

Pensieve header: Proof of the main \mathfrak{g}_0 lemma and a poly-time program to compute the \mathfrak{g}_0 invariant.

Reminder

Make sure that you have Mathematica and that you play with these programs!

Representing $\mathfrak{g}_0 = \langle h, e, l, f \rangle / ([e, l] = -e, [f, l] = f, [e, f] = h, [h, *] = 0)$

$$\rho h = \begin{pmatrix} 0 & 0 & -1 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}; \rho e = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}; \rho l = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}; \rho f = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}; \rho \theta = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix};$$

`B[x_?MatrixQ, y_?MatrixQ] := x.y - y.x;`

`{B[ρe, ρl] == -ρe, B[ρf, ρl] == ρf, B[ρe, ρf] == ρh, B[ρh, ρe] == ρθ, B[ρh, ρl] == ρθ, B[ρh, ρf] == ρθ}`
`{True, True, True, True, True, True}`

The Main \mathfrak{g}_0 Theorem

Raw Version. The \mathfrak{g}_0 invariant of any S-component tangle T can be written in the form $Z(T) = \mathcal{O}(\omega e^{L+Q} \mid \prod_{i \in S} e_i l_i f_i)$, where ω is a scalar (meaning, a rational function in the variables h_i and their exponentials $t_i = e^{h_i}$), where $L = \sum a_{ij} h_i l_j$ is a balanced quadratic in the variables h_i and l_j with integer coefficients a_{ij} and where $Q = \sum b_{ij} e_i f_j$ is a balanced quadratic in the variables e_i and f_j with scalar coefficients b_{ij} . Furthermore, after setting $h_i = h$ and $t_i = t$ for all i , the invariant $Z(T)$ is poly-time computable.

Proof. Indeed, as shown below,

$$0. R^s = e^{s(h \otimes l + e \otimes f)} = \mathcal{O}(\exp(s h l + \frac{e^{s h} - 1}{h} e f \mid e \otimes l),$$

$$1. \mathcal{O}(e^{\gamma l + \beta e} \mid l e) = \mathcal{O}(e^{\gamma l + e^\gamma \beta e} \mid e l),$$

$$2. \mathcal{O}(e^{\gamma l + \beta f} \mid f l) = \mathcal{O}(e^{\gamma l + e^\gamma \beta f} \mid l f),$$

$$3. \mathcal{O}(e^{\beta e + \alpha f + \delta e f} \mid f e) = \mathcal{O}(v e^{\gamma l + \beta e + \alpha f + \delta e f} \mid e f), \text{ with } v = (1 + h \delta)^{-1},$$

and the rest is straight-forward.

Proofs of the \mathfrak{g}_0 lemmas

$$(* \ 0 \ *) \text{MatrixForm} \ /@ \ { \text{MatrixExp}[h \rho l + e \rho f], \text{MatrixExp}[h \rho l] . \text{MatrixExp}[\frac{e^h - 1}{h} e \rho f] }$$

$$\left\{ \begin{pmatrix} 1 & \frac{e(-1+e^h)}{h} & 0 \\ 0 & e^h & 0 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & \frac{e(-1+e^h)}{h} & 0 \\ 0 & e^h & 0 \\ 0 & 0 & 1 \end{pmatrix} \right\}$$

$$(* \ 1 \ *) \text{MatrixForm} \ /@ \ { \text{MatrixExp}[\gamma \rho l] . \text{MatrixExp}[\beta \rho e], \text{MatrixExp}[e^\gamma \beta \rho e] . \text{MatrixExp}[\gamma \rho l] }$$

$$\left\{ \begin{pmatrix} 1 & 0 & 0 \\ 0 & e^\gamma & e^\gamma \beta \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 \\ 0 & e^\gamma & e^\gamma \beta \\ 0 & 0 & 1 \end{pmatrix} \right\}$$

$$(* \ 2 \ *) \text{MatrixForm} \ /@ \ { \text{MatrixExp}[\beta \rho f] . \text{MatrixExp}[\gamma \rho l], \text{MatrixExp}[\gamma \rho l] . \text{MatrixExp}[e^\gamma \beta \rho f] }$$

$$\left\{ \begin{pmatrix} 1 & e^\gamma \beta & 0 \\ 0 & e^\gamma & 0 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & e^\gamma \beta & 0 \\ 0 & e^\gamma & 0 \\ 0 & 0 & 1 \end{pmatrix} \right\}$$

