

Plan. UFO blunder, JCF abstractly & in practice.

I said "I think in a UFO every prime ideal is maximal"

JCF. V a f.d.v.s, $A: V \rightarrow V$ linear, makes V a module over $F[x]$ via $xu = Au$. Then

$$V \cong \bigoplus F[x]/(x-\lambda_i)^{s_i}. \text{ What's } \frac{F[x]}{(x-\lambda_i)^{s_i}}?$$

UFO Blunder. The above statement is nonsense.

In $\mathbb{Q}[x,y] = \mathbb{Q}[x][y]$, $\langle x \rangle$ is prime but not maximal.

Basis: $1, x-\lambda, (x-\lambda)^2, \dots, (x-\lambda)^{s-1}$

$A-\lambda$ acts by "shift to the right" $\begin{pmatrix} 0 & 0 \\ 1 & 0 \\ \vdots & \vdots \end{pmatrix}$

so A acts by $\begin{pmatrix} \lambda & & \\ & \lambda & \\ & & \lambda \end{pmatrix}$

Now lets do that in practice....

step 1. Find a presentation matrix for $V \in R\text{-mod}$.

w.l.o.g $V = F^n$ and $A \in M_{n \times n}(F)$.

$$\ker \pi = \zeta$$

$$r_i = x e_i - A e_i \in \ker \pi$$

claim $\langle r_i \rangle = \ker \pi$

PF Consider

$$\begin{array}{l} R^n \xrightarrow{xI-A} R^n \xrightarrow{\pi} F^n \\ e_i \mapsto e_i \\ x^k e_i \mapsto A^k e_i \end{array}$$

$$F^n \xrightarrow[\text{v.s. map } \beta]{\text{onto? } \alpha} R^n / \langle r_i \rangle \xrightarrow[\text{well-def!}]{1-I} R^n / \ker \pi \xrightarrow{\sim} F^n$$

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We want to know if α is 1-1; it is enough to show that β is onto; i.e., that any $x^k e_i$ can be written, modulo $\langle r_i \rangle$,

Corollary 2. Over an algebraically closed field \mathbb{F} , every square matrix

$$A \text{ is conjugate to a block diagonal matrix } B = \begin{pmatrix} B_1 & 0 & \dots & 0 \\ 0 & B_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & B_n \end{pmatrix},$$

where each B_i is either a 1×1 matrix (λ_i) for some $\lambda_i \in \mathbb{F}$, or an $s_i \times s_i$ matrix with λ_i 's on the diagonals, 1's right below the diagonal, and 0's elsewhere,

$$\begin{pmatrix} \lambda_i & 0 & \dots & \dots & 0 & 0 \\ 1 & \lambda_i & \ddots & & & 0 \\ 0 & \ddots & \ddots & \ddots & & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & & & & \lambda_i & 0 \\ 0 & 0 & \dots & 0 & 1 & \lambda_i \end{pmatrix},$$

for some $\lambda_i \in \mathbb{F}$ and for some $s_i \geq 2$. Furthermore, B is unique up to a permutation of its blocks B_i .

(Corollary: good old diagonalization.)

as a combination of e_j 's. Indeed,

$$x^k e_i = x^{k-1}(x e_i) = x^{k-1} A e_i = \dots = A^k e_i$$

Go over handout, first in the distinct-eigenval's case:

Row and Column Operations

Row operations are performed by left-multiplying N by some properly-positioned 2×2 matrix and at the same time left-multiplying the "tracking matrix" P by the same 2×2 matrix. Column operations are similar, with left replaced by right and P by Q .

```

RowOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[TT.NN]; PP = Simplify[TT.PP];
];
ColOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[NN.TT]; QQ = Simplify[QQ.TT];
];

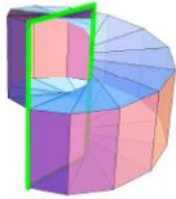
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Swapping Rows and Columns

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SwapRows[i_, j_] := RowOp[i, j, {{0 1}, {1 0}}];
SwapColumns[i_, j_] := ColOp[i, j, {{0 1}, {1 0}}];
SwapBoth[i_, j_] := {SwapRows[i, j], SwapColumns[i, j]};

```



?

The "GCD" Trick

If $q = \gcd(a, b) = sa + tb$, the equality $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} q \\ 0 \end{pmatrix}$ allows us to replace pairs of entries in the same column by their greatest common divisor (and a zero!), using invertible row operations. A similar trick works for rows.

```

GCDTrick[i_, j_, k_] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[i, k]], b = NN[[j, k]], x];
  RowOp[i, j, {{s, t}, {-b/q, a/q}}];
];
GCDTrick[k_, {i_, j_}] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[k, i]], b = NN[[k, j]], x];
  ColOp[i, j, {{s, t}, {-b/q, a/q}}];
];

```

Factoring Diagonal Entries

If $1 = \gcd(a, b) = sa + tb$, the equality $\begin{pmatrix} sa & 1 \\ -tb & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & ab \end{pmatrix} \begin{pmatrix} a & -b \\ t & s \end{pmatrix} = \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$ is an invertible row-column-operations proof of the isomorphism $\frac{R}{(a)} \oplus \frac{R}{(b)} = \frac{R}{(ab)}$.

```

SplitToSum[i_, j_, a_, b_] := Module[{
  {q, s, t, T1, T2},
  {q, {s, t}} = PolynomialExtendedGCD[a, b, x];
  If[q == 1,
    RowOp[i, j, {{sa, 1}, {-tb, 1}}]; ColOp[i, j, {{a, -b}, {t, s}}];
  ]
];

```

Recovering C From P ?

$$\begin{array}{ccc}
 R^n \xrightarrow[\mathcal{M}]{Ix - A} R^n \xrightarrow{T_A} F^n & & \\
 \uparrow Q & & \downarrow P \\
 R^n \xrightarrow{Ix - B} R^n \xrightarrow{T_B} F^n & & \downarrow C
 \end{array}$$

$$\begin{aligned}
 C e_i &= T_B(P e_i) \\
 &= T_B(\sum x^k P_k e_i) \\
 &= \sum x^k T_B(P_k e_i) \\
 &= \sum B^k P_k e_i
 \end{aligned}$$

$$\Rightarrow C = \sum B^k P_k \dots \text{complete run 1}$$

The "Jordan Trick":

Then go through run 2
& run 3

A repeated application of the identity $\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{1+k} & 0 \\ 1 & p \end{pmatrix}$ will bring a matrix like

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix}$$

to the "Jordan" form of $\begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}$, using invertible row and column operations.

```

JordanTrick[i_, j_, p_, s_] := {RowOp[i, j, {{p^{s-1}, -1}, {1, 0}}], ColOp[i, j, {{1, p}, {0, 1}}]};

```

done line

^(debt)
Theorem. The universal property for tensor products.

1. Holds
2. Determines $M \otimes N$ up to a unique isomorphism.

