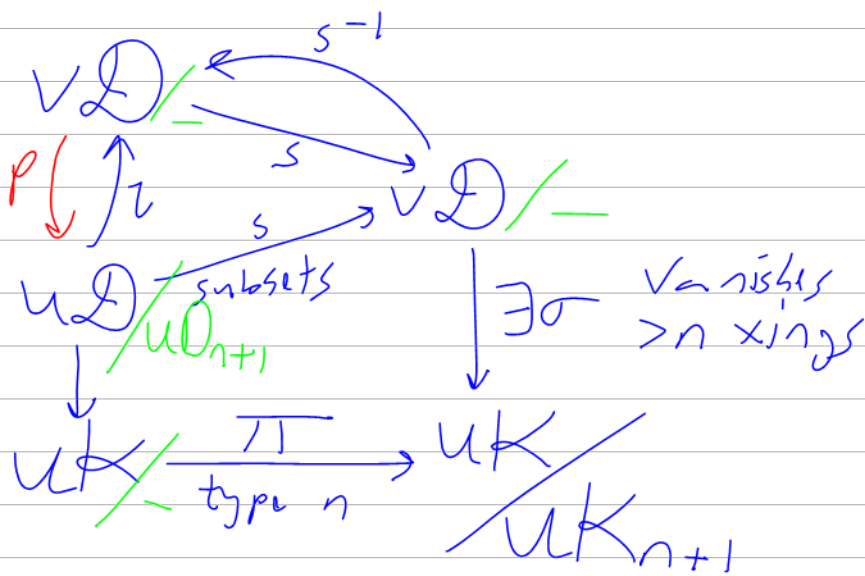
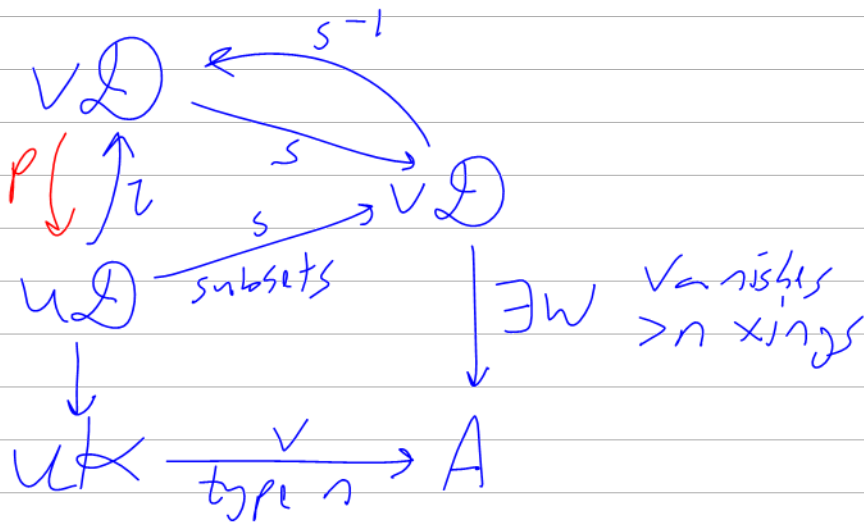
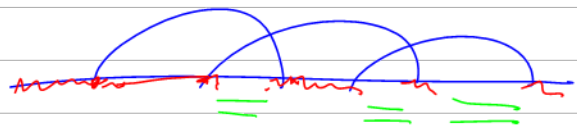
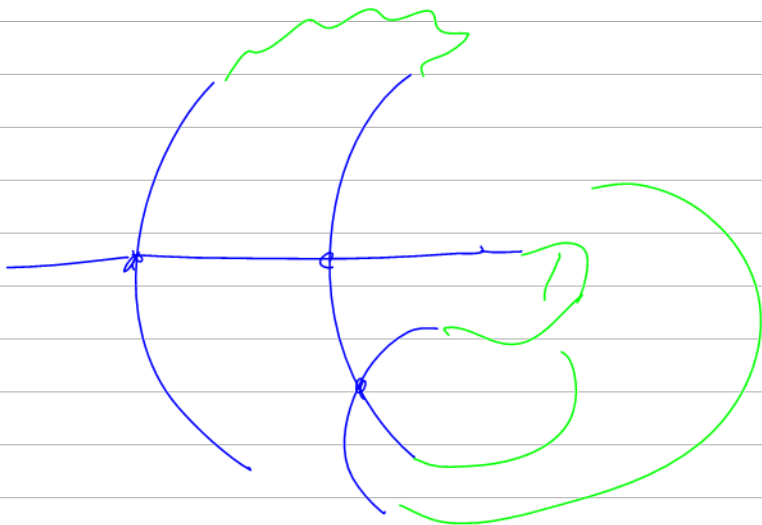


