...

**The Big  $\mathfrak{g}_1$  Lemma.** Parts 1 and 6 are the same, yet

$$\otimes (e^{\alpha w + \beta u + \delta u v} | wu) = \otimes (v(1 + \epsilon v \Lambda) e^{v(-b\alpha\beta + \alpha w + \beta u + \delta u v)} | ucw)$$

Here  $\Lambda$  is for  $\Lambda\delta\gamma\sigma\varsigma$ , "a principle of order and knowledge", a balanced quartic in  $\alpha, \beta, u, c$ , and  $w$ :

$$\begin{aligned} \Lambda = & -bv(\alpha^2\beta^2v^2 + 4\alpha\beta\delta v + 2\delta^2)/2 + \beta^2\delta v^3(b\delta + 2)u^2/2 \\ & + \delta^3v^3(3b\delta + 4)u^2w^2/2 + \beta\delta^2v^3(2b\delta + 3)u^2w \\ & + \alpha\delta^2v^3(2b\delta + 3)uw^2 + 2\delta v^2(b\delta + 2)(\alpha\beta v + \delta)uw \\ & + \alpha^2\delta v^3(b\delta + 2)w^2/2 + 2(\alpha\beta v + \delta)c + 2\beta\delta vuc + 2\delta^2vucw \\ & + 2\alpha\delta v cw + \beta v^2(\alpha\beta v + 2\delta)u + \alpha v^2(\alpha\beta v + 2\delta)w. \end{aligned}$$

**Proof.** A lengthy computation. (Verification:  $\omega\epsilon\beta/\text{Big}$ )

**Problem.** We now need to normal-order perturbed Gaussians!

**Solution.** Borrow some tactics from QFT:

$$\otimes (\epsilon P(c, u) e^{\gamma c + \beta u} | uc) = \otimes (\epsilon P(\partial_\gamma, \partial_\beta) e^{\gamma c + \beta u} | uc) = \otimes (\epsilon P(\partial_\gamma, \partial_\beta) e^{\gamma c + e^{-\gamma} \beta u} | cu),$$

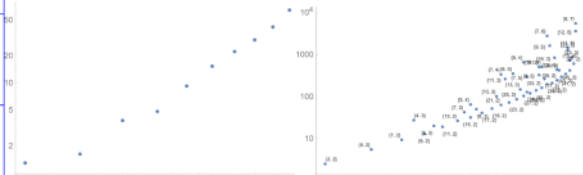
and likewise  $\otimes (\epsilon P(u, w) e^{\alpha w + \beta u + \delta u v} | wu) = \otimes (\epsilon P(\partial_\beta, \partial_\alpha) v e^{(-b\alpha\beta + \alpha w + \beta u + \delta u v)} | ucw)$

Finally, the values of the generators  $\gamma, \beta, \alpha, \delta$ , and  $\underline{u}$ , are set by solving many equations, non-uniquely.

**Pragmatic Simplifications.** Set  $t := e^b$ , work with  $v := (t - 1)u/b$ , and set  $\mathbb{E}(\omega, L, Q, P) := \otimes (\omega^{-1} e^{L+Q/\omega} (1 + \epsilon \omega^{-4} P) : (i: v_i c_i w_i))$ . Now  $\omega \in R_S := \mathbb{Z}[t_i, t_i^{-1}]$  is Laurent,  $L = \sum l_{ij} \log(t_i) c_j$  with  $l_{ij} \in \mathbb{Z}$ ,  $Q = \sum q_{ij} v_i w_j$  with  $q_{ij} \in R_S$ , and  $P$  is a quartic polynomial in  $v_i, c_j, w_k$  with coefficients in  $R_S$ . The operations are lightly modified, and the  $\Lambda\delta\gamma\sigma\varsigma$  and the values of the generators become somewhat simpler, as in the implementation below.