(* 3 at $\delta=0$ *)

`MatrixForm /@ {MatrixExp[$\alpha \rho f$].MatrixExp[$\beta \rho e$], MatrixExp[$-\alpha \beta \rho h$].MatrixExp[$\beta \rho e$].MatrixExp[$\alpha \rho f$]}`

$$\left\{ \begin{pmatrix} 1 & \alpha & \alpha \beta \\ 0 & 1 & \beta \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & \alpha & \alpha \beta \\ 0 & 1 & \beta \\ 0 & 0 & 1 \end{pmatrix} \right\}$$

For the full proof of 3, see the blackboard and then check:

`With[{ $\psi = v e^{v(t e^{f-\alpha \beta h+\alpha f+\beta e})}$ /. $v \rightarrow (1+t h)^{-1}$ }, Simplify@{ $\partial_t \psi - \partial_{\alpha, \beta} \psi$, ψ /. $t \rightarrow 0$ }]`
`{0, $e^{f \alpha + e \beta - h \alpha \beta}$ }`

Implementation

```
CF[E[ $\omega$ _, L_, Q_]] := E[Simplify[ $\omega$ ], Simplify[L], Simplify[Q]];
E /: E[ $\omega 1$ _, L1_, Q1_] E[ $\omega 2$ _, L2_, Q2_] := CF@E[ $\omega 1 \omega 2$ , L1 + L2, Q1 + Q2];
E[ $\omega 1$ _, L1_, Q1_]  $\equiv$  E[ $\omega 2$ _, L2_, Q2_] := Simplify[ $\omega 1 == \omega 2 \wedge L1 == L2 \wedge Q1 == Q2$ ];
```

$$0. R = e^{h\otimes + e\otimes f} = \mathcal{O}(\exp(hl + \frac{e^{h-1}}{h} ef \mid e\otimes l f):$$

```
E[X_{i,j}^+] := E[1, h_i l_j, h_i^{-1} (e^{h_i} - 1) e_i f_j];
E[X_{i,j}^-] := E[1, -h_i l_j, h_i^{-1} (e^{-h_i} - 1) e_i f_j];
E[p_Times] := E /@ p;
```

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+]$$

- $\mathcal{O}(e^{v l + \beta e} \mid e) = \mathcal{O}(e^{v l + e^v \beta e} \mid e l)$,
- $\mathcal{O}(e^{v l + \beta f} \mid f l) = \mathcal{O}(e^{v l + e^v \beta f} \mid l f)$:

```
NO_{(x:f|e)_i 1_j} [E[ $\omega$ _, L_, Q_]] := CF[E[ $\omega$ , L, e^y  $\alpha x_i + (Q / . x_i \rightarrow \theta)$  /. { $y \rightarrow \partial_{1_j} L$ ,  $\alpha \rightarrow \partial_{x_i} Q$ }]];
ANO_{(x:f|e)_i 1_j} [E[ $\omega$ _, L_, Q_]] := CF[E[ $\omega$ , L, e^{-y}  $\alpha x_i + (Q / . x_i \rightarrow \theta)$  /. { $y \rightarrow \partial_{1_j} L$ ,  $\alpha \rightarrow \partial_{x_i} Q$ }]];
```

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+]$$

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // ANO_{e_2 1_3}$$

$$(\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // ANO_{e_2 1_3} // NO_{e_2 1_3}) == \mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+]$$

$$3. \mathcal{O}(e^{\beta e + \alpha f + \delta e f} \mid f e) = \mathcal{O}(v e^{v(-\alpha \beta h + \beta e + \alpha f + \delta e f)} \mid e f), \text{ with } v = (1 + h\delta)^{-1}:$$

```
NO_{f_i e_j \rightarrow k} [E[ $\omega$ _, L_, Q_]] := CF[
  E[v  $\omega$ , L, v (- $\alpha \beta h_k + \beta e_k + \alpha f_k + \delta e_k f_k$ ) + (Q / . f_i | e_j \rightarrow \theta)]
  /. v \rightarrow (1 + h_k \delta)^{-1} /. { $\alpha \rightarrow \partial_{f_i} Q / . e_j \rightarrow \theta$ ,  $\beta \rightarrow \partial_{e_j} Q / . f_i \rightarrow \theta$ ,  $\delta \rightarrow \partial_{f_i, e_j} Q$ }];
```