$u\mathcal{D}, v\mathcal{D}, u\mathcal{D}, v\mathcal{D} \dots$

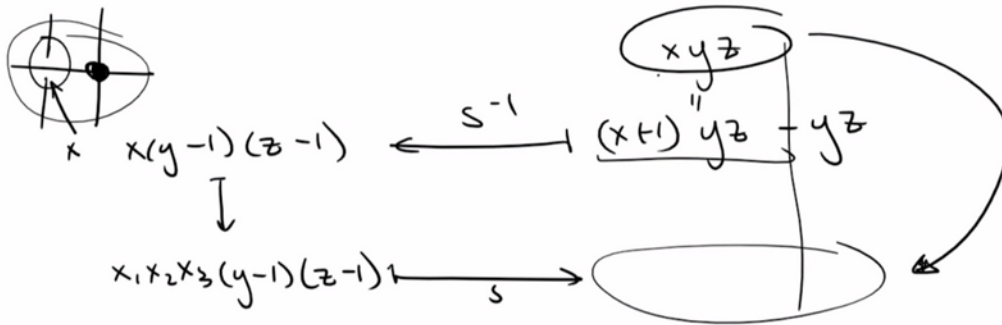
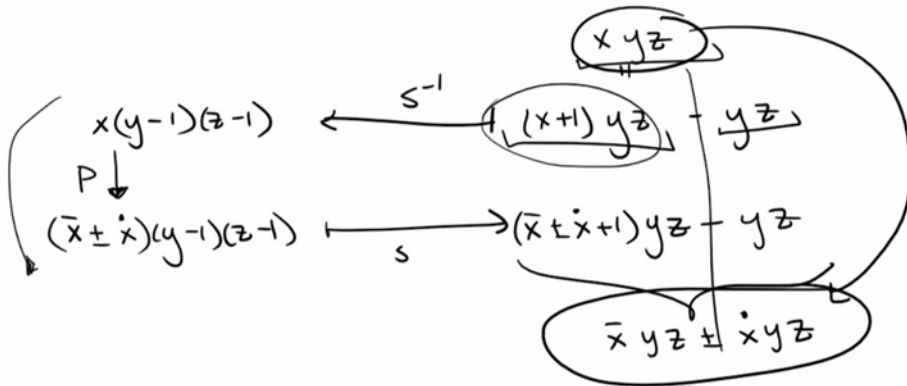
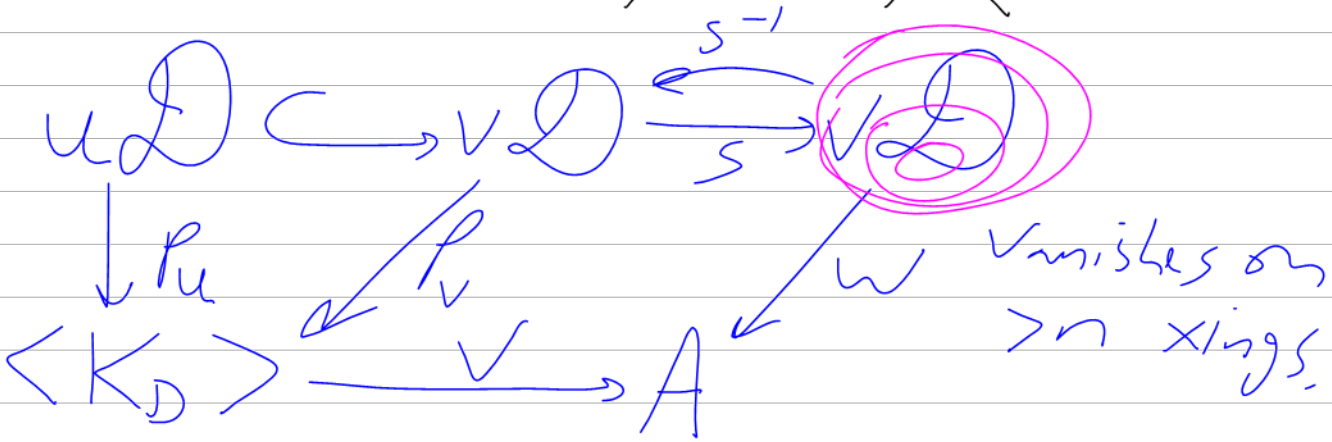


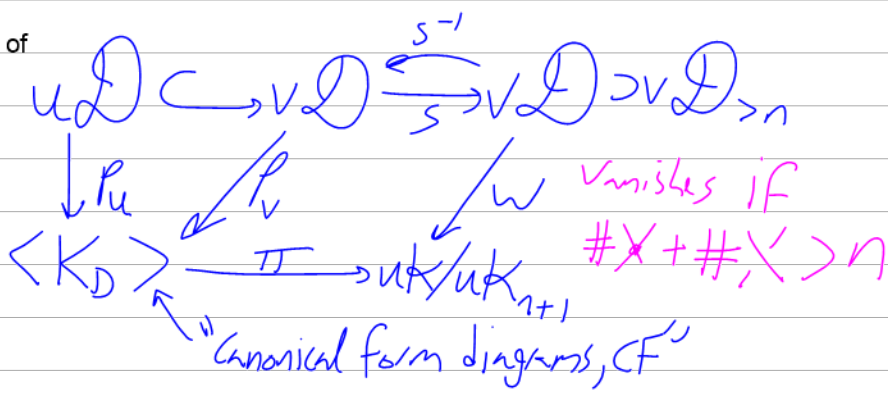
Does σ determine P
 which σ 's come
 from $\sim P$?



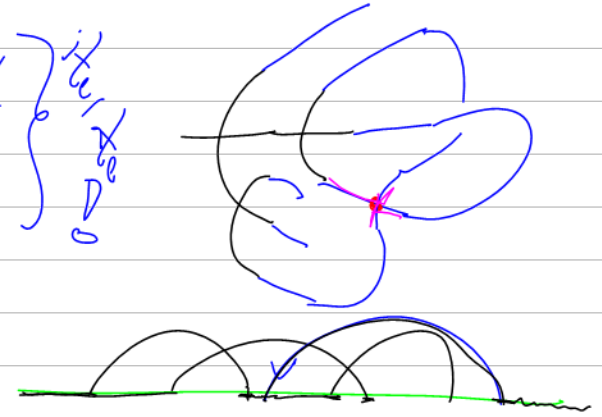
$$D \mapsto K_D$$

~~$$X = X$$~~





1. Brush xings off the maximal tree for \mathcal{D}
2. Flip the 1st non-dec xing, paying w/ X
3. Realize (V -xings) as descending.



NTS: $W = S^{-1} // P_v // \Pi$ vanishes on $v\mathcal{D}_{>n}$

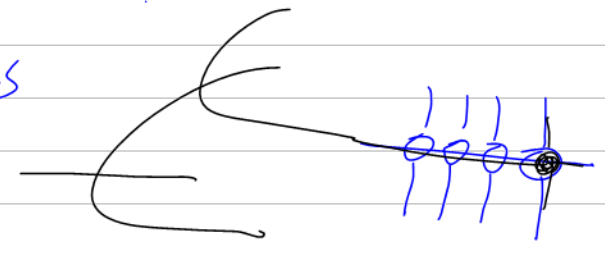
use downward induction on $\#X$.

$$\text{Xing} = \text{Xing} - \text{Xing}$$

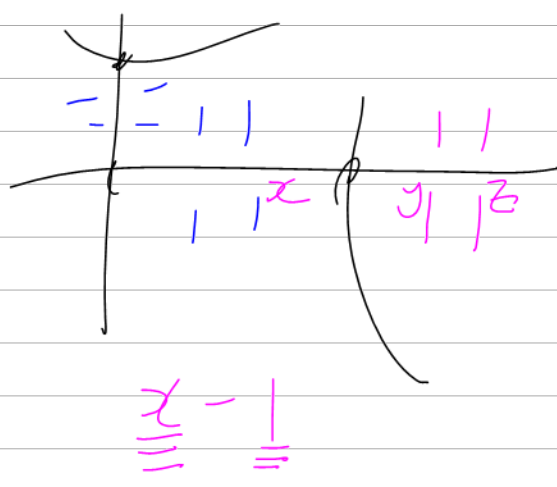
step 1 W vanishes on CF/VCF dings w/ $\#X' > 0$