**Rough complexity estimate,** after  $t_k \rightarrow t$ :  $n$ : xing  $\frac{n}{A} \sum_{d=0}^4 \frac{B}{E} \frac{W^4 - d}{F} \frac{n^2}{G} = n^3 W^4 \in [n^5, n^7]$  number;  $w$ : width, maybe  $\sim \sqrt{n}$ .  $A$ : go over stitchings in order.  $B$ : multiplication ops per  $N^{u_i w_j}$ .  $d$ : deg of  $u_i, w_j$  in  $P$ .  $E$ : #terms of deg  $d$  in  $P$ .  $F$ : ops per term.  $G$ : cost per polynomial multiplication op.

**Experimental Analysis ( $\omega\epsilon\beta/\text{Exp}$ ).** Log-log plots of computation time (sec) vs. crossing number, for all knots with up to 12 crossings (mean times) and for all torus knots with up to 48 crossings:



**Conjecture** (checked on the 12 collections). Given a knot  $K$  with Alexander polynomial  $A$ , there is a polynomial  $\rho_1$  such that

$$P = A^2 \frac{(t-1)^3 \rho_1 + t^2(2vw + (1-t)(1-2c))AA'}{(1-t)t}$$

Furthermore,  $A$  and  $\rho_1$  are symmetric under  $t \rightarrow t^{-1}$ , so let  $A^+$  and  $\rho_1^+$  be their "positive parts", so e.g.,  $\rho_1(t) = \rho_1^+(t) + \rho_1^+(t^{-1}) - \rho_1^+(0)$ .

**Power.** On the 250 knots with at most 10 crossings, the pair  $(A, \rho_1)$  attains 250 distinct values, while (Khovanov, HOMFLY-PT) attains only 249 distinct values. To 11 crossings the numbers are (802, 788, 772) and to 12 they are (2978, 2883, 2786).

**Genus.** Up to 12 xings, always  $\deg \rho_1^+ \leq 2g - 1$ , where  $g$  is the 3-genus of  $K$  (equality for 2530 knots). This gives a lower bound on  $g$  in terms of  $\rho_1$  (conjectural, but undoubtedly true). This bound is often weaker than the Alexander bound, yet for 10 of the 12-xing Alexander failures it does give the right answer.

Switch to ucw conventions here too!

Demo Programs for 0-Co.

ωβ/θ/Demo

$R_{\theta, i, j}^+ := \mathbb{E}[b_i c_j + b_i^{-1} (e^{b_i} - 1) u_i w_j];$   
 $R_{\theta, i, j}^- := \mathbb{E}[-b_i c_j + b_i^{-1} (e^{-b_i} - 1) u_i w_j];$

The R-matrices

CF[ω<sub>-</sub>. E[Q<sub>-</sub>]] := Simplify[ω<sub>-</sub> E[Simplify[Q<sub>-</sub>]]; Utilities

E /: E[Q1<sub>-</sub>] E[Q2<sub>-</sub>] := CF@E[Q1<sub>-</sub> + Q2<sub>-</sub>];

ω<sub>1</sub><sub>-</sub>. E[Q1<sub>-</sub>] = ω<sub>2</sub><sub>-</sub>. E[Q2<sub>-</sub>] := Simplify[ω<sub>1</sub><sub>-</sub> = ω<sub>2</sub><sub>-</sub> ∧ Q1 = Q2];

N<sub>(k:w|u)</sub><sub>i</sub> c<sub>j</sub> → h<sub>-</sub> [ω<sub>-</sub>. E[Q<sub>-</sub>]] := CF[ Normal Ordering Operators

ω E[e<sup>α</sup> X<sub>h</sub> + γ C<sub>h</sub> + (Q / . C<sub>j</sub> | X<sub>i</sub> → θ)] / . {γ → ∂<sub>C<sub>j</sub></sub> Q, α → ∂<sub>X<sub>i</sub></sub> Q};

N<sub>w<sub>i</sub></sub> u<sub>j</sub> → h<sub>-</sub> [ω<sub>-</sub>. E[Q<sub>-</sub>]] := CF[

v ω E[-b<sub>h</sub> v α β + v β u<sub>h</sub> + v α w<sub>h</sub> + v δ u<sub>h</sub> w<sub>h</sub> + (Q / . w<sub>i</sub> | u<sub>j</sub> → θ)] / .

v → (1 + b<sub>h</sub> δ)<sup>-1</sup> / .