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // NO_{f_3 e_4 \rightarrow 7}$$

The Stitching Formula

```
m_{i,j \rightarrow k} [Z_] := Module[{x, z}, CF[(Z // NO_{f_i e_j \rightarrow x} // NO_{1_i e_x} // NO_{f_x 1_j}) /. z_{-i|j|x} \rightarrow z_k]]
```

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // m_{1,2 \rightarrow 1}$$

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // m_{1,2 \rightarrow 1} // m_{1,3 \rightarrow 1}$$

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // m_{1,2 \rightarrow 1} // m_{1,3 \rightarrow 1} // m_{1,4 \rightarrow 1}$$

$$\mathbb{E}[X_{4,1}^+ X_{2,5}^+ X_{6,3}^+] // m_{1,2 \rightarrow 1} // m_{1,3 \rightarrow 1} // m_{1,4 \rightarrow 1} // m_{1,5 \rightarrow 1} // m_{1,6 \rightarrow 1}$$

Independent Proof of Invariance

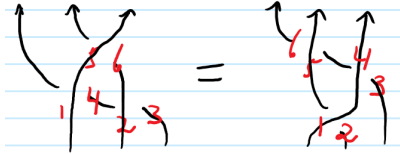
Meta-Associativity:

$$\xi = \mathbb{E} \left[\omega, \sum_{i=1}^4 \sum_{j=1}^4 a_{i,j} h_i l_j, \sum_{i=1}^4 \sum_{j=1}^4 b_{i,j} e_i f_j \right]$$

$$\xi // m_{1,2 \rightarrow 1}$$

$$\xi // m_{1,2 \rightarrow 1} // m_{1,3 \rightarrow 1}$$

$$(\xi // m_{1,2 \rightarrow 1} // m_{1,3 \rightarrow 1}) \equiv (\xi // m_{2,3 \rightarrow 2} // m_{1,2 \rightarrow 1})$$



Reidemeister 3:

$$\text{lhs} = \mathbb{E} [X_{1,4}^+, X_{2,3}^+, X_{5,6}^+] // m_{1,5 \rightarrow 1} // m_{2,6 \rightarrow 2} // m_{3,4 \rightarrow 3}$$

$$\text{rhs} = \mathbb{E} [X_{1,2}^+, X_{4,3}^+, X_{5,6}^+] // m_{1,4 \rightarrow 1} // m_{2,5 \rightarrow 2} // m_{3,6 \rightarrow 3}$$

$$\text{lhs} \equiv \text{rhs}$$

Homework.

1. Use the same methodology to verify $m_{a,b \rightarrow c} // m_{d,e \rightarrow f} \equiv m_{d,e \rightarrow f} // m_{a,b \rightarrow c}$.
2. Likewise, verify the two types of R2 moves.
3. Make sure that R1 gives no trouble.
4. Implement the “polished version” of the main theorem below, and verify that everything works.

The Main g_0 Theorem, Polished Version

Polished Version. With $\bar{e} = \frac{(e^h - 1)}{h} e$, the g_0 invariant of any S-component tangle T can be written in the form $Z(T) = \mathcal{O}(\omega^{-1} e^{L + \omega^{-1} Q} \mid \prod_{i \in S} \bar{e}_i l_i f_i)$, where ω is a scalar (meaning, a **polynomial** in the variables $t_i = e^{h_i}$), where $L = \sum a_{ij} h_i l_j$ is a balanced quadratic in the variables h_i and l_j with integer coefficients a_{ij} and where $Q = \sum b_{ij} \bar{e}_i f_j$ is a balanced quadratic in the variables \bar{e}_i and f_j with scalar coefficients b_{ij} . Furthermore, after setting $t_i = t$ for all i , the invariant $Z(T)$ is poly-time computable.