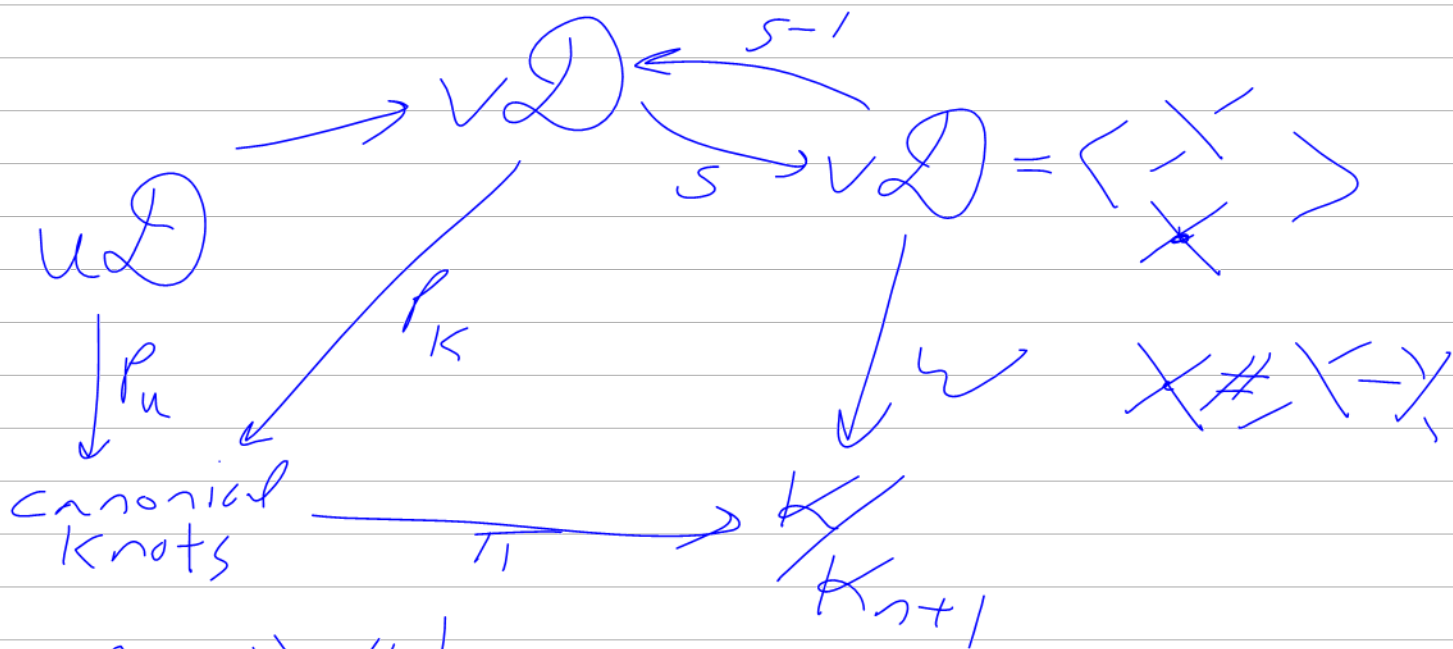
step 2 W vanishes on clear-tree diagrams

with a non-descending xing.



step 3 The rest.





$$P_K = \text{"} P \text{"} \mid \# \times \geq K$$

Assume P_{K+1} is def & has right props, def P_K

$$V \langle \text{diagram with slash and cross}, \text{diagram with slash and cross}, \text{diagram with slash and cross} \rangle \xrightarrow{P} U \langle \text{diagram with slash and cross}, \text{diagram with slash and cross} \rangle \mid \# \times > n$$

s.t. 1. $P \langle \text{diagram with slash and cross}, \text{diagram with slash and cross} \rangle$ should respect R-moves & $\times = \text{diagram with slash and cross} - \text{diagram with slash and cross}$

$$2. P \cdot \{ \# \times + \# \times > n \} \rightarrow \{ \# \times > n \}$$

Goal: Construct

$$\cancel{X} := \cancel{X} - X$$

$$P: vD \langle \cancel{X}, \cancel{X}' \rangle \rightarrow uD \langle \cancel{X}, \cancel{X}' \rangle / \# \cancel{X} > n$$

s.t. 1. $P|_{u0} = Id \pmod{R\text{-moves}}$ & $\cancel{X} = \cancel{X} - \cancel{X}'$

$$2. P|_{\# \cancel{X} + \# \cancel{X}' > n} = 0$$

IF $\# \cancel{X} > n$ in D $P(D) = 0$; Suppose P is defined & satisfies 1 & 2 IF $\# \cancel{X} > K$, and D has $\# \cancel{X} = K$.

* For every D choose a tree $T(D)$ (for \cancel{X}), index of \cancel{X} .



* Aside A "good choice function" on

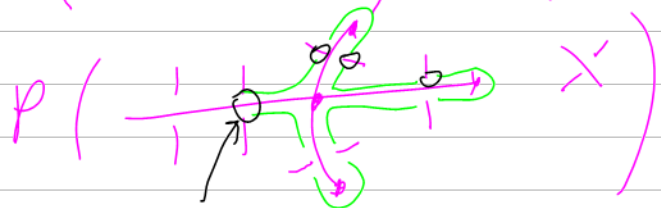
$S(X)$, the set of non-empty subsets of a set X , is $c: S(X) \rightarrow X$ s.t. 1. $c(A) \in A$

2. IF $b \in A \setminus \{c(A)\}$ then $c(A \setminus \{b\}) = c(A)$.

Thm {good choice Fncns} \leftrightarrow {total orderings of X }

* Define P by Fixing the First issue.

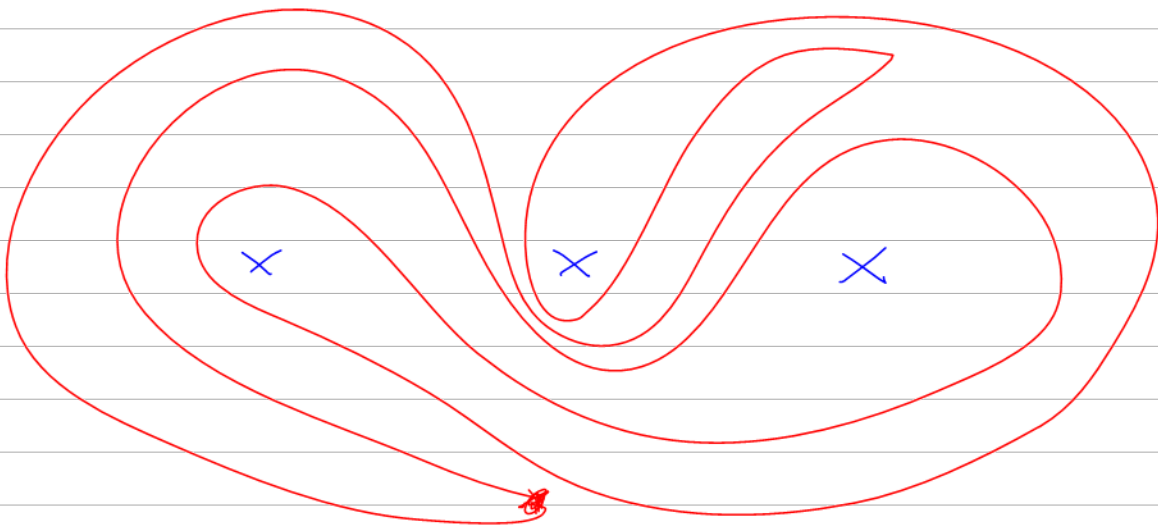
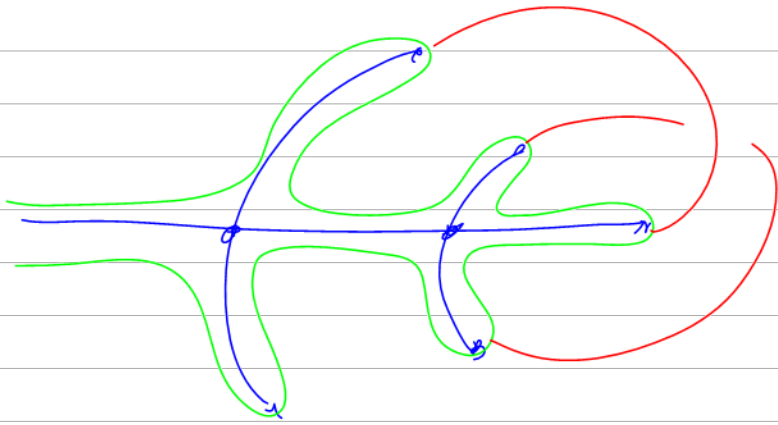
$$P(\dots \leftarrow) = P(\dots \cap \dots) = P(\dots \rightarrow)$$