{α → ∂<sub>w<sub>i</sub></sub> Q / . u<sub>j</sub> → θ, β → ∂<sub>u<sub>j</sub></sub> Q / . w<sub>i</sub> → θ, δ → ∂<sub>w<sub>i</sub></sub> u<sub>j</sub> Q};

m<sub>i</sub> c<sub>j</sub> → h<sub>-</sub> [Z<sub>-</sub>] := Module[{X, Z}, Stitching

CF[(Z // N<sub>w<sub>i</sub></sub> u<sub>j</sub> → x // N<sub>c<sub>i</sub></sub> v<sub>x</sub> → x // N<sub>w<sub>x</sub></sub> c<sub>j</sub> → x) / . Z<sub>-</sub> i | j | x → z<sub>h</sub>]]

T<sub>θ</sub> = R<sub>θ, 5, 1</sub> R<sub>θ, 2, 4</sub> R<sub>θ, 3, 6</sub> Some calculations for T<sub>θ</sub>

$\mathbb{E}\left[b_5 c_1 + b_2 c_4 - b_3 c_6 + \frac{(-1 \cdot e^{b_5}) u_5 w_1}{b_5} + \frac{(-1 \cdot e^{b_2}) u_2 w_4}{b_2} + \frac{(-1 \cdot e^{-b_3}) u_3 w_6}{b_3}\right]$

$T_\theta // m_{1,2 \rightarrow 1} // m_{3,4 \rightarrow 3} // m_{3,5 \rightarrow 3} // m_{3,6 \rightarrow 3}$

$\frac{1}{1 - (-1 \cdot e^{b_1})} \mathbb{E}\left[b_3 c_1 + b_1 c_3 - b_3 c_3 + \frac{e^{b_3} (-1 \cdot e^{b_1}) (-1 \cdot e^{b_3}) u_1 w_1}{(-e^{b_3} - e^{b_1} - e^{b_3}) b_1} - \frac{e^{b_1} (-1 \cdot e^{b_3}) u_3 w_1}{(-1 \cdot (-1 \cdot e^{b_1}) (-1 \cdot e^{b_3})) b_3} - \frac{e^{-b_3} (-1 \cdot e^{b_3}) u_3 w_3}{b_3} - \frac{e^{-b_3} (-1 \cdot e^{b_1}) (-e^{b_3} b_3 u_1 \cdot e^{b_1} (-1 \cdot e^{b_3}) b_1 u_3 w_3)}{b_1 | b_3 - (-1 \cdot e^{b_1}) (-1 \cdot e^{b_3}) b_3}\right]$

Verifying meta-associativity

Q<sub>0</sub> = E[Sum[f<sub>i</sub> c<sub>i</sub>, {i, 3}] + Sum[f<sub>i, j</sub> u<sub>i</sub> w<sub>j</sub>, {i, 3}, {j, 3}]]

E[C<sub>1</sub> f<sub>1</sub> + C<sub>2</sub> f<sub>2</sub> + C<sub>3</sub> f<sub>3</sub> + u<sub>1</sub> w<sub>1</sub> f<sub>1,1</sub> + u<sub>1</sub> w<sub>2</sub> f<sub>1,2</sub> + u<sub>1</sub> w<sub>3</sub> f<sub>1,3</sub> + u<sub>2</sub> w<sub>1</sub> f<sub>2,1</sub> + u<sub>2</sub> w<sub>2</sub> f<sub>2,2</sub> + u<sub>2</sub> w<sub>3</sub> f<sub>2,3</sub> + u<sub>3</sub> w<sub>1</sub> f<sub>3,1</sub> + u<sub>3</sub> w<sub>2</sub> f<sub>3,2</sub> + u<sub>3</sub> w<sub>3</sub> f<sub>3,3</sub>]

