



Shifted Partial Quadratics, their Pushforwards, and Signature Invariants for Tangles

Abstract. Following a general discussion of the computation of zombians of unfinished columbaria (with examples), I will tell you about my recent joint work w/ Jessica Liu on what we feel is the “textbook” extension of knot signatures to tangles, which for unknown reasons, is not in any of the textbooks that we know.



Jessica Liu



Columbaria in an East Sydney Cemetery



Jacobian, Hamiltonian, Zombian

Kashaev's Conjecture [Ka]

$$\text{For knots, } \sigma_{Kas} = 2\sigma_{TL}.$$

Liu's Theorem [Li].

A **Partial Quadratic (PQ)** on V is a quadratic Q defined only on a subspace $\mathcal{D}_Q \subset V$. We add PQs with $\mathcal{D}_{Q_1+Q_2} := \mathcal{D}_{Q_1} \cap \mathcal{D}_{Q_2}$. Given a linear $\psi: V \rightarrow W$ and a PQ Q on W , there is an obvious pullback ψ^*Q , a PQ on V .

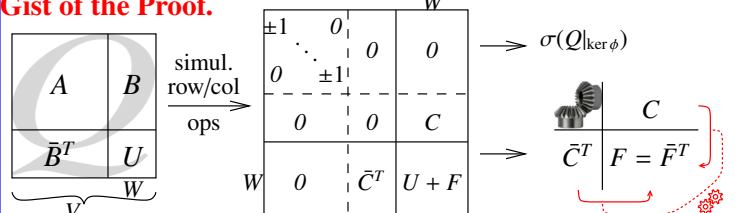
Theorem 1. Given a linear $\phi: V \rightarrow W$ and a PQ Q on V , there is a unique pushforward PQ ϕ_*Q on W such that for every PQ U on W , $\sigma_V(Q + \phi^*U) = \sigma_{\ker \phi}(Q|_{\ker \phi}) + \sigma_W(U + \phi_*Q)$.

(If you must, $\mathcal{D}(\phi_*Q) = \phi(\text{ann}_Q(\mathcal{D}(Q) \cap \ker \phi))$ and $(\phi_*Q)(w) = Q(v)$, where v is s.t. $\phi(v) = w$ and $Q(v, \text{rad } Q|_{\ker \phi}) = 0$).

Nick's proof!

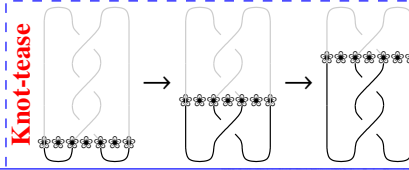
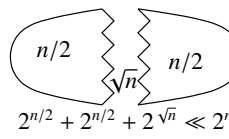
Prior Art on signatures for tangles / braids. Gambaudo and Ghys [GG], Cimasoni and Conway [CC], Conway [Co], Merz [Me]. All define signatures of tangles / braids by first closing them to links and then work hard to derive composition properties.

Gist of the Proof.



Why Tangles? • Faster!

- Conceptually clearer proofs of invariance (and of skein relations).
- Often fun and consequential:
 - The Jones Polynomial \rightsquigarrow The Temperley-Lieb Algebra.
 - Khovanov Homology \rightsquigarrow “Unfinished complexes”, complexes in a category.
 - The Kontsevich Integral \rightsquigarrow Associators.
 - HFK \rightsquigarrow OMG, type D , type A , $\mathcal{A}_\infty, \dots$



Computing Zombians of Unfinished Columbaria.

- Must be no slower than for finished ones.
- Future zombies must be able to complete the computation.
- Future zombies must not even know the size of the task that today's zombies were facing.
- We must be able to extend to ZPUCs, Zombie Processed Unfinished Columbaria!

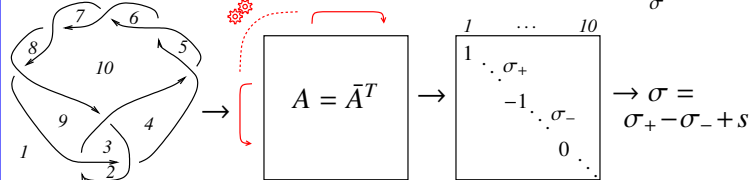


Columbarium near Assen

Example / Exercise. Compute the determinant of a $1,000 \times 1,000$ matrix in which 50 entries are not yet given.

Homework / Research Projects. • What with ZPUCs? • Use this to get an Alexander tangle invariant.

Reminders. {knots} \rightleftharpoons {matrices / quadratic forms} $\xrightarrow{\text{signature } \sigma}$ \mathbb{Z} :



With $|\omega| = 1, t = 1 - \omega, r = t + \bar{t}, v = \text{Re}(\omega),$ and $u = \text{Re}(\omega^{1/2})$:

$X_{-i,j,k,-l}$	Tristram-Levine (TL)	Kashaev (Kas)
	$A += \begin{pmatrix} -r & -t & 2t & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ 2\bar{t} & t & -r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A += \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$
$\bar{X}_{-i,j,k,-l}$	$A += \begin{pmatrix} r & -t & -2\bar{t} & \bar{t} \\ -\bar{t} & 0 & \bar{t} & 0 \\ -2t & t & r & -\bar{t} \\ t & 0 & -t & 0 \end{pmatrix}$	$A -= \begin{pmatrix} v & u & 1 & u \\ u & 1 & u & 1 \\ 1 & u & v & u \\ u & 1 & u & 1 \end{pmatrix}$

A **Shifted Partial Quadratic (SPQ)** on V is a pair $S = (s \in \mathbb{Z}, Q \text{ a PQ on } V)$. addition also adds the shifts, pullbacks keep the shifts, yet $\phi_*S := (s + \sigma_{\ker \phi}(Q|_{\ker \phi}), \phi_*Q)$ and $\sigma(S) := s + \sigma(Q)$.

Theorem 1' (Reciprocity). Given $\phi: V \rightarrow W$, for SPQs S on V and U on W we have $\sigma_V(S + \phi^*U) = \sigma_W(U + \phi_*S)$ (and this characterizes ϕ_*S).