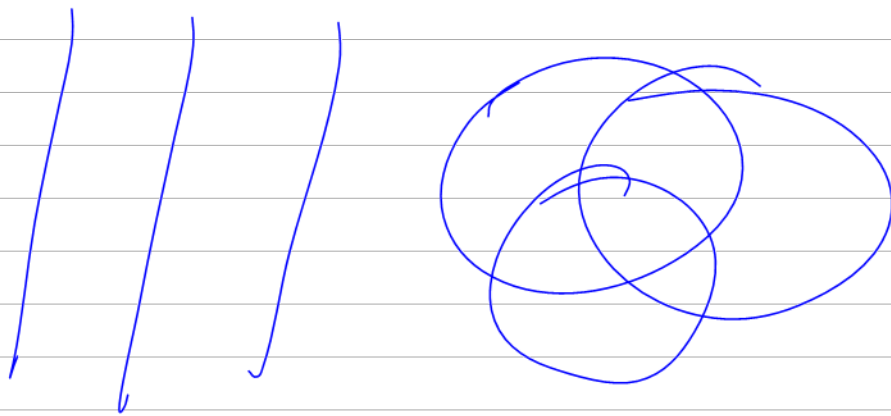
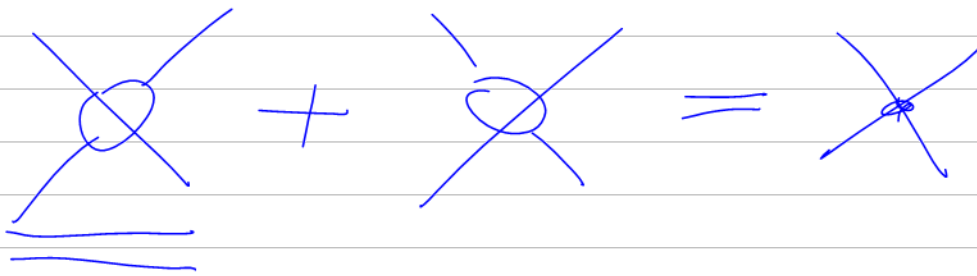
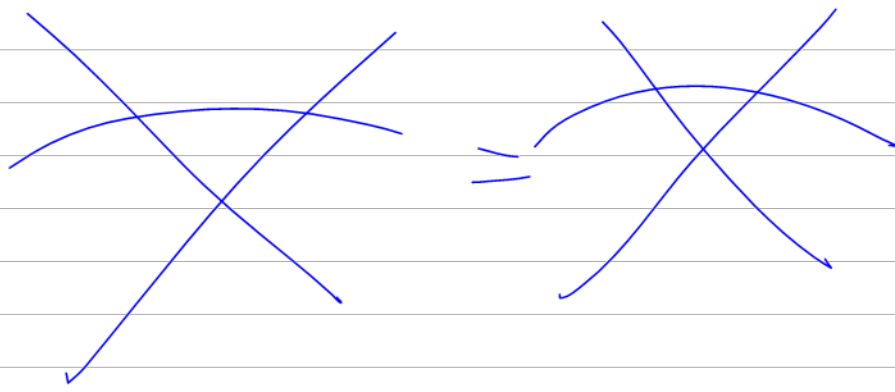
$$v(D) = v(PD) = v\left(\sum_{D' \subset D} D'|_{X \rightarrow \cancel{X}}\right) =$$

$$= v \left(P \sum_{\substack{D' \subset D \\ |D'| \leq n}} D' \mid_{X \rightarrow X} \right)$$

$$= \sum_{\substack{D' \subset D \\ |D'| \leq n}} v(P(D' \mid_{X \rightarrow X})) = \sum_{\sim} w(D')$$



$ab^{-1}cabb^{-1}c$



A poly of deg 3 on \mathbb{R}^{26} .

1. a, b, c, q, r, s .

2. $F: \mathbb{R}^{26} \rightarrow \mathbb{R}$ s.t. $\partial_{a,q,z,v} F = 0$

$$\mathbb{Z} \setminus W: \left\{ \begin{array}{l} \emptyset, \textcircled{+5} \\ \textcircled{\frac{5}{5}}, \textcircled{\frac{4}{4}}, \textcircled{\frac{3}{3}} \end{array} \right\} \rightarrow \textcircled{\times}$$

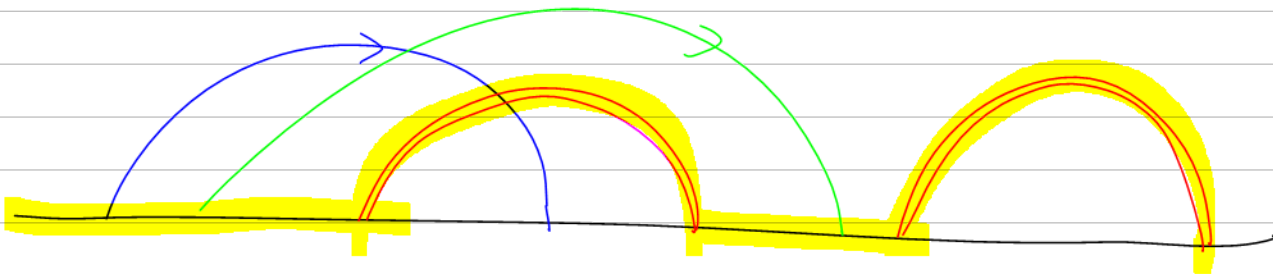
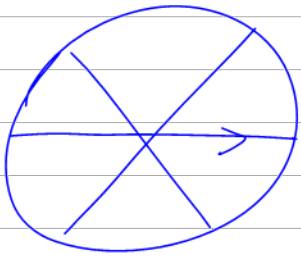
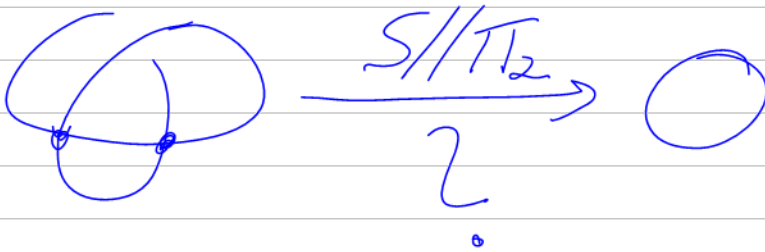
S.t.