(Q<sub>0</sub> // m<sub>1,2 → 1</sub> // m<sub>1,3 → 1</sub>) = (Q<sub>0</sub> // m<sub>2,3 → 2</sub> // m<sub>1,2 → 1</sub>)

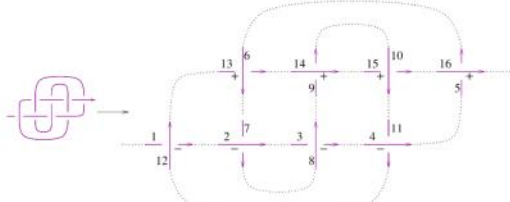
True

t<sub>1</sub> = R<sub>θ, 1, 2</sub> R<sub>θ, 3, 4</sub> R<sub>θ, 5, 6</sub> // m<sub>3,5 → x</sub> // m<sub>1,6 → y</sub> // m<sub>2,4 → z</sub> Testing R<sub>3</sub>

E[b<sub>x</sub> c<sub>y</sub> + b<sub>x</sub> c<sub>z</sub> + b<sub>y</sub> c<sub>z</sub> +  $\frac{e^{b_x} (-1 \cdot e^{b_y}) u_y w_z}{b_y} + \frac{(-1 \cdot e^{b_x}) u_x (w_y w_z)}{b_x}$ ]

t<sub>1</sub> = (R<sub>θ, 1, 2</sub> R<sub>θ, 3, 4</sub> R<sub>θ, 5, 6</sub> // m<sub>1,3 → x</sub> // m<sub>2,5 → y</sub> // m<sub>4,6 → z</sub>)

True



z<sub>1</sub> = R<sub>θ, 12, 1</sub> R<sub>θ, 2, 7</sub> R<sub>θ, 8, 3</sub> R<sub>θ, 4, 11</sub> R<sub>θ, 16, 5</sub> R<sub>θ, 6, 13</sub> R<sub>θ, 14, 9</sub> R<sub>θ, 10, 15</sub>;

Do[z<sub>1</sub> = (z<sub>1</sub> // m<sub>1, n → 1</sub>) / . b<sub>-</sub> → b, {n, 2, 16}];

{CF@z<sub>1</sub>, KnotData[{8, 17}, "AlexanderPolynomial"]}[t]

$\left\{-\frac{e^{3b} b | \theta |}{1 - 4 e^{b_8} e^{2b_{11}} e^{3b_8} e^{4b_4} e^{5b_6} e^{6b_5}}, 11 - \frac{1}{t^3} + \frac{4}{t^2} - \frac{8}{t} - 8t + 4t^2 - t^3\right\}$

Demo Programs for 1-Co.

ωβ/θ/Demo

$\Lambda[k_-] := ((t_h - 1) (2 (\alpha \beta + \delta \mu)^2 - \alpha^2 \beta^2) - 4 v_h c_h w_h \delta^2 \mu^2 - \delta (1 + \mu) (w_h^2 \alpha^2 + v_h^2 \beta^2) - v_h^2 w_h^2 \delta^3 (1 + 3 \mu) - 2 (\alpha \beta + 2 \delta \mu + v_h w_h \delta^2 (1 + 2 \mu) + 2 c_h \delta \mu^2) (w_h \alpha + v_h \beta) - 4 (c_h \mu^2 + v_h w_h \delta (1 + \mu)) (\alpha \beta + \delta \mu)) (1 + t_h) / 4;$  The Λόγος

$R_{i, j}^+ := \mathbb{E}[1, \text{Log}[t_i] c_j, v_i w_j, v_i c_i w_j + c_i c_j + v_i^2 w_j^2 / 4];$

$R_{i, j}^- := \mathbb{E}[1, -\text{Log}[t_i] c_j, -t_i^{-1} v_i w_j,$  The Generators

$t_i^{-1} v_i c_j w_j - c_i c_j - t_i^{-2} v_i^2 w_j^2 / 4];$

$(ur_i := \mathbb{E}[t_i^{1/2}, \theta, \theta, c_i t_i^2]; nr_i := \mathbb{E}[t_i^{1/2}, \theta, \theta, -c_i t_i^2];)$

Differential Polynomials

DP<sub>x → 0, y → 0, z → 0</sub>[P<sub>-</sub>] [f<sub>-</sub>] := (\* means P[∂<sub>α</sub>, ∂<sub>β</sub>] [f] \*)

Total[CoefficientRules[P, {x, y}]] / .

