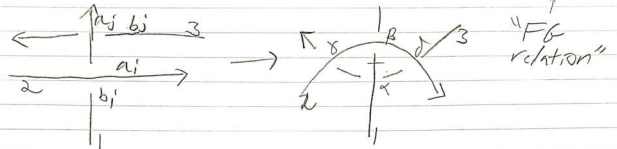




Requirement on  $A, B, R = \sum a_i \otimes b_i \in A \otimes B$



□

- $A$  &  $B$  must be  $\checkmark$  <sup>Fid.</sup> algebras.  $A^* \rightarrow B$  <sup>onto</sup>,  $B^* \rightarrow A$  <sup>onto</sup>
- The  $b_j$ 's must span all of  $B$   $\Rightarrow$   $R$  is of max rank,  $\dim A = \dim B$ ,  $\exists$  pairing. The  $a_j$ 's must span all of  $A$ .  $\Rightarrow \exists B^* \otimes A^*$  dual to  $R$ .

$$\sum_{i=1}^n \langle a_i, \cdot \rangle = \sum_{j=1}^n \langle \cdot, b_j \rangle = a \quad \sum \langle b_j, a_i \rangle = b$$

$$\begin{array}{c} \diagup \\ \diagdown \end{array} \rightarrow \sum a_i a_j \otimes b_i \otimes b_j \quad b \mapsto \sum \langle b, a_i a_j \rangle b_i \otimes b_j$$

$a$  co-product on  $B$   
s.t.  $(\Delta \otimes 1)R = R^{\otimes 2} R^3$

$$\begin{array}{c} \diagdown \\ \diagup \end{array} \rightarrow \sum a_i \otimes a_j \otimes b_j \quad a \mapsto \sum \langle a, a_j \rangle \langle b_j, \cdot \rangle = a$$

$a$  co-product on  $A$  s.t.  
 $(\Delta \otimes 1)R = R^{\otimes 2} R^3$

$\Rightarrow A, B$  are (dual) bi-algebras.

4.  $R^{-1} = \sum a'_i \otimes b'_i$   
 $R$  is invertible in  $A \otimes B$

$\Rightarrow$  s.t.  $(\Delta \otimes 1)R = R^{-1}$  ... likewise for  $B$ . an anti-pode for  $A \rightarrow$

The F&A relation:  
 $\sum b_i \cdot a_j \otimes a_i \otimes b_j = \sum a_\alpha b_\beta \otimes a_\gamma a_\delta \otimes b_\epsilon b_\zeta$

Pair w/  $b \otimes a$  on strands 2,3:  
 $\Delta(b) = b_1 \otimes b_2 \otimes b_3$   
 $\Delta(a) = a_1 \otimes a_2 \otimes a_3$   
 $b \cdot a := \sum a_\alpha b_\beta \otimes \langle b_1, a_\alpha \rangle \langle a_2, b_\beta \rangle \langle a_3, \cdot \rangle$   
 $\langle b_1, a_1 \rangle \langle b_2, a_2 \rangle \langle b_3, a_3 \rangle$   
 $= \langle b_1, a_3 \rangle \langle b_3, a_1 \rangle a_2 b_2$

Example  $A = U\langle x, y \rangle / [x, y] = x$   $B = U\langle b, c \rangle / [b, c] = -cy$

$\Delta(y, b, a, x) = y_1 + b_1 y_2$   $\Delta(b) = b_1 + b_2$   $\Delta(c) = c_1 + c_2$   
 $\Delta(x) = x_1 + a_1 x_2$   $\Delta(a) = a_1 + a_2$   $\Delta(x) = x_1 + a_1 x_2$   
 $S(y, b, a, x) = -b^{-1} y, -b_1, -a_1, -A^{-1} x$   
 $\Rightarrow \langle b^m, x^n \rangle = h^{nm} n! \langle y^n, x^m \rangle = h^{-nm} n! [y^n]_q!$   
 $\rho = e^{h\epsilon}$

$\Rightarrow R = \sum \frac{a^m b^n}{n!} \frac{x^m y^n}{[m]_q!}$   
 $ba = ab$   
 $xy = (yx + (1 - AB)/h)$

Conversion starts  
Impractical lecture on OV tangles and Drinfeld's Double Formula, and quantization of Lie bialgebras.  
Define OV tangles.  
Examples:  $X, Y = \begin{array}{c} \diagup \\ \diagdown \end{array}$

Thm Every tangle is an OV tangle.  
Proof  $\begin{array}{c} \text{---} \\ | \end{array} \rightarrow \begin{array}{c} \diagup \\ \diagdown \end{array}$

Corollary Knot theory is trivial!  
Too good to be true? What about  $\rho$ ?  
Too good to be ignored? Perhaps true in some quotient/completion/image?

- Quotient-completion: complete to virtual knots, divide by OC, get my paper "Knotted surfaces & hoops". There may be more...
- Completion like in  $\mathbb{Q} \hookrightarrow \mathbb{A}$ : see Ogden's talk.
- Completion like in  $\mathbb{Q}[x] \hookrightarrow \mathbb{Q}[[x]]$  or  $\mathbb{K} \hookrightarrow \widehat{\mathbb{K}} \cong \mathbb{K}^n$ .

I wish I knew

4i represent  $\begin{array}{c} \diagup \\ \diagdown \end{array}$   $R \in A \otimes B = \sum a_i \otimes b_i$   
 $b \cdot a := \sum a_i b_i$   
A "bald" formula.  
need a rule making  $A \otimes B$  into  $\mathbb{K}$  need  $b \cdot a := \sum a_i b_i$