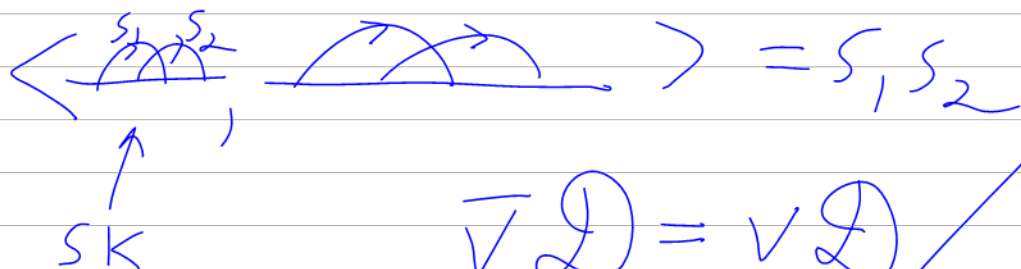
$$V_2(D) = W \circ S(D)$$

↑
knotting

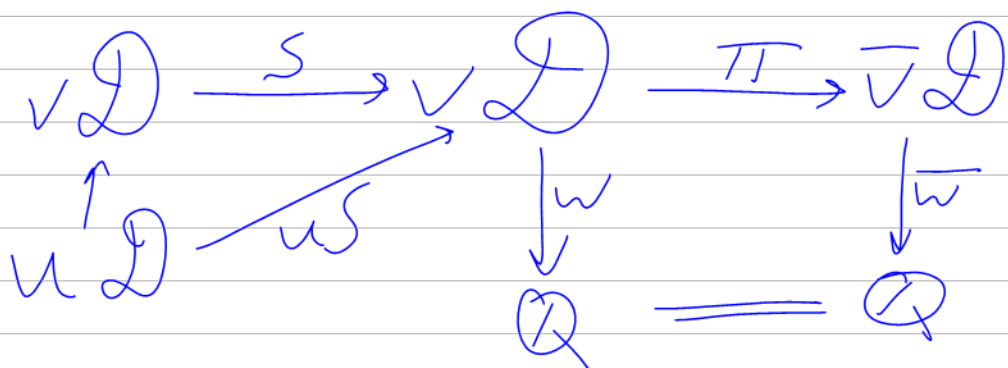


$$V_2(\textcircled{\times}) = 1$$





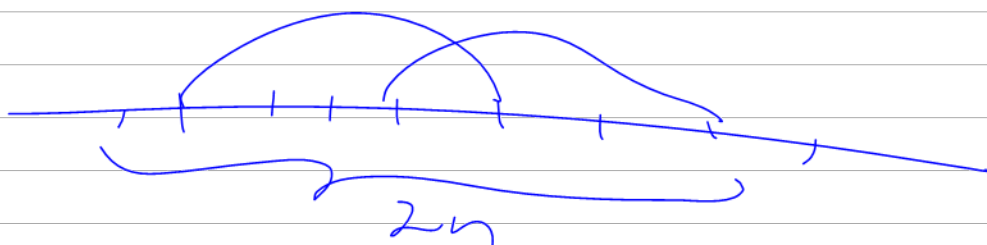
$$\overline{vQ} = vQ / \text{arc } s = s \text{ arc}$$



$$\text{im}(us)^* \stackrel{\cong}{=} \text{im}(us // \pi)^*$$

$$\text{Ker}(us)^\perp \stackrel{\cong}{\subset} \text{Ker}(us // \pi)^\perp$$

$$\text{Ker } us \supset \text{Ker}(us // \pi)$$



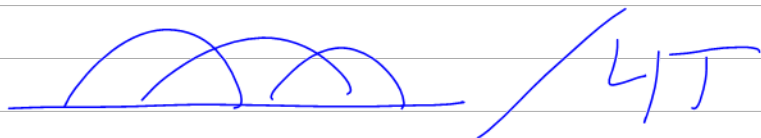
$$\binom{2n}{n} n! 2^n \quad \frac{2n!}{n!} 2^n$$

$$6 \cdot 7 \cdot 8 \cdot 9 \cdot 10 \cdot 32$$

$$Z: X \rightarrow A$$

HIS

"Expansion"
VFTI



Exercise Prove GPV for w-knots.

$$PwB_n = \langle \sigma_{ij} : 1 \leq i \neq j \leq n \rangle \left. \begin{array}{l} \text{R-moves} \\ \text{OC} \\ \text{"overcrossings commute"} \end{array} \right\rangle$$

= motion group of
horizontal flying rings in \mathbb{R}^3

σ_{ij}^{\pm} = "ring j flies through ring i in the \pm direction."

Relations 1. $\sigma_{ij} \sigma_{kl} = \sigma_{kl} \sigma_{ij} \quad | \{i, j, k, l\} | = 4$

2. $\sigma_{ij}^+ \sigma_{ij}^- = 1$

3. $\sigma_{jk} \sigma_{ik} \sigma_{ij} = \sigma_{ij} \sigma_{ik} \sigma_{jk}$

$i = \text{big}$
 $j = \text{med}$
 $k = \text{small}$

4 OC: $\sigma_{ij} \sigma_{jk} = \sigma_{jk} \sigma_{ij}$



$$G \sim \mathbb{Q}G = \langle \sum a_i g_i \rangle \quad I = \langle \sum a_i g_i, \sum a_i = 0 \rangle$$

