

Riddle Along.

$$\left(\begin{array}{c} 4 \\ p \end{array} \right) v_L = v_S$$

Read Along. 1.4-1.6.

Web Fact: (not visible) = (doesn't exist).

Life Fact. No "teaching over email".

Reminders. We seek "basis"; l.c.; span; "generates"

DEF A subset $S \subset V$ is "lin. dep" if it is "wasteful".I.e., if $\exists a_i \in F$ not all 0 & $\sum a_i v_i = 0$

Otherwise, it is "lin. indep."

Examples $\{e_i\}$ ✓, $\left\{ \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix}, \begin{pmatrix} 4 \\ 5 \\ 6 \end{pmatrix}, \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix} \right\}$ Comments 1. \emptyset is lin. indep.2. $\{u\}$ is lin indep iff $u \neq 0$.3. Suppose $S_1 \subset S_2 \subset V$. Thena. If S_1 is dep, so is S_2 b. If S_2 is indep, so is S_1 4. If S is lin indep in V and $v \in V \setminus S$, then $S \cup \{v\}$ is lin. dep. iff $v \in \text{span}(S)$. } skippedDEF Basis $\beta \subset V$ Examples: 1. \emptyset for $\{0\}$.2. e_i for F^n 3. E^{ij} for $M_{m \times n}(F)$ 4. $(1, x, \dots, x^n)$ for $P_n(F)$ 5. $(1, x, \dots)$ for $P(F)$ 6. $\left\{ \begin{pmatrix} 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ -1 \end{pmatrix} \right\}$ for \mathbb{R}^2 . $\begin{pmatrix} a \\ b \end{pmatrix} = \frac{a+b}{2} \begin{pmatrix} 1 \\ 1 \end{pmatrix} + \frac{a-b}{2} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$ Thm A subset $\beta \subset V$ is a basis iff every $v \in V$ can be expressed in a unique way as a l.c. of elements of β .Thm If a finite set S generates a v.s. V ,

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then there is a subset $B \subseteq S$ which is a basis

of V

PF Let B be a lin indep subset of S which is of maximal size. The way $v \in S \setminus B$ satisfies $v \in \text{span } B$, so $S \subseteq \text{span } B$, so $\text{span } S \subseteq \text{span } B$.

Our first non-language theorem:

Thm If a v.s. V has a finite basis, then every other basis of V has the same number of elements in it.

target line

Def If V has a finite basis, we say that it is "finite-dimensional" and let

$$\dim V := \left(\begin{array}{l} \text{The number of elements} \\ \text{in (any) basis of } V \end{array} \right)$$

Examples as above:

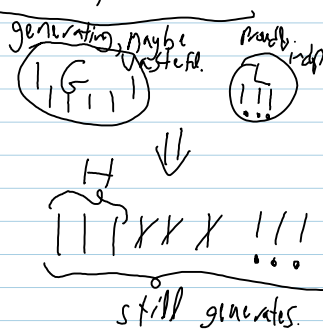
- $\{0\}, F^n,$
- $M_{m \times n}, P_n(F),$
- $P(F)$

Lemma (the replacement lemma)

$|G| = n, \text{span } G = V, L$ lin indep

$\Rightarrow |L| \leq n$ & $\exists H \subseteq G$ with

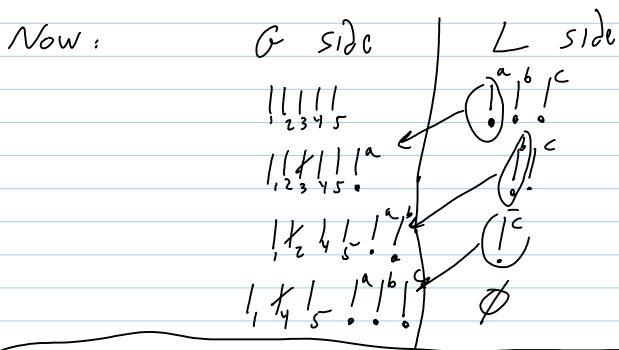
$$|H| = n - |L| \text{ and } \text{span}(H \cup L) = V$$



PF of Theorem from Lemma.

Informal proof of Lemma First of all, if $\sum a_i u_i = 0$, then any vector that appears in this dependency with non-zero coeff is a l.c. of the others.

09 hour 13



Formal proof: Induction on $|L|$. $|L|=0$: trivial.

Now $|L|=m+1$; $L = \{v_1, \dots, v_{m+1}\}$. Use $L' = \{v_1, \dots, v_m\}$,

Find $H' = \{u_1, \dots, u_{n-m}\} \subseteq G$ s.t. $\{u_1, \dots, u_{n-m}, v_1, \dots, v_m\}$

spans. write

$$v_{m+1} = a_1 u_1 + \dots + a_{n-m} u_{n-m} + b_1 v_1 + \dots + b_m v_m$$

\therefore Not all $a_i = 0$. $\hookrightarrow n-m > 0$. so $m+1 \leq n$.

\therefore w.l.o.g. $a_1 \neq 0$, so $u_1 \in \text{span}(u_2, \dots, u_{n-m}, v_1, \dots, v_{m+1})$,
so take $H = \{u_2, \dots, u_{n-m}\}$.

Corollaries: 1. IF V has a finite basis β_1 then every other basis β_2 of V is also finite & $|\beta_1| = |\beta_2|$.

2. "dim V " makes sense.

3. Assume $\dim V = n$. Then $\frac{\text{hour } 13}{\text{hour } 14}$

a. IF G generates V , $|G| \geq n$ & if also $|G| = n$, then G is a basis.

b. IF L is linearly indep in V , then $|L| \leq n$;
if also $|L| = n$, L is a basis.

if also $|L| < n$, L can be extended to a basis.

4. IF V is finite-dimensional and $W \subset V$ is a subspace, then W is f.d. and $\dim W \leq \dim V$.

IF also $\dim W = \dim V$, then $W = V$.

IF also $\dim W < \dim V$, then any basis of W can be extended to a basis of V .