{(m<sub>-</sub>, n<sub>-</sub>) → c<sub>-</sub>} → c D[f, {α, m}, {β, n}]]

CF[ε<sub>-</sub> E] := Expand /@ Together /@ ε;

E /: E[ω<sub>1</sub><sub>-</sub>, L<sub>1</sub><sub>-</sub>, Q<sub>1</sub><sub>-</sub>, P<sub>1</sub><sub>-</sub>] E[ω<sub>2</sub><sub>-</sub>, L<sub>2</sub><sub>-</sub>, Q<sub>2</sub><sub>-</sub>, P<sub>2</sub><sub>-</sub>] :=

CF@E[ω<sub>1</sub> ω<sub>2</sub>, L<sub>1</sub> + L<sub>2</sub>, ω<sub>2</sub> Q<sub>1</sub> + ω<sub>1</sub> Q<sub>2</sub>, ω<sub>2</sub><sup>4</sup> P<sub>1</sub> + ω<sub>1</sub><sup>4</sup> P<sub>2</sub>];

Normal Ordering Operators

N<sub>c<sub>j</sub></sub> (x:v|w) i → h<sub>-</sub> [E[ω<sub>-</sub>, L<sub>-</sub>, Q<sub>-</sub>, P<sub>-</sub>]] := With[{q = e<sup>γ</sup> β X<sub>h</sub> + γ C<sub>h</sub>}, CF

E[ω, γ C<sub>h</sub> + (L / . C<sub>j</sub> → θ), ω e<sup>γ</sup> β X<sub>h</sub> + (Q / . X<sub>i</sub> → θ),

e<sup>-Q</sup> DP<sub>c<sub>j</sub> → 0, x<sub>i</sub> → 0, β</sub>[e<sup>Q</sup>] / . {γ → ∂<sub>C<sub>j</sub></sub> L, β → ω<sup>-1</sup> ∂<sub>X<sub>i</sub></sub> Q}];

N<sub>w<sub>i</sub></sub> v<sub>j</sub> → h<sub>-</sub> [E[ω<sub>-</sub>, L<sub>-</sub>, Q<sub>-</sub>, P<sub>-</sub>]] :=

With[{q = ((1 - t<sub>h</sub>) α β + β v<sub>h</sub> + α w<sub>h</sub> + δ v<sub>h</sub> w<sub>h</sub>) / μ}, CF

E[μ ω, L, μ ω q + μ (Q / . w<sub>i</sub> | v<sub>j</sub> → θ),

μ<sup>4</sup> e<sup>-Q</sup> DP<sub>w<sub>i</sub> → 0, v<sub>j</sub> → 0, β</sub>[e<sup>Q</sup>] + ω<sup>4</sup> Λ[k<sub>-</sub>] / . μ → 1 + (t<sub>h</sub> - 1) δ / .

{α → ω<sup>-1</sup> (∂<sub>w<sub>i</sub></sub> Q / . v<sub>j</sub> → θ), β → ω<sup>-1</sup> (∂<sub>v<sub>j</sub></sub> Q / . w<sub>i</sub> → θ),

δ → ω<sup>-1</sup> ∂<sub>w<sub>i</sub></sub> v<sub>j</sub> Q}];

m<sub>i</sub> c<sub>j</sub> → h<sub>-</sub> [Z<sub>-</sub> E] := Module[{X, Z},

CF[(Z // N<sub>w<sub>i</sub></sub> v<sub>j</sub> → x // N<sub>c<sub>i</sub></sub> v<sub>x</sub> → x // N<sub>w<sub>x</sub></sub> c<sub>j</sub> → x) / . Z<sub>-</sub> i | j | x → z<sub>h</sub>]]