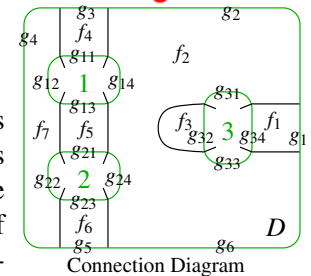
Theorem 2. ψ^* and ϕ_* are functorial.

Theorem 3. If, as on the right, $\beta\alpha = \delta\gamma$ and α and γ are surjective, then $\alpha_*\gamma^* = \beta^*\delta_*$.

FIX!

Definition. $S \begin{pmatrix} \text{cyclic set} \\ \text{on } \langle g_i \rangle \end{pmatrix} := \left\{ \begin{matrix} \text{SPQ } S \\ \text{on } \langle g_i \rangle \end{matrix} \right\}$.

Theorem 4 ~~X~~. $\{S(\text{cyclic sets})\}$ is a planar algebra, with compositions $S(D)((S_i)) := \phi_*^D(\psi_D^*(\bigoplus_i S_i))$, where $\psi_D: \langle f_i \rangle \rightarrow \langle g_{ai} \rangle$ maps every face of D to the sum of the input gaps adjacent to it and $\phi^D: \langle f_i \rangle \rightarrow \langle g_i \rangle$ maps every face to the sum of the output gaps adjacent to it. So for our D , ψ_D is $f_1 \mapsto g_{34}, f_2 \mapsto g_{31} + g_{14} + g_{24} + g_{33}, f_3 \mapsto g_{32}, f_4 \mapsto g_{11}, f_5 \mapsto g_{13} + g_{21}, f_6 \mapsto g_{23}, f_7 \mapsto g_{12} + g_{22}$ and ϕ^D is $f_1 \mapsto g_1, f_2 \mapsto g_2 + g_6, f_3 \mapsto 0, f_4 \mapsto g_3, f_5 \mapsto 0, f_6 \mapsto g_5, f_7 \mapsto g_4$.



Theorem 5. TL and Kas, defined on X and \bar{X} as before, extend to planar algebra morphisms $\{\text{tangles}\} \rightarrow \{S\}$.



Proof of Theorem 1'. Fix W and consider triples $(V, S, \phi: V \rightarrow W)$ where $S = (s, D, Q)$ is an SPQ on V . Say that two triples are "push-equivalent", $(V_1, S_1, \phi_1) \sim (V_2, S_2, \phi_2)$ if for every quadratic U on W ,

$$\sigma_{V_1}(S_1 + \phi_1^*U) = \sigma_{V_2}(S_2 + \phi_2^*U).$$

Given our (V, S, ϕ) , we need to show:

1. There is an SPQ S' on W such that $(V, S, \phi) \sim (W, S', I)$.
2. If $(W, S', I) \sim (W, S'', I)$ then $S' = S''$.

Property 2 is easy (Exercises 1, 2). Property 1 follows from the following three claims, each of which is easy.

Claim 1. If $v \in \ker \phi \cap D(S)$, and $\lambda := Q(v) \neq 0$, then $(V, S, \phi) \sim (V/\langle v \rangle, (s + \text{sign}(\lambda), V/\langle v \rangle, Q - \frac{Q(-, v) \otimes Q(v, -)}{|\lambda|^2}), \phi/\langle v \rangle)$.

So wlog $Q|_{\ker \phi} = 0$ (meaning, $Q|_{\ker \phi \otimes \ker \phi} = 0$). \square

Claim 2. If $Q|_{\ker \phi} = 0$ and $v \in \ker \phi \cap D(S)$, let $V' = \ker Q(v, -)$ and then $(V, S, \phi) \sim (V', S|_{V'}, \phi|_{V'})$ so wlog $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$. \square

Claim 3. If $Q|_{V \otimes \ker \phi + \ker \phi \otimes V} = 0$ then $S = \phi^*S'$ for some SPQ S' on $\text{im } \phi$ and then $(V, S, \phi) \sim (W, S', I)$. $\square \square$

Proof of Theorem 2. The functoriality of pullbacks needs no proof. Now assume $V_0 \xrightarrow{\alpha} V_1 \xrightarrow{\beta} V_2$ and that S is an SPQ on V_0 . Then for every SPQ U on V_2 we have, using reciprocity three times, that $\sigma(\beta_*\alpha_*S + U) = \sigma(\alpha_*S + \beta^*U) = \sigma(S + \alpha^*\beta^*U) = \sigma(S + (\beta\alpha)^*U) = \sigma((\beta\alpha)_*S + U)$. Hence $\beta_*\alpha_*S = (\beta\alpha)_*S$. \square

Lemma 1. $\phi_*\phi^*S = S|_{\text{im } \phi}$.

Proof. For every PQ U with $D(U) = \text{im } \phi$ we have $\sigma(S|_{\text{im } \phi} + U) = \sigma(S + U) = \sigma(\phi^*(S + U)) = \sigma(\phi^*S + \phi^*U) = \sigma(\phi_*\phi^*S + U)$ where for the second equality we use the fact that a pullback by a surjective map does not change the signature, and the last equality is the reciprocity property. \square

Lemma 2. Under the conditions of Theorem 3, if S_i is an SPQ on V_i for $i = 1, 2$, then $\sigma(\alpha^*S_1 + \gamma^*S_2) = \sigma(\beta_*S_1 + \delta_*S_2)$.

Proof. Let $\pi := \beta\alpha = \delta\gamma$. Then (in order) by reciprocity with $U = 0$, by the functoriality of pushforwards, by Lemma 1, and using the surjectivity of α and of γ , $\sigma(\alpha^*S_1 + \gamma^*S_2) = \sigma(\pi_*(\alpha^*S_1 + \gamma^*S_2)) = \sigma(\beta_*\alpha_*\alpha^*S_1 + \delta_*\gamma_*\gamma^*S_2) = \sigma(\beta_*(S_1|_{\text{im } \alpha}) + \delta_*(S_2|_{\text{im } \gamma})) = \sigma(\beta_*S_1 + \delta_*S_2)$. \square

Proof of Theorem 3. Given S on V_2 , for every U on V_1 we have using reciprocity, Lemma 2, and reciprocity again, that $\sigma(\alpha_*\gamma^*S + U) = \sigma(\gamma^*S + \alpha^*U) = \sigma(\delta_*S + \beta_*U) = \sigma(\beta^*\delta_*S + U)$. Hence $\alpha_*\gamma^*S = \beta^*\delta_*S$. \square

Homework.