$$\langle g-1 \rangle$$

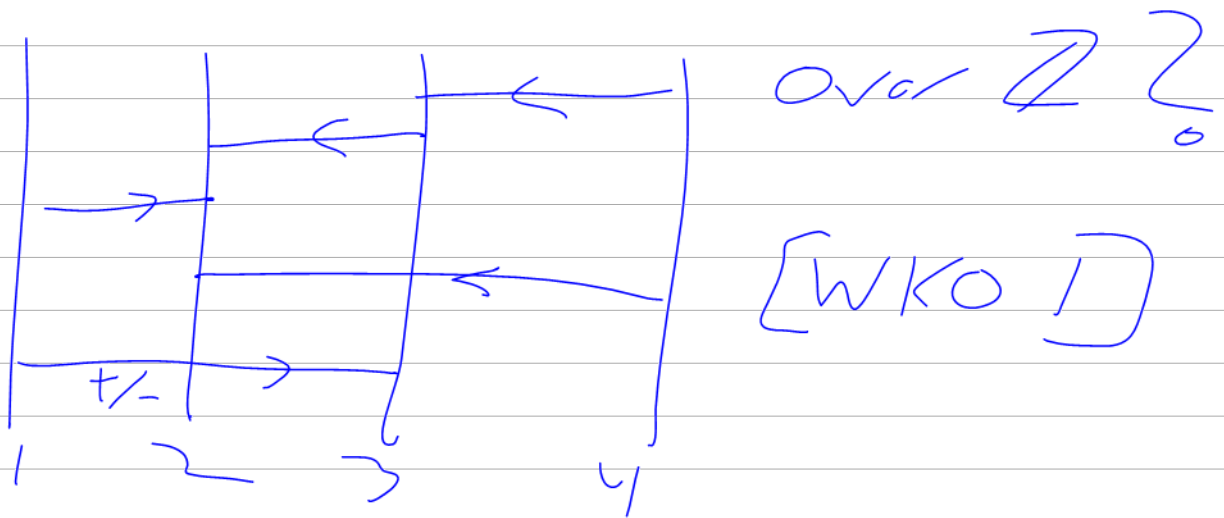
$$\mathbb{Q}G = I^0 \supset I^1 \supset I^2 \supset I^3 \supset \dots$$

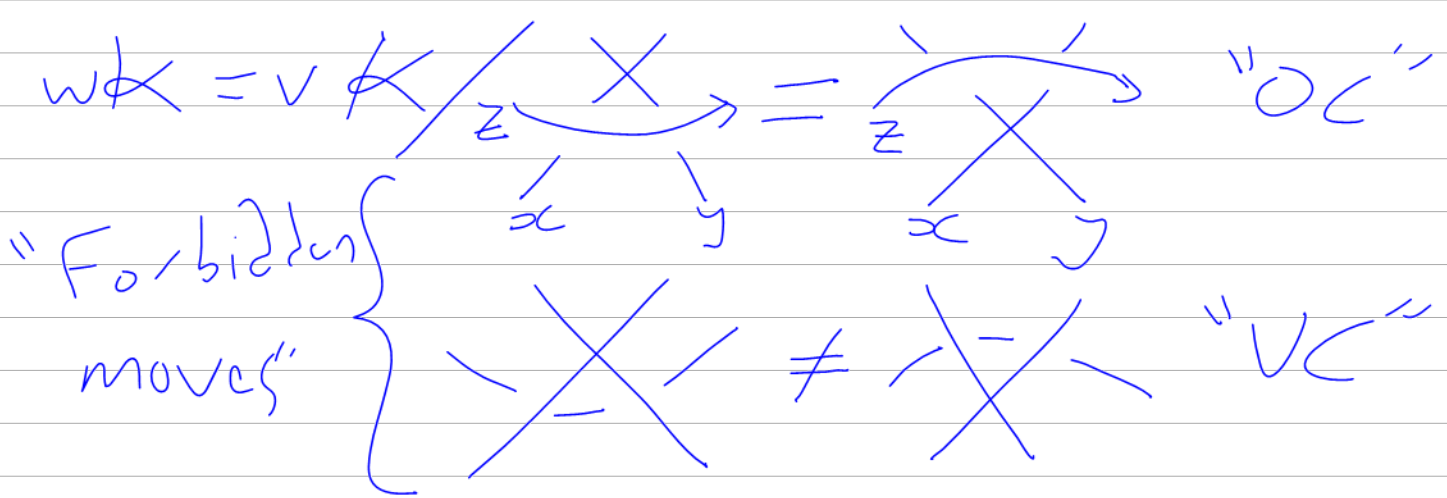
Def $\varphi: G \rightarrow \mathbb{Q}$ is of type m if

$$\varphi|_{I^{m+1}} \equiv 0$$

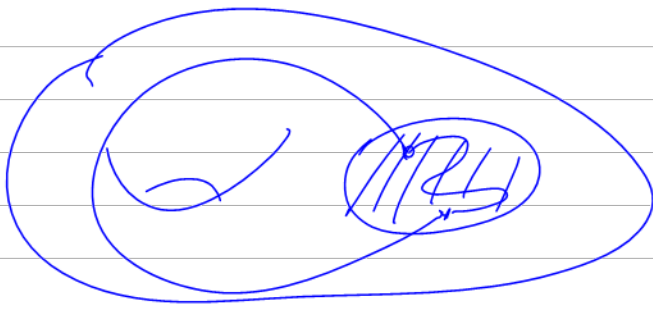
$$\varphi(\dots (g_1-1) \dots (g_2-1) \dots (g_{m+1}-1) \dots) = 0$$

Exercise GPV holds over \mathbb{Q}





$$PwB_n = \pi_1 \left(\begin{array}{l} \text{Conf space} \\ \text{of } \underbrace{\text{rings}}_{\text{bar}} \text{ in } \mathbb{R}^3 \end{array} \right)$$



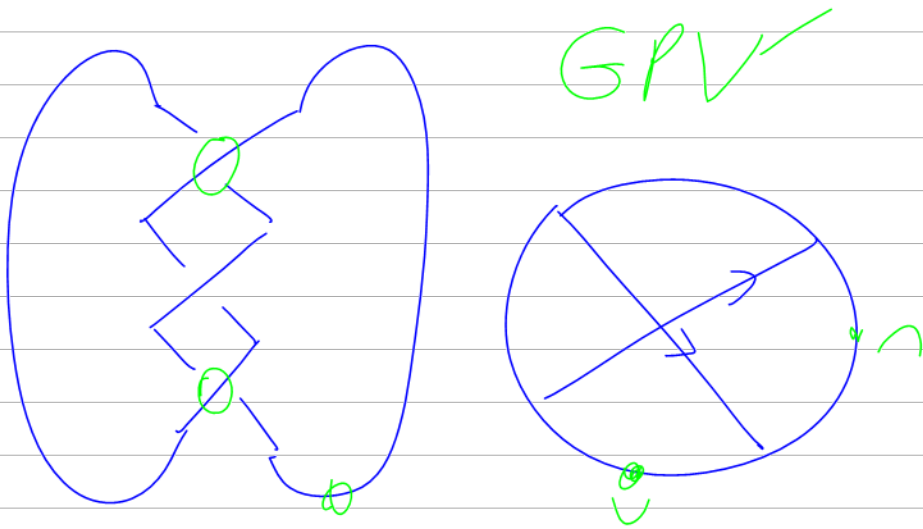
$$\pi_1 \left(\left(\text{rings} \right)^c \right)$$

