

Pensieve header: The “speedy” engine.

## The “Speedy” Engine

Program

### Internal Utilities

Program

Canonical Form:

Program

```
In[*]:= CCF[ $\mathcal{E}$ _] := PP_CCF@ExpandDenominator@ExpandNumerator@PP_Together@Together[PP_Exp[
Expand[ $\mathcal{E}$ ] /. ex-ey->ex+y /. ex->eCCF[x]]];
CF[ $\mathcal{E}$ _List] := CF/@ $\mathcal{E}$ ;
CF[ $sd\_SeriesData$ ] := MapAt[CF,  $sd$ , 3];
CF[ $\mathcal{E}$ _] := PP_CF@Module[
{ $vs$  = Cases[ $\mathcal{E}$ , { $y$  |  $b$  |  $t$  |  $a$  |  $x$  |  $\eta$  |  $\beta$  |  $\tau$  |  $\alpha$  |  $\xi$ }_ ,  $\infty$ ] U { $y$ ,  $b$ ,  $t$ ,  $a$ ,  $x$ ,  $\eta$ ,  $\beta$ ,  $\tau$ ,  $\alpha$ ,  $\xi$ }},
Total[CoefficientRules[Expand[ $\mathcal{E}$ ],  $vs$ ] /. ( $ps\_ \rightarrow c\_$ ) -> CCF[ $c$ ] (Times@@ $vs^{ps}$ )]
];
CF[ $\mathcal{E}\_E$ ] := CF/@ $\mathcal{E}$ ; CF[Esp___[ $\mathcal{ES}$ ___]] := CF/@Esp[ $\mathcal{ES}$ ];
```

Program

The Kronecker  $\delta$ :

Program

```
In[*]:= K $\delta$  /: K $\delta$ i,j := If[i === j, 1, 0];
```

Program

Equality, multiplication, and degree-adjustment of perturbed Gaussians;  $E[L, Q, P]$  stands for  $e^{L+Q} P$ :

Program

```
In[*]:= E /: E[L1_, Q1_, P1_] == E[L2_, Q2_, P2_] :=
CF[L1 == L2] ^ CF[Q1 == Q2] ^ CF[Normal[P1 - P2] == 0];
E /: E[L1_, Q1_, P1_] E[L2_, Q2_, P2_] := E[L1 + L2, Q1 + Q2, P1 * P2];
E[L_, Q_, P_]$_k := E[L, Q, Series[Normal@P, { $\epsilon$ , 0, $k}]]];
```

Program

```
In[*]:= E3@E[ $\omega$ _, L_, Q_,  $Ps$ _] := CF/@E[L,  $\omega^{-1} Q$ ,  $\omega^{-1} (\omega^{-4} \epsilon)^{-1+Range@Length@Ps} . Ps]$_k;
E4@E[L_, Q_, P_] := Module[
{ $\omega$  = Normal[P]-1 /.  $\epsilon \rightarrow 0$ ,  $Ps$  = CoefficientList[P,  $\epsilon$ ]},
CF/@E[ $\omega$ , L,  $\omega Q$ ,  $\omega^{-3+4 Range@Length@Ps} Ps$ ]];
E3@Esp___[ $as$ ___] := E3@E[ $as$ ] /. E -> Esp;
E4@Esp___[ $as$ ___] := E4@E[ $as$ ] /. E -> Esp;$ 
```

Program

### Zip and Bind

Program

Variables and their duals:

Program

```
In[ ]:= {t*, b*, y*, a*, x*, z*} = {τ, β, η, α, ξ, ζ};
        {τ*, β*, η*, α*, ξ*, ζ*} = {t, b, y, a, x, z}; (u_{-i})* := (u*)_i;
```

Program

Finite Zips:

Program

```
In[ ]:= collect[sd_SeriesData, ζ_] := MapAt[collect[#, ζ] &, sd, 3];
        collect[ε_, ζ_] := PPCollect@Collect[ε, ζ];
        Zip[_][P_] := P;
        Zip_{ζs}_[Ps_List] := Zip_{ζs} /@ Ps;
        Zip_{ζs, ζs_...}[P_] := PPZip[
            (collect[P // Zip_{ζs}, ζ] /. f_ . ζ^{d_} .> ∂_{ζ^*, d} f) /. ζ* -> θ]
```

Program

QZip implements the “Q-level zips” on  $\mathbb{E}(L, Q, P) = e^{L+Q} P(\epsilon)$  and/or on  $\mathbb{E}(\omega, L, Q, P) = \omega^{-1} e^{L+\omega^{-1}Q} P(\omega^{-4} \epsilon)$ . Such zips regard the  $L$  variables as scalars.