Exercise 1. Show that if two SPQ's S_1 and S_2 on V satisfy $\sigma(S_1 + U) = \sigma(S_2 + U)$ for every quadratic U on V , then they have the same shift and the same domains.

Exercise 2. Show that if two full quadratics Q_1 and Q_2 satisfy $\sigma(Q_1 + U) = \sigma(Q_2 + U)$ for every U , then $Q_1 = Q_2$.

Exercise 3. By taking $U = 0$ in the reciprocity statement, prove that always $\sigma(\phi_*S) = \sigma(S)$. But that seems wrong, if $\phi = 0$. What saves the day?

Exercise 4. By taking $S = 0$ in the reciprocity statement, prove that always $\sigma(\phi^*U) = \sigma(U)$. But wait, this is nonsense! What went wrong?

Exercise 5. There are 11 types or irreducible commutative squa-

res: $1 \succ 0, 0 \succ 1, 0 \succ 0, 0 \succ 0, 1 \succ 1, 0 \succ 1, 0 \succ 1,$
 $\downarrow \downarrow$
 $0 \succ 0 \quad 0 \succ 0 \quad 1 \succ 0 \quad 0 \succ 1 \quad 0 \succ 0 \quad 0 \succ 1 \quad 0 \succ 1$
 $0 \succ 0, 0 \succ 1, 1 \succ 1,$ and $1 \succ 1$. Show that pushing commutes
 $\downarrow \downarrow$
 $1 \succ 1 \quad 1 \succ 1 \quad 1 \succ 0 \quad 1 \succ 1$
 with pulling for all but \checkmark of them. Compare with the statement of Theorem 3.

Solutions / Hints.

Hint 1. On a vector in the domain of one but not the other, take an outrageous value for U , that will raise or lower the signature.

Hint 2. WLOG, Q_1 is diagonal and $Q_1 = 0$.


Hint 3. The "shift" part of 0_*S is $\sigma(S)$.


Hint 4. ϕ_*S isn't 0, it's the partial quadratic "0 on $\text{im } \phi$ " (and indeed, $\sigma(\phi^*U) = \sigma(U)$ if ϕ is surjective).

Hint 5. The exceptions are ${}_{11}^{01}$ and ${}_{10}^{11}$. \checkmark \rightarrow $01, 10$

Add to The HW section:

1. Any lemma that's no longer needed.

2.  analyze diagrams of this form from the perspective of the [push, pull] theorem.

3. Does [push, pull] hold for  4. Understand from the perspective of the list of the proof box.


$S_i \xrightarrow{\phi} S_e$
 $\downarrow \psi \quad \downarrow \tau \text{ injective}$
 $\Rightarrow \phi_*\psi^*S = \tau^*S$

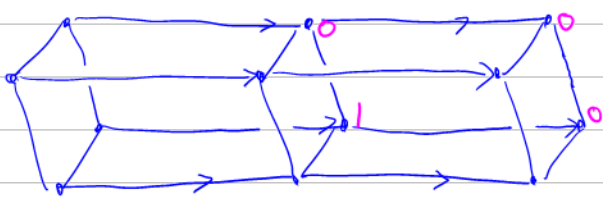
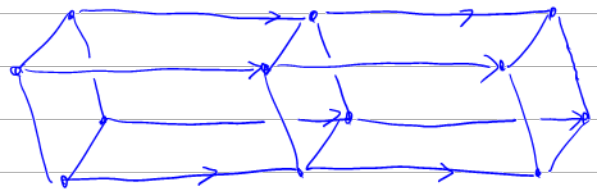
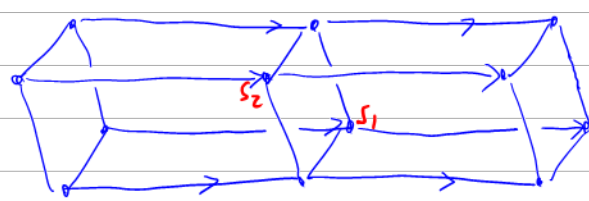
6. Pull back of pushforward scores are pushforward scores.

7_0^0 claim: $\pi_*\alpha^*S = \beta_*S$ if $\beta\alpha = \pi$
 $\text{proof } \sigma(\pi_*\alpha^*S + U) = \sigma(\alpha^*S + \pi^*U) = \sigma(\alpha^*S + \alpha^*\beta^*U) = \sigma(S + \beta^*U) = \sigma(\beta_*S + U)$ for every U , and so $\beta_*S = \pi_*\alpha^*S$

8_0^0 claim: $S \circ \alpha \rightarrow 0$ If α is 1-1, then $\alpha^*\alpha^*S = S$

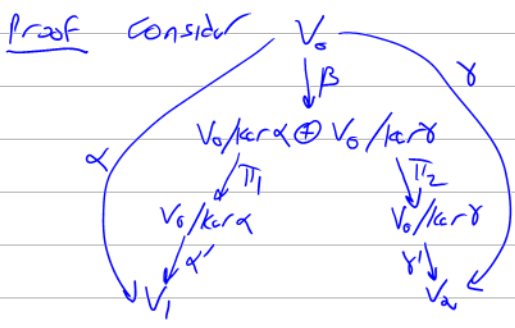
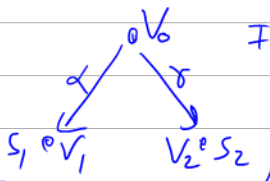
The classification of irreducible commutative squares $\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow g & & \downarrow h \\ Z & \xrightarrow{k} & W \end{array}$:

	$\begin{array}{c} 1 \\ 0 \\ 0 \end{array}$	$\begin{array}{c} 0 \\ 0 \\ 1 \end{array}$	$\begin{array}{c} 0 \\ 1 \\ 0 \end{array}$	$\begin{array}{c} 0 \\ 0 \\ 1 \end{array}$	$\begin{array}{c} 1 \\ 0 \\ 0 \end{array}$	$\begin{array}{c} 0 \\ 1 \\ 1 \end{array}$	$\begin{array}{c} 1 \\ 0 \\ 0 \end{array}$	$\begin{array}{c} 0 \\ 0 \\ 1 \end{array}$	$\begin{array}{c} 1 \\ 1 \\ 0 \end{array}$	$\begin{array}{c} 0 \\ 1 \\ 1 \end{array}$	$\begin{array}{c} 1 \\ 1 \\ 1 \end{array}$
[push, pull]:	✓	✗	✗	✓	✓	✓	✓	✓	✗	✗	✓
Equilizers?	✗	✗	✗	✓	✓	✓	✓	✓	✗	✗	✓
$V \leftarrow V \oplus A$ $W \leftarrow W \oplus A$	✗	✗	✗	✗	✓	✓	✗	✓	✗	✗	✓
$V \oplus A \xrightarrow{\pi_1} V$ $W \oplus A \xrightarrow{\pi_1} W$	✗	✗	✗	✗	✓	✓	✓	✗	✗	✗	✓
$A \oplus B \oplus C \rightarrow A \oplus B$ $A \oplus C \rightarrow A$	✗	✗	✗	✗	✓	✗	✓	✗	✗	✗	✓
 scene	✗	✗	✗	✗	✓	✗	✓	✗	✗	✗	✓
$\ker \alpha \cap \ker \beta = 0$	✗	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
$(\text{im } \alpha > \ker \beta) \wedge (\text{im } \beta > \ker \alpha)$	✓	✗	✗	✓	✓	✓	✓	✓	✓	✓	✓
$\ker T = \ker \alpha + \ker \beta$	✓	✓	✓	✓	✓	✓	✓	✓	✗	✓	✓
$\text{im } T = \text{im } \alpha \wedge \text{im } \beta$	✓	✓	✓	✓	✓	✓	✓	✓	✓	✗	✓
$A \oplus D \oplus E \rightarrow A \oplus B$ $A \oplus C \oplus D \oplus F \rightarrow A \oplus C$	✓		✓		✓	✓	✓				✓



It would be nice to have some restricted additivity for pushforwards.
 Perhaps first, additivity for signatures.

claim $\text{IF } \text{Ker } \alpha + \text{Ker } \delta = V_0$, Then $\sigma(\alpha^*S_1 + \delta^*S_2) = \sigma(\alpha^*S_1) + \sigma(\delta^*S_2)$



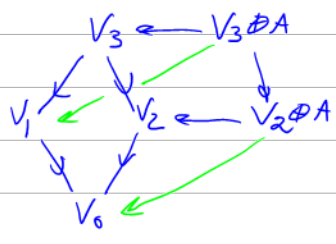
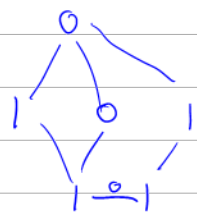
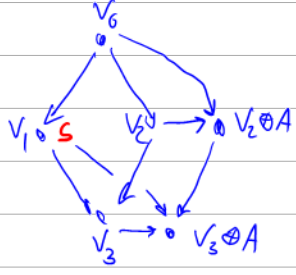
β is a surjection for if $(\bar{u}_1, \bar{u}_2) \in V_0/\text{ker } \alpha \oplus V_0/\text{ker } \delta$, write $u_1 - u_2 = w_2 - w_1$ w/ $w_1 \in \text{ker } \alpha$, $w_2 \in \text{ker } \delta$, and then $\beta(u_1 + w_1) = \bar{u}_1 + w_1 \oplus \bar{u}_2 + w_2 = \bar{u}_1 + \bar{u}_2$. β is a surjection.

Hence $\sigma(\alpha^*S_1 + \delta^*S_2) = \sigma(\beta^*(\pi_1^* \alpha^* S_1 + \pi_2^* \delta^* S_2)) \stackrel{\beta \text{ is surjection}}{=} \sigma(\alpha^*S_1 + \delta^*S_2)$

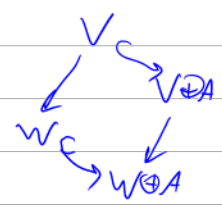
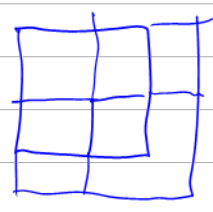
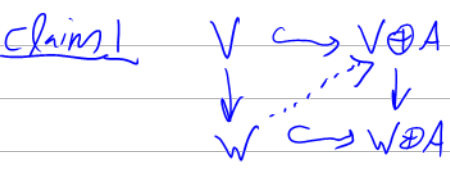
$$\sigma(\alpha^*S_1 \oplus \delta^*S_2) = \sigma(\alpha^*S_1) + \sigma(\delta^*S_2) \stackrel{\text{The } \beta \text{ are surjections}}{=} \sigma(\beta^* \pi_1^* \alpha^* S_1) + \sigma(\beta^* \pi_2^* \delta^* S_2) = \sigma(\alpha^*S_1) + \sigma(\delta^*S_2) \quad \square$$

claim(?) suppose [push, pull] holds for a diagram . Then it still holds if you add:

$\begin{smallmatrix} 1 & 0 \\ 0 & 0 \end{smallmatrix} : \checkmark$ $\begin{smallmatrix} 0 & 0 \\ 0 & 1 \end{smallmatrix} : \checkmark$ $\begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix} : \times$ $\begin{smallmatrix} 1 & 0 \\ 0 & 0 \end{smallmatrix} : \times$ $\begin{smallmatrix} 0 & 1 \\ 1 & 0 \end{smallmatrix} : \checkmark$ (using claim 1) $\begin{smallmatrix} 1 & 0 \\ 0 & 1 \end{smallmatrix} : \checkmark$ (modulo claim 2)

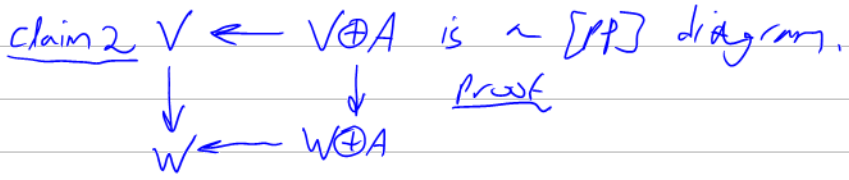


Is $(\phi + \psi)_* = \phi_* + \psi_*$? $\sigma((\phi + \psi)_* S + U) = \sigma(S + (\phi + \psi)^* U) = \sigma(S + \phi^* U + \psi^* U)$ [unlikely]