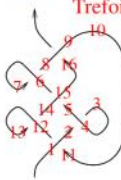
Stitching

z<sub>2</sub> = R<sub>i, 11</sub> R<sub>i, 2</sub> nr<sub>3</sub> R<sub>i, 5, 5</sub> R<sub>6, 8</sub> ur<sub>7</sub> R<sub>9, 16</sub> nr<sub>10</sub> R<sub>i, 12, 14</sub> ur<sub>13</sub>;

(Do[z<sub>2</sub> = z<sub>2</sub> // m<sub>1, k → 1</sub>, {k, 2, 16}];

z<sub>2</sub> = z<sub>2</sub> / . a<sub>-1</sub> → a)

The 0-Framed Trefoil



E[-1 + 1/t + t, θ, θ,

16 + 2c/t<sup>4</sup> - 1/t<sup>3</sup> - 6c/t<sup>3</sup> + 4/t<sup>2</sup> + 10c/t<sup>2</sup> - 10/t - 8c/t - 18t + 8ct +

14t<sup>2</sup> - 10ct<sup>2</sup> - 7t<sup>3</sup> + 6ct<sup>3</sup> + 2t<sup>4</sup> - 2ct<sup>4</sup> + 2vw -

2vw/t<sup>4</sup> + 4vw/t<sup>3</sup> - 6vw/t<sup>2</sup> + 2vw - 6tvw + 4t<sup>2</sup>vw - 2t<sup>3</sup>vw]

Questions and To Do List. • Clean up and write up. • Implement well, compute for everything in sight. • Why are our quantities polynomials rather than just rational functions? • Bounds on their degrees? • Their integrality (Z) properties? • Can everything be re-stated using integrals (∫)? • Find the 2-variable version (for knots). How complex is it? • What about links / closed components? • Fully digest the “expansion” theorem; include cuaps. • Explore the (non-)dependence on R. • Is there a canonical R? • What does “group like” mean? • Strand removal? Strand doubling? Strand reversal? • Say something about knot genus. • Find the EK/AT/KV “vertex”. • Use as a playground to study associators/braidors. • Restate in topological language. • Study the associated (v-)braid representations. • Study mirror images and the b<sup>+</sup> ↔ b<sup>-</sup> involution. • Study ribbon knots. • Make precise the relationship with Γ-calculus and Alexander. • Relate to the coloured Jones polynomial. • Relate with “ordinary” q-algebra. • k-smidgen sl<sub>n</sub>, etc. • Are there “solvable” CYBE algebras not arising from semi-simple algebras? • Categorify and appease the Gods.