$$PwB_n \rightarrow \text{Aut}(F_n)$$

$$C(\mathcal{O}^k) = \delta_{k,1}$$

$$C(\begin{array}{c} \nearrow \\ \searrow \end{array}) - C(\begin{array}{c} \nearrow \\ \nearrow \end{array}) = z \cdot C(\begin{array}{c} \uparrow \\ \uparrow \end{array})$$

$$C(K) = \sum_{m=0}^{\infty} \underbrace{V_m(K)}_{V_2} z^m$$



GPV ✓

$$V = W \circ S'$$

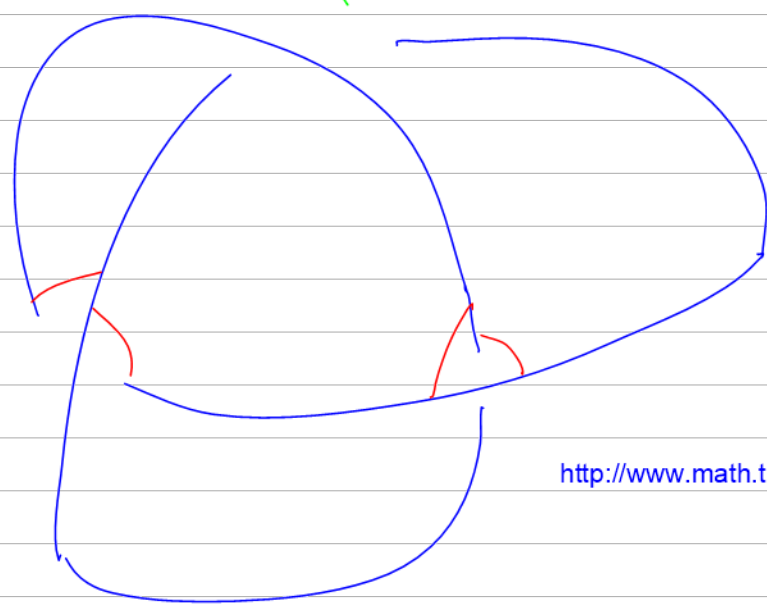
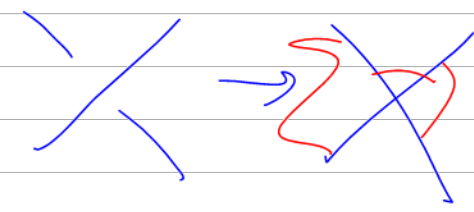
Expansion \Leftrightarrow UFTI

$$Z: X \rightarrow A \quad Z(K_D) = D + \dots$$

$$D \rightarrow \Sigma \otimes$$



$$Z(\cdot) = X + X + \dots$$



$$\begin{aligned} \text{ad}_a x &:= [a, x] = ax - xa \\ &= L_a x - R_a x \end{aligned}$$

$$\begin{aligned} e^{\text{ad}_a} x &= x + [a, x] + \frac{1}{2}[a, [a, x]] + \dots \\ &\parallel e^{L_a - R_a} x = e^{L_a} e^{R_a} x \\ &= \underline{e^a x e^{-a}} \end{aligned}$$

$$\dot{g}_s(x_i) = \frac{d}{ds} g_s(x_i)$$

$$L_s \in M_{n \times n}(\mathbb{R})$$

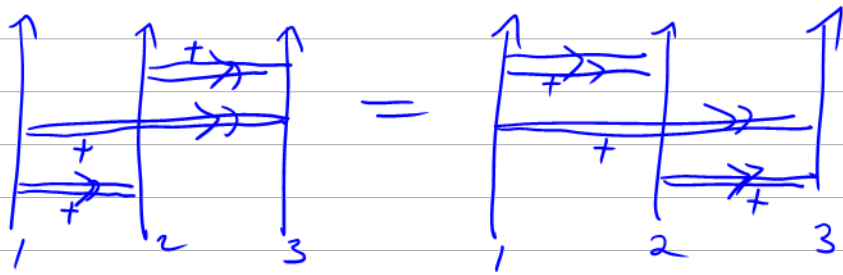
$$\rho(L_s)(V) = \frac{d}{ds}(L_s V)$$

$$\parallel \dot{E} = \sum x_i \frac{\partial}{\partial x_i}$$

$$A_s(x^d) = s^d x^d \quad A_s^{-1}(x^d) = s^{-d} x^d$$

$$\dot{A}_s(x^d) = d s^{d-1} x^d$$

$$\dot{A}_s A_s^{-1}(x^d) = \dot{A}_s (s^{-d} x^d) = d s^{-1} x^d$$



$$\underline{R_{12} R_{13} R_{23}} = R_{23} R_{13} R_{12}$$

$$| \Rightarrow | \xrightarrow{Z} \mathcal{A}^V(\uparrow \uparrow) = \langle \uparrow \uparrow \rangle / \text{GT}$$

$$Z(| \Rightarrow |) = \uparrow \uparrow + \uparrow \rightarrow \uparrow + \text{h.o.} = R$$



$$\square(V) = V \otimes V \quad V = e^p \text{ up to incl. deg 7.}$$

$$V e^{-p} = 1 + \underset{\text{deg 8}}{p} + \text{h.o.}$$

$$\square(V e^{-p}) = V e^{-p} \otimes V e^{-p}$$

$$\square(1+p) = (1+p) \otimes (1+p) \text{ up to deg 8}$$

$$\square p = | \otimes p + p \otimes | \quad \text{so } p \text{ is primitive.}$$

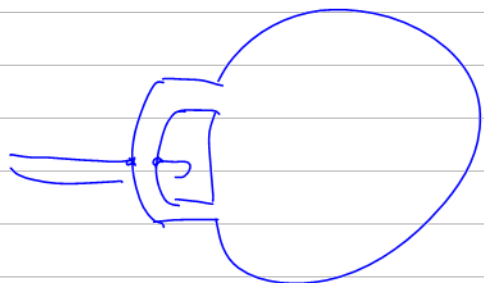
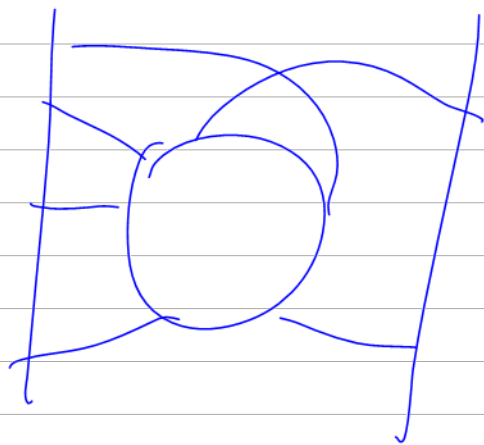
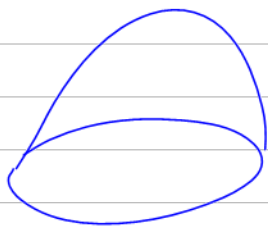
$$p' = p + p \quad V e^{-p'} = V e^{-p} e^{-p} = (1+p)(1-p) = 1$$



$$g^* \otimes g^* \otimes g^*$$

$$\langle A, B \rangle = \text{tr}(A \cdot B)$$

$$\langle X_{ij}, X_{kl} \rangle = \delta_{jk} \delta_{li}$$



[1] K_n : knots in comp of n poles.