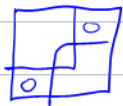


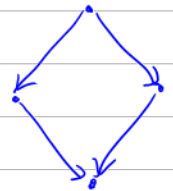
$$\begin{smallmatrix} 1 & 0 \\ 0 & 1 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 0 \\ 0 & 0 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 0 \\ 1 & 1 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 0 \\ 1 & 1 \end{smallmatrix}$$

is a [PP] diagram.
Proof: The indicated direction is easier than its opposite.



$$\begin{smallmatrix} 1 & 0 \\ 0 & 1 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 0 \\ 0 & 0 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 1 \\ 0 & 1 \end{smallmatrix} \quad \begin{smallmatrix} 1 & 0 \\ 1 & 1 \end{smallmatrix}$$

In the  context, the pushforward is additive. What's the abstraction?



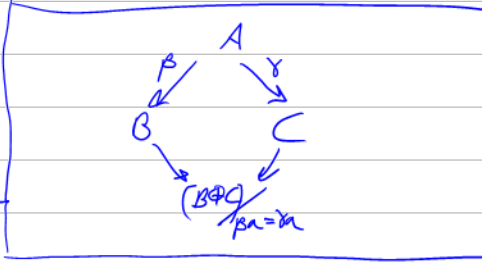
Claim(?) suppose $A_1, A_2 \in V$ $A_1 \cap A_2 = \{0\}$, $\pi_1: V \rightarrow V/A_1$, $\pi_2: V \rightarrow V/A_2$

S_i is an SRQ on V/A_i . Then $\pi_*(\pi_1^* S_1 + \pi_2^* S_2) = \pi_* \pi_1^* S_1 + \pi_* \pi_2^* S_2$.

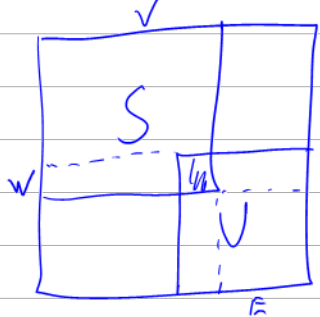
NTS, $\forall U$

$$\sigma(\pi_* \pi_1^* S_1 + \pi_* \pi_2^* S_2 + U) = \sigma(\pi_1^* S_1 + \pi_2^* S_2 + \pi^* U)$$

does not satisfy the conditions likewise



The stupidest extension case:



$$S \in V \leftarrow V \oplus E$$

$$\phi \downarrow \swarrow \downarrow \phi \oplus I$$

$$W \in W \leftarrow W \oplus E$$

Claim $\phi_* S = \pi_* \alpha^* S$

$$\text{PF } \sigma(\pi_* \alpha^* S + U) = \sigma(\alpha^* S + \pi^* U)$$

$$= \sigma(\alpha^* S + \alpha^* \phi^* U) = \sigma(\alpha^*(S + \phi^* U))$$

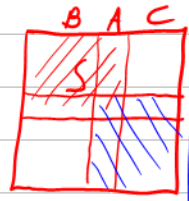
α is surjective $\Rightarrow \sigma(S + \phi^* U) = \sigma(\phi_* S + U)$

Q. What's the pushforward of a direct sum? Given $A_1 \oplus A_2 \xrightarrow{\phi = \phi_1 \oplus \phi_2} B$, is $\phi_*(S_1 \oplus S_2) = \phi_{1*}(S_1) + \phi_{2*}(S_2)$?

False already for $A \oplus B \xrightarrow{\text{id} \oplus 0} A \oplus B$ perhaps true if both ϕ_1 & ϕ_2 are surjective?

$$\sigma(\phi_{1*}(S_1) + \phi_{2*}(S_2) + U) =$$

HW. Do the pathetic case, $S \in A \oplus B \xleftarrow{\alpha} A \oplus B \oplus C$, using $\beta \downarrow$ and $\delta \downarrow$, only the reciprocal definition of pushforwards.



Lemma(?) If $S_1, S_2 \in A \oplus B$ w/ $B \subset \text{rad } S_1$ and $\forall U \in A$, $\sigma(S_1 + \pi_1^* U) = \sigma(S_2 + \pi_1^* U)$, then $S_1 = S_2$.

$$\sigma(\beta_* S + U) = \sigma(\beta_* S + \delta_* U) = \sigma(S + \beta^* \delta_* U)$$

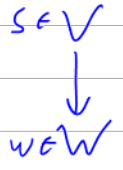
$$\sigma(\gamma_* \alpha^* S + U) =$$

Lemma(?) $C \subset \text{rad } \delta_* \alpha^* S$

Lemma(?). Given $\phi: V \rightarrow W$, if $w \in W$ is such that $\phi^{-1}(w) \subset \text{rad } S$, then $w \in \text{rad}(\phi_* S)$.

PF Assume $w \notin \text{rad}(\phi_* S)$.