$$\begin{aligned} \left\langle P(z_i, \zeta^j) e^{c+\eta^i z_i + y_j \zeta^j + q_i^j z_i \zeta^j} \right\rangle &= |\tilde{q}| \left\langle P(z_i, \zeta^j) e^{c+\eta^i z_i} \Big|_{z_i \rightarrow \tilde{q}_i^k(z_k + y_k)} \right\rangle \\ &= |\tilde{q}| e^{c+\eta^i \tilde{q}_i^k y_k} \left\langle P\left(\tilde{q}_i^k(z_k + y_k), \zeta^j + \eta^i \tilde{q}_i^j\right) \right\rangle. \end{aligned}$$

Program

```
$QZipFail = False;
QZip_{ζs_List}@E[L_, Q_, P_] := PPQZip@Module[{ζ, z, zs, c, ys, ηs, qt, zrule, ζrule, out},
    zs = Table[ζ*, {ζ, ζs}];
    c = CF[Q /. Alternatives @@ (ζs ∪ zs) -> θ];
    ys = CF@Table[∂_ζ (Q /. Alternatives @@ zs -> θ), {ζ, ζs}];
    ηs = CF@Table[∂_z (Q /. Alternatives @@ ζs -> θ), {z, zs}];
    qt = CF@Inverse@Table[Kδ_{z, ζ*} - ∂_{z, ζ} Q, {ζ, ζs}, {z, zs}];
    zrule = Thread[zs -> CF[qt.(zs + ys)]];
    ζrule = Thread[ζs -> ζs + ηs.qt];
    out = CF /@ E[L, c + ηs.qt.ys, Det[qt] Zip_{ζs}[P /. (zrule ∪ ζrule)]];
    If[¬ ($QZipFail ∨ TrueQ[out ≡ E3@QZip_{ζs}@E4@E[L, Q, P]]),
        $QZipFail = True; Print["QZip4 fail at {L,Q,P}=", {L, Q, P}];
    ];
    out
];
```

Program

```

In[ ]:= $QZipFail = False;
QZip $\zeta$ s_List@E[ $\omega$ _, L_, Q_, Ps_] := PPQZip4@Module[{ $\zeta$ , z, zs, c, ys,  $\eta$ s, qt, zrulerule,  $\zeta$ rulerule},
  zs = Table[ $\zeta^*$ , { $\zeta$ ,  $\zeta$ s}];
  c = CF[Q /. Alternatives@@( $\zeta$ s  $\cup$  zs)  $\rightarrow$  0];
  ys = CF@Table[ $\partial_{\zeta}$ (Q /. Alternatives@@zs  $\rightarrow$  0), { $\zeta$ ,  $\zeta$ s}];
   $\eta$ s = CF@Table[ $\partial_z$ (Q /. Alternatives@@ $\zeta$ s  $\rightarrow$  0), {z, zs}];
  qt = CF@Inverse@Table[K $\delta_{z, \zeta^*} - \partial_{z, \zeta} Q$ , { $\zeta$ ,  $\zeta$ s}, {z, zs}];
  zrulerule = Thread[zs  $\rightarrow$  CF[qt.(zs + ys)]];
   $\zeta$ rulerule = Thread[ $\zeta$ s  $\rightarrow$   $\zeta$ s +  $\eta$ s.qt];
  CF /@ E[ $\omega$  Det[qt /  $\omega$ ], L, c +  $\eta$ s.qt.yz, Zip $\zeta$ s[Ps /. (zrulerule  $\cup$   $\zeta$ rulerule)]]];

```

Program

Upper to lower and lower to Upper:

Program

```

In[ ]:= U21 = {B $\zeta$   $\rightarrow$  e $^{-p \hbar \gamma b_i}$ , B $\zeta^*$   $\rightarrow$  e $^{-p \hbar \gamma b}$ , T $\zeta$   $\rightarrow$  e $^{p \hbar t_i}$ , T $\zeta^*$   $\rightarrow$  e $^{p \hbar t}$ ,  $\mathcal{A}^{\zeta}$   $\rightarrow$  e $^{p \gamma \alpha_i}$ ,  $\mathcal{A}^{\zeta^*}$   $\rightarrow$  e $^{p \gamma \alpha}$ };
L2U = {e $^{c_- \cdot b_i + d_-}$   $\rightarrow$  B $\zeta^{c / (\hbar \gamma)}$  e $^d$ , e $^{c_- \cdot b + d_-}$   $\rightarrow$  B $^{-c / (\hbar \gamma)}$  e $^d$ ,
  e $^{c_- \cdot t_i + d_-}$   $\rightarrow$  T $\zeta^{c / \hbar}$  e $^d$ , e $^{c_- \cdot t + d_-}$   $\rightarrow$  T $^{c / \hbar}$  e $^d$ ,
  e $^{c_- \cdot \alpha_i + d_-}$   $\rightarrow$   $\mathcal{A}^{\zeta / \gamma}$  e $^d$ , e $^{c_- \cdot \alpha + d_-}$   $\rightarrow$   $\mathcal{A}^{c / \gamma}$  e $^d$ ,
  e $^{\zeta}$   $\rightarrow$  e $^{\text{Expand}[\zeta]}$ };

```

Program

LZip implements the “L-level zips” on  $\mathbb{E}(L, Q, P) = P e^{L+Q}$ . Such zips regard all of  $P e^Q$  as a single “P”. Here the z’s are  $b$  and  $\alpha$  and the  $\zeta$ ’s are  $\beta$  and  $a$ .