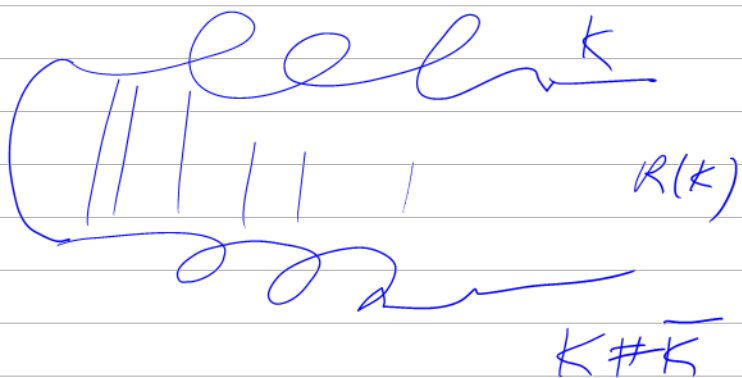
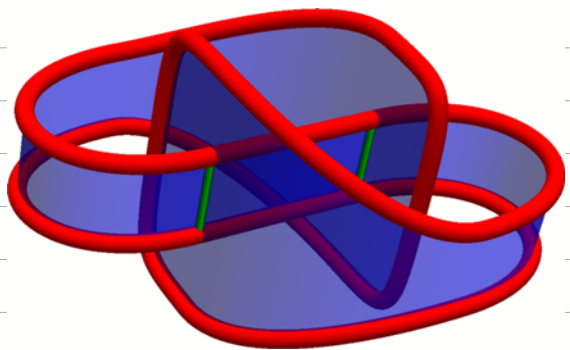
[1] A_n : $gr K_n = [1 \circlearrowleft | 0 |] / rels$

[1] M_n : A_n in $gl(N)$: surface diagrams / rels

m_n / hc
 $\uparrow \varphi$

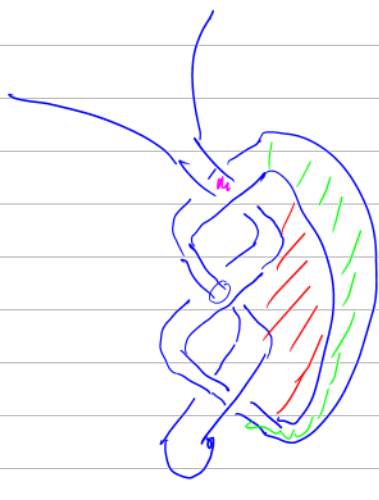
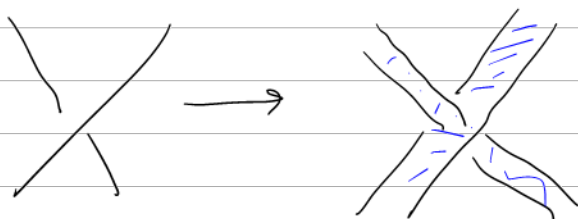
$$[\varphi(w_1), \varphi(w_2)] = \varphi([w_1, w_2])$$

[1] C_n : cyclic words

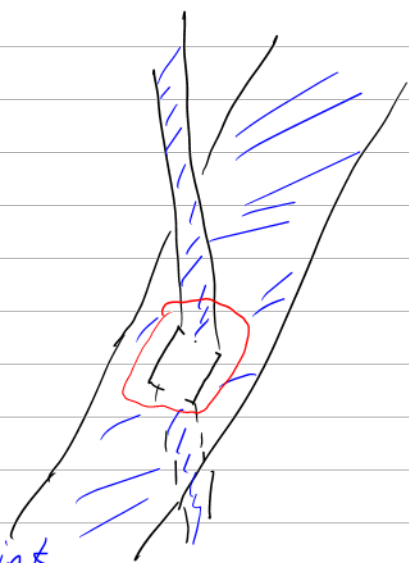


$$\Delta(R(K)) = \Delta(K)\Delta(\bar{K}) = \Delta^2$$

$W1 \rightarrow u$



	small loops	faces
small loops	O	A
faces	A ^T	C



I think

$$\Delta(R(K)) = |\Delta(K)|^2$$

$$\begin{array}{ccc} \gamma^* B & \xrightarrow{\alpha} & B \\ \downarrow & & \downarrow \beta \\ C & \xrightarrow{\gamma} & A \end{array} \quad \text{"Equalizer"}$$

$$\gamma^* B = \{(c, b) \in C \times B : \gamma c = \beta b\} \stackrel{\alpha}{\cong} \beta^* C = B \oplus_A C$$

Claim? pullbacks of pushforward scenes are pushforward scenes

"push & pull commute"

$$\begin{array}{ccc} \alpha^* S \supset B \oplus_A C & \xrightarrow{\alpha} & B \supset S \\ \downarrow & \downarrow & \downarrow \beta \\ \gamma^* \beta^* S \subset C & \xrightarrow{\gamma} & A \supset \beta^* S \end{array}$$

$$\text{if } \begin{array}{ccc} B \oplus_A C & \rightarrow & B \\ \downarrow & & \downarrow \\ C & \rightarrow & A \end{array}$$

I think $\begin{array}{ccc} D & \xrightarrow{\alpha} & \circ B \\ \downarrow \delta & \searrow \pi & \downarrow \beta \\ C & \xrightarrow{\gamma} & \circ A \end{array}$ is a pullback scene, ie, $D \cong C \oplus_A B$

iff

- $\ker \alpha \cap \ker \gamma = 0$
- $\text{im } \alpha \supset \ker \beta$ and $\text{im } \delta \supset \ker \gamma$
- $\ker \pi = \ker \alpha + \ker \delta$

Proof?

$$\begin{array}{ccc} A \oplus E \oplus F & & \\ \alpha \swarrow & \downarrow \pi & \searrow \gamma \\ A \oplus B \oplus E & & A \oplus C \oplus F \\ \beta \downarrow & & \downarrow \gamma \\ A \oplus B \oplus C \oplus D & & \end{array}$$

This is the most general pullback diagram!

$$\begin{array}{ccc} V_1 & & V_2 \\ \beta \searrow & & \searrow \delta \\ V_0 & & \end{array}$$