Take U w/ $U(w, w) \neq 0$ and $U(w, \text{rad}(\phi_* S)) = 0$. Then

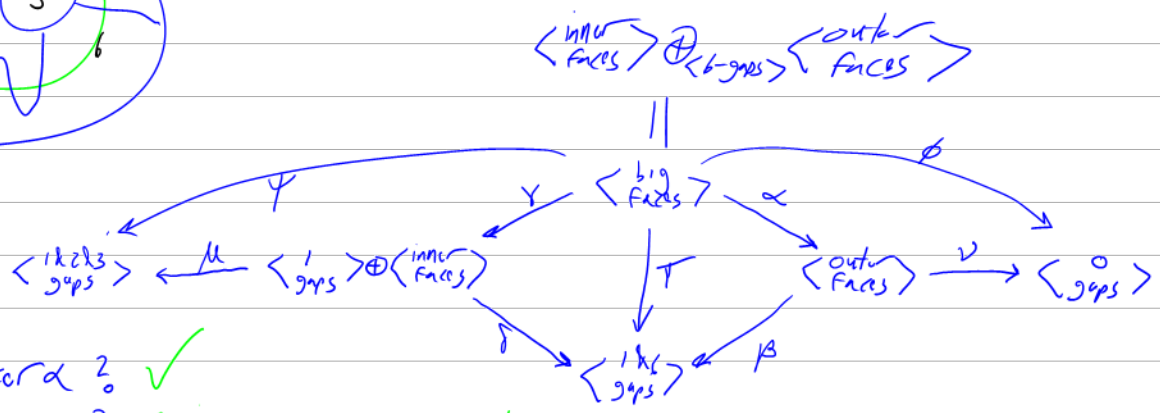
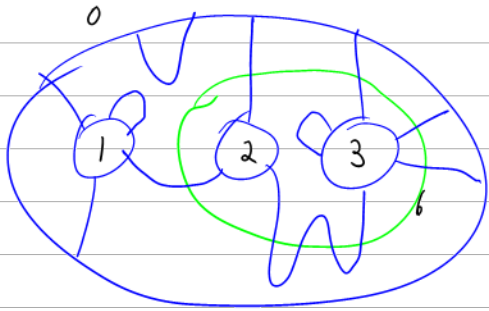


$$\sigma(\phi_* S + U) = \sigma(S + \phi^* U)$$

Lemma/Exercise. Given Q on V , let $\pi: V \rightarrow V/\text{rad } Q$. Then $\pi_* Q = \underline{Q/\text{rad } Q}$.
 PF Follows from the surjectivity of π and from $\pi^*(Q/\text{rad } Q) = Q$.
 w/ obvious definition

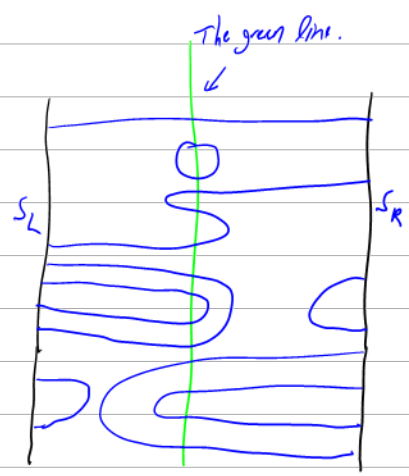
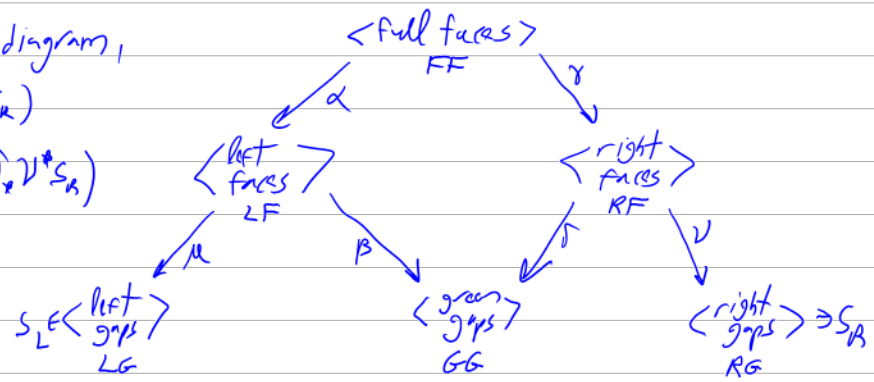
Indeed, $\sigma(Q/\text{rad } Q + U) = \sigma(\pi^*(Q/\text{rad } Q) + U) = \sigma(Q + \pi^* U)$

If ϕ is surjective, $\sigma(S+\phi^*U) = \sigma(\phi_*S+U) = \sigma(\phi^*\phi_*S+\phi^*U)$

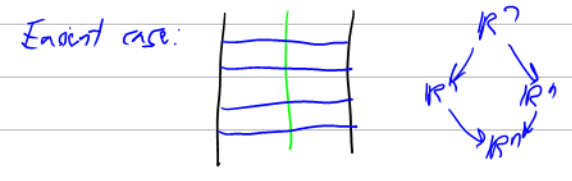


$\ker T = \ker \delta + \ker \alpha$? ✓
 $\text{Im } T = \text{Im } \delta \cap \text{Im } \beta$? looks good, needs confirmation.

Then in this diagram,
 $\sigma(\alpha^* \mu^* S_L + \beta^* \nu^* S_R)$
 $= \sigma(\beta^* \mu^* S_L + \delta^* \nu^* S_R)$



claim $\text{Im } \alpha + \ker \mu = \text{LF}$



Implementation (sources: <http://drorbn.net/icerm23/ap>). I like it most when the implementation matches the math perfectly. We failed here.

Once[<< KnotTheory`];

Loading KnotTheory` version

of February 2, 2020, 10:53:45.2097.

Read more at <http://katlas.org/wiki/KnotTheory>.

Utilities. The step function, algebraic numbers, canonical forms.

$\theta[x_]$ /; NumericQ[x] := UnitStep[x]

```
 $\omega 2[v\_][p\_]$  := Module[{q = Expand[p], n, c},
  If[q == 0, 0,
    c = Coefficient[q,  $\omega$ , n = Exponent[q,  $\omega$ ]];
     $c v^n + \omega 2[v][q - c(\omega + \omega^{-1})^n]$ ];
```

```
sign[ $\mathcal{E}$ _] := Module[{n, d, v, p, rs, e, k},
  {n, d} = NumeratorDenominator[ $\mathcal{E}$ ];
  {n, d} /=  $\omega^{\text{Exponent}[n, \omega]/2 + \text{Exponent}[n, \omega, \text{Min}]/2}$ ;
  p = Factor[ $\omega 2[v]@n * \omega 2[v]@d / . v \rightarrow 4 u^2 - 2$ ];
  rs = Solve[p == 0, u, Reals];
  If[rs == {}, Sign[p /. u -> 0],
    rs = Union@{u /. rs};
    Sign[(-1)e=Exponent[p, u] Coefficient[p, u, e]] + Sum[
      k = 0;
      While[{d = RootReduce[ $\partial_{\{u, ++k\}} p / . u \rightarrow r$ ]} == 0];
      If[EvenQ[k], 0, 2 Sign[d]] *  $\theta[u - r]$ ,
      {r, rs}]]
  ]
]
```