Program

```

In[ ]:= LZip $\zeta$ s_List@E[L_, Q_, P_] :=
  PPLZip@Module[{ $\zeta$ , z, zs, c, ys,  $\eta$ s, lt, zrulerule, Zrulerule,  $\zeta$ rulerule, Q1, EEQ, EQ},
    zs = Table[ $\zeta^*$ , { $\zeta$ ,  $\zeta$ s}];
    c = L /. Alternatives@@( $\zeta$ s  $\cup$  zs)  $\rightarrow$  0;
    ys = Table[ $\partial_{\zeta}$ (L /. Alternatives@@zs  $\rightarrow$  0), { $\zeta$ ,  $\zeta$ s}];
     $\eta$ s = Table[ $\partial_z$ (L /. Alternatives@@ $\zeta$ s  $\rightarrow$  0), {z, zs}];
    lt = Inverse@Table[K $\delta_{z, \zeta^*} - \partial_{z, \zeta} L$ , { $\zeta$ ,  $\zeta$ s}, {z, zs}];
    zrulerule = Thread[zs  $\rightarrow$  lt.(zs + ys)];
    Zrulerule = zrulerule /. r_Rule  $\rightarrow$ 
      ((U = r[[1]] /. {b  $\rightarrow$  B, t  $\rightarrow$  T,  $\alpha$   $\rightarrow$   $\mathcal{A}$ })  $\rightarrow$  (U /. U21 /. r // L2U)); (* not used *)
     $\zeta$ rulerule = Thread[ $\zeta$ s  $\rightarrow$   $\zeta$ s +  $\eta$ s.lt];
    Q1 = Q /. U21 /. (zrulerule  $\cup$   $\zeta$ rulerule);
    EEQ[ps___] := EEQ[ps] = PP $^{\text{EEQ}}$ @(CF[e $^{-Q1}$  D[e $^{Q1}$ , Sequence@@Thread[{zs, {ps}}]]] /.
      Alternatives@@zs  $\rightarrow$  0 // L2U);
    CF /@ ((CF/@*)E[
      c +  $\eta$ s.lt.yz, Q1 /. Alternatives@@zs  $\rightarrow$  0,
      Det[lt] (Zip $\zeta$ s[EQ@@zs] (P /. U21 /. (zrulerule  $\cup$   $\zeta$ rulerule)))] /.
      Derivative[ps___][EQ][___]  $\rightarrow$  EEQ[ps] /. _EQ  $\rightarrow$  1)
    ] // L2U)
  ];

```

Program

```

In[*]:=
B_{ } [L_, R_] := L R;
B_{is_} [L_E, R_E] := PP_B@Module[{n},
  Times[
    L /. Table[(v : b | B | t | T | a | x | y)_i -> v_{nei}, {i, {is}}],
    R /. Table[(v : beta | tau | alpha | A | xi | eta)_i -> v_{nei}, {i, {is}}]
  ] // LZipJoin@Table[{beta_{nei}, tau_{nei}, alpha_{nei}}, {i, {is}}] // QZipJoin@Table[{xi_{nei}, eta_{nei}}, {i, {is}}] ];
B_{is_} [L_, R_] := B_{is_} [L, R];

```

Program

## E morphisms with domain and range.

Program

```

In[*]:=
B_{is_List} [E_{d1 -> r1} [L1_, Q1_, P1_], E_{d2 -> r2} [L2_, Q2_, P2_]] :=
  E_{(d1 U Complement[d2, is]) -> (r2 U Complement[r1, is])} @@ B_{is} [E [L1, Q1, P1], E [L2, Q2, P2]];
E_{d1 -> r1} [L1_, Q1_, P1_] // E_{d2 -> r2} [L2_, Q2_, P2_] :=
  B_{r1 n d2} [E_{d1 -> r1} [L1, Q1, P1], E_{d2 -> r2} [L2, Q2, P2]];
E_{d1 -> r1} [L1_, Q1_, P1_] == E_{d2 -> r2} [L2_, Q2_, P2_] ^:=
  (d1 == d2) ^ (r1 == r2) ^ (E [L1, Q1, P1] == E [L2, Q2, P2]);
E_{d1 -> r1} [L1_, Q1_, P1_] E_{d2 -> r2} [L2_, Q2_, P2_] ^:=
  E_{(d1 U d2) -> (r1 U r2)} @@ (E [L1, Q1, P1] E [L2, Q2, P2]);
E_{d -> r} [L_, Q_, P_] $k_ := E_{d -> r} @@ E [L, Q, P] $k;
E_{S_} [i_] := {S} [[i]];

```