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diagram	$n_k^i$ Alexander's $A_+$ Today's / Rozansky's $\rho_1^+$	genus / ribbon unknotting number / amphicheiral	diagram	$n_k^i$ Alexander's $A_+$ Today's / Rozansky's $\rho_1^+$	genus / ribbon unknotting number / amphicheiral
	$0_1^1$ 1 0	0 / ✓ 0 / ✓		$3_1^1$ $t - 1$ $t$	1 / ✗ 1 / ✗
	$4_1^1$ $3 - t$ 0	1 / ✗ 1 / ✓		$5_1^1$ $t^2 - t + 1$ $2t^3 + 3t$	2 / ✗ 2 / ✗
	$5_2^1$ $2t - 3$ $5t - 4$	1 / ✗ 1 / ✗		$6_1^1$ $5 - 2t$ $t - 4$	1 / ✓ 1 / ✗
	$6_2^1$ $-t^2 + 3t - 3$ $t^3 - 4t^2 + 4t - 4$	2 / ✗ 1 / ✗		$6_3^1$ $t^2 - 3t + 5$ 0	2 / ✗ 1 / ✓
	$7_1^1$ $t^3 - t^2 + t - 1$ $3t^5 + 5t^3 + 6t$	3 / ✗ 3 / ✗		$7_2^1$ $3t - 5$ $14t - 16$	1 / ✗ 1 / ✗
	$7_3^1$ $2t^2 - 3t + 3$ $-9t^3 + 8t^2 - 16t + 12$	2 / ✗ 2 / ✗		$7_4^1$ $4t - 7$ $32 - 24t$	1 / ✗ 2 / ✗
	$7_5^1$ $2t^2 - 4t + 5$ $9t^3 - 16t^2 + 29t - 28$	2 / ✗ 2 / ✗		$7_6^1$ $-t^2 + 5t - 7$ $t^3 - 8t^2 + 19t - 20$	2 / ✗ 1 / ✗
	$7_7^1$ $t^2 - 5t + 9$ $8 - 3t$	2 / ✗ 1 / ✗		$8_1^1$ $7 - 3t$ $5t - 16$	1 / ✗ 1 / ✗
	$8_1^1$ $-t^3 + 3t^2 - 3t + 3$ $2t^5 - 8t^4 + 10t^3 - 12t^2 + 13t - 12$	3 / ✗ 2 / ✗		$8_2^1$ $9 - 4t$ 0	1 / ✗ 2 / ✓
	$8_3^1$ $-2t^2 + 5t - 5$ $3t^3 - 8t^2 + 6t - 4$	2 / ✗ 2 / ✗		$8_3^1$ $-t^3 + 3t^2 - 4t + 5$ $-2t^5 + 8t^4 - 13t^3 + 20t^2 - 22t + 24$	3 / ✗ 2 / ✗
	$8_4^1$ $-2t^2 + 6t - 7$ $5t^3 - 20t^2 + 28t - 32$	2 / ✗ 2 / ✗		$8_4^1$ $t^3 - 3t^2 + 5t - 5$ $-t^5 + 4t^4 - 10t^3 + 12t^2 - 13t + 12$	3 / ✗ 1 / ✗
	$8_5^1$ $2t^2 - 6t + 9$ $-t^3 + 4t^2 - 12t + 16$	2 / ✓ 2 / ✗		$8_5^1$ $-t^3 + 3t^2 - 5t + 7$ 0	3 / ✓ 1 / ✓
	$8_6^1$ $t^3 - 3t^2 + 6t - 7$ $-t^5 + 4t^4 - 11t^3 + 16t^2 - 21t + 20$	3 / ✗ 2 / ✗		$8_6^1$ $-2t^2 + 7t - 9$ $5t^3 - 24t^2 + 39t - 44$	2 / ✗ 1 / ✗
	$8_7^1$ $t^2 - 7t + 13$ 0	2 / ✗ 2 / ✓		$8_7^1$ $2t^2 - 7t + 11$ $-t^3 + 4t^2 - 14t + 20$	2 / ✗ 1 / ✗
	$8_8^1$ $-2t^2 + 8t - 11$ $5t^3 - 28t^2 + 57t - 68$	2 / ✗ 1 / ✗		$8_8^1$ $3t^2 - 8t + 11$ $21t^3 - 64t^2 + 120t - 140$	2 / ✗ 2 / ✗
	$8_9^1$ $t^3 - 4t^2 + 8t - 9$ $t^5 - 6t^4 + 17t^3 - 28t^2 + 35t - 36$	3 / ✗ 2 / ✗		$8_9^1$ $-t^3 + 4t^2 - 8t + 11$ 0	3 / ✗ 1 / ✓
	$8_{10}^1$ $-t^3 + 5t^2 - 10t + 13$ 0	3 / ✗ 2 / ✓		$8_{10}^1$ $t^3 - t^2 + 1$ $-3t^5 - 4t^2 - 3t$	3 / ✗ 3 / ✗
	$8_{11}^1$ $t^2 - 2t + 3$ $4t - 4$	2 / ✓ 1 / ✗		$8_{11}^1$ $-t^2 + 4t - 5$ $t^5 - 8t^2 + 16t - 20$	2 / ✗ 1 / ✗