SetAttributes[B, Orderless];

```
CF[b_B] := RotateLeft[#, First@Ordering[#] - 1] & /@
DeleteCases[b, {}]
```

```
CF[ $\mathcal{E}$ _] := Module[{ $\gamma$ s = Union@Cases[ $\mathcal{E}$ ,  $\gamma$ _ |  $\bar{\gamma}$ _,  $\infty$ ]},
  Total[CoefficientRules[ $\mathcal{E}$ ,  $\gamma$ s] /.
    (ps_ -> c_) := Factor[c]  $\times$  Times@@ $\gamma$ sps]]
```

CF[{}] = {};

CF[C_List] :=

```
Module[{ $\gamma$ s = Union@Cases[C,  $\gamma$ _,  $\infty$ ],  $\gamma$ },
  CF /@ DeleteCases[0] [
    RowReduce[Table[ $\partial_{\gamma} r$ , {r, C}, { $\gamma$ ,  $\gamma$ s}]] .  $\gamma$ s ]
```

(\mathcal{E} _)* := $\mathcal{E} / . \{\bar{\gamma} \rightarrow \gamma, \gamma \rightarrow \bar{\gamma}, \omega \rightarrow \omega^{-1}, c_Complex \rightarrow c^*\}$;

r_Rule* := {r, r*}

RulesOf[γ_i + rest_] := ($\gamma_i \rightarrow -rest$)*;

```
CF[PQ[C_, q_]] := Module[{nC = CF[C]},
  PQ[nC, CF[q /. Union@@RulesOf /@nC]] ]
```

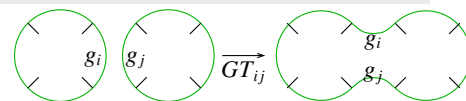
CF[Σ_b [σ _, pq_]] := $\Sigma_{CF[b]}$ [σ , CF[pq]]

Pretty-Printing.

```
Format[ $\Sigma_{b,B}[\sigma_, PQ[C_, q_]]$ ] := Module[{ $\gamma$ s},
   $\gamma$ s =  $\gamma$ # & /@ Join@@b;
  Column[{TraditionalForm@ $\sigma$ ,
    TableForm[Join[
      Prepend[""] /@ Table[TraditionalForm[ $\partial_c r$ ],
        {r, C}, {c,  $\gamma$ s}],
      {Prepend[""] [
        Join@@
          (b /. {L_, m___, r_} =>
            {DisplayForm@RowBox[{"(", L}],
              m, DisplayForm@RowBox[{r, ")"}]}) / .
            i_Integer =>  $\gamma_i$  ]},
      MapThread[Prepend,
        {Table[TraditionalForm[ $\partial_{r,c} q$ ], {r,  $\gamma$ s*},
          {c,  $\gamma$ s}],  $\gamma$ s*}]]
    ], TableAlignments -> Center]
  ], Center] ];
```

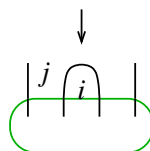
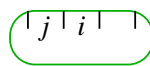
The Face-Centric Core.

```
 $\Sigma_{b1}[\sigma_1, PQ[C1_, q1_]] \oplus \Sigma_{b2}[\sigma_2, PQ[C2_, q2_]] \wedge :=$ 
  CF@ $\Sigma_{\text{Join}[b1, b2]}$ [ $\sigma_1 + \sigma_2, PQ[C1 \cup C2, q1 + q2]$ ];
```



GT for Gap Touch:

```
GTi,j@ $\Sigma_B$ [{li___, i_, ri___}, {lj___, j_, rj___}, bs___][ $\sigma$ _,
  PQ[C_, q_]] :=
  CF@ $\Sigma_B$ [{ri, li, j, rj, lj, i}, bs][ $\sigma$ _, PQ[C  $\cup$  { $\gamma_i - \gamma_j$ }, q]]
```



cor·don (kôr'dn)
n.



1. A line of people, military posts, or ships stationed around an area to enclose or guard it: a *police cordon*.
2. A rope, line, tape, or similar border stretched around an area, usually by the police, indicating that access is restricted.

$$s \begin{pmatrix} 0 & \phi C_{\text{rest}} \\ \bar{\phi}^T & \lambda \theta \\ \bar{C}_{\text{rest}}^T & \bar{\theta}^T A_{\text{rest}} \end{pmatrix} \rightarrow \begin{cases} \exists p \phi_p \neq 0 & \text{use } \phi_p \text{ to kill its row and} \\ \phi = 0, \lambda \neq 0 & \text{column, drop a } \begin{pmatrix} 01 \\ 10 \end{pmatrix} \text{ summand} \\ \phi = 0, \lambda = 0 & \text{use } \lambda \text{ to kill } \theta, \text{ let } s += \text{sign}(\lambda) \\ & \text{append } \theta \text{ to } C_{\text{rest}}. \end{cases}$$

```
Cordoni@ $\Sigma_B$ [{li___, i_, ri___}, bs___][ $\sigma$ _, PQ[C_, q_]] :=
```

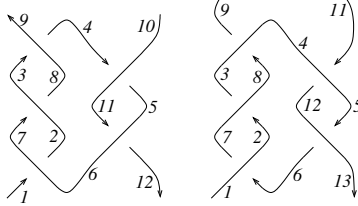
```
Module[{ $\phi = \partial_{\gamma_i} C$ ,  $\lambda = \partial_{\bar{\gamma}_i, \gamma_i} q$ , n $\sigma = \sigma$ , nC, nq, p},
  {p} = FirstPosition[ (# != 0) & /@  $\phi$ , True, {0}];
  {nC, nq} = Which[
    p > 0, {C, q} /. ( $\gamma_i \rightarrow -C[[p]] / \phi[[p]]$ )* /. ( $\gamma_i \rightarrow \theta$ )*,
     $\lambda \neq 0$ , (n $\sigma += \text{sign}[\lambda]$ ;
      {C, q /. ( $\gamma_i \rightarrow -(\partial_{\bar{\gamma}_i} q) / \lambda$ )* /. ( $\gamma_i \rightarrow \theta$ )*}),
     $\lambda == 0$ , {C  $\cup$  { $\partial_{\bar{\gamma}_i} q$ }, q /. ( $\gamma_i \rightarrow \theta$ )*}];
  CF@ $\Sigma_B$ [Most@{ri, li}, bs][n $\sigma$ ,
    PQ[nC, nq] /. ( $\gamma_{\text{Last}@{ri, li}} \rightarrow \gamma_{\text{First}@{ri, li}}$ )* ] ]
```


The Conway-Kinoshita-Terasaka Tangles.



$$T1 = PD[\bar{X}_{-6,2,7,-1}, \bar{X}_{-2,8,3,-7}, \bar{X}_{-8,4,9,-3}, X_{-11,6,12,-5}, X_{-4,11,5,-10}];$$

$$T2 = PD[X_{-6,2,7,-1}, X_{-2,8,3,-7}, X_{-8,4,9,-3}, \bar{X}_{-12,6,13,-5}, \bar{X}_{-4,12,5,-11}, \bar{X}_{-10,15,11,-14}, \bar{X}_{-15,10,16,-9}];$$



Column@{TL [T1], Kas [T1]}

$$-2\theta(u - \frac{\sqrt{3}}{2}) + 2\theta(u + \frac{\sqrt{3}}{2}) - 1$$

(Y_{-10})	Y_9	Y_{-1}	Y_{12}
\bar{Y}_{-10}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega^2-\omega+1}$	$\frac{\omega-1}{\omega^2-\omega+1}$
\bar{Y}_9	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{12}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$

(Y_{-10})	Y_9	Y_{-1}	Y_{12}
\bar{Y}_{-10}	$2(u-1)(u+1)(4u^2-3)$	0	$-2(u-1)(u+1)(4u^2-3)$
\bar{Y}_9	0	$\frac{1}{2(4u^2-3)}$	0
\bar{Y}_{-1}	$-2(u-1)(u+1)(4u^2-3)$	0	$2(u-1)(u+1)(4u^2-3)$
\bar{Y}_{12}	0	$-\frac{1}{2(4u^2-3)}$	0

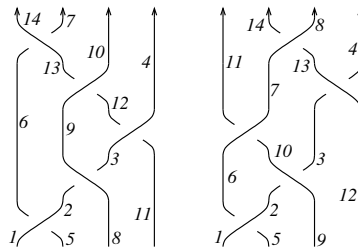
Column@{TL [T2], Kas [T2]}

(Y_{-14})	Y_{16}	Y_{-1}	Y_{13}
\bar{Y}_{-14}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{16}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{13}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$

Examples with non-trivial co-dimension.

$$B1 = PD[X_{-5,2,6,-1}, \bar{X}_{-8,3,9,-2}, X_{-11,4,12,-3}, X_{-12,10,13,-9}, \bar{X}_{-13,7,14,-6}];$$

$$B2 = PD[X_{-5,2,6,-1}, \bar{X}_{-9,3,10,-2}, X_{-10,7,11,-6}, \bar{X}_{-12,4,13,-3}, X_{-13,8,14,-7}];$$



Column@{TL [B1], Kas [B1]}

(Y_{-11})	Y_4	Y_{10}	Y_7	Y_{14}	Y_{-1}	Y_{-5}	Y_{-8}
\bar{Y}_{-11}	0	0	0	0	0	0	0
\bar{Y}_4	0	0	0	0	0	0	0
\bar{Y}_{10}	0	0	0	0	0	0	0
\bar{Y}_7	0	0	0	0	0	0	0
\bar{Y}_{14}	0	0	0	0	0	0	0
\bar{Y}_{-1}	0	0	0	0	0	0	0
\bar{Y}_{-5}	0	0	0	0	0	0	0
\bar{Y}_{-8}	0	0	0	0	0	0	0

Column@{TL [B2], Kas [B2]}

(Y_{-12})	Y_4	Y_8	Y_{14}	Y_{11}	Y_{-1}	Y_{-5}	Y_{-9}
\bar{Y}_{-12}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_4	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_8	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{14}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{11}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-5}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-9}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$

(Y_{-12})	Y_4	Y_8	Y_{14}	Y_{11}	Y_{-1}	Y_{-5}	Y_{-9}
\bar{Y}_{-12}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_4	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_8	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{14}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{11}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-1}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-5}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$
\bar{Y}_{-9}	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$	$\frac{\omega-1}{\omega}$

$\begin{pmatrix} A & B \\ C & U \end{pmatrix} \xrightarrow{\det(A)} \begin{pmatrix} I & A^{-1}B \\ C & U \end{pmatrix} \xrightarrow{1} \begin{pmatrix} I & A^{-1}B \\ 0 & U - CA^{-1}B \end{pmatrix}$. Roughly, $\det(A)$ is "det on ker", $-CA^{-1}B$ is "a pushforward of $\begin{pmatrix} A & B \\ C & U \end{pmatrix}$ ".

so $\det \begin{pmatrix} A & B \\ C & U \end{pmatrix} = \det(A) \det(U - CA^{-1}B)$. (what if $\mathbb{A}A^{-1}$?)

Questions. 1. Does this have a topological meaning? 2. Is there a version of the Kashaev Conjecture for tangles? 3. Find all solutions of R123 in our "algebra". 4. Braids and the Burau representation. 5. Recover the work in "Prior Art". 6. Are there any concordance properties? 7. What is the "SPQ group"? 8. The jumping points of signatures are the roots of the Alexander polynomial. Does this generalize to tangles? 9. Which of the three Cordon cases is the most common? 10. Are there interesting examples of tangles for which rels is non-trivial? 11. Is the pq part determined by Γ -calculus? 12. Is the pq part determined by finite type invariants? 13. Does it work with closed components / links? 14. Strand-doubling formulas? 15. A multivariable version? 16. Mutation invariance? 17. Ribbon knots? 18. Are there "face-virtual knots"? 19. Does the pushforward story extend to ranks? To formal Gaussian measures? To super Gaussian measures?

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Acknowledgement. This work was partially supported by NSERC grant RGPIN-2018-04350 and by the Chu Family Foundation (NYC).