

xtx handout on 151102-09:40

Monday, November 2, 2015 9:37 AM

Dror Bar-Natan: Talks: MoscowByWeb-1511:
weβ:=http://drorbn.net/mbw

Abstract. The subject will be very close to Manturov's representation of vB_n into $\text{Aut}(FG_{n+1})$ — I'll describe how I think about it in terms of a very simple minded map \mathcal{K} from n -component v-tangles to $(n+1)$ -component w-tangles. It is possible that you all know this already. Possibly my talk will be very short — it will be as long as it is necessary to describe \mathcal{K} and say a few more words, and if this is little, so be it.

- All you need is \mathcal{K} ... • What is its domain? • What is its target?
• Why should one care?

Virtual Knots. Virtual knots are the algebraic structure underlying the Reidemeister presentation of ordinary knots, without the topology. Locally they are knot diagrams modulo the Reidemeister relations; globally, who cares? So,

$$vT = CA < \times, \times ! - - - - >$$

Flying Pogs

Note that also

$$vT = PA < - - - - >$$

But I have a prejudice, or deeply held belief, that this is morally wrong!

$$\mathcal{K} : \begin{array}{c} \nearrow \searrow \\ \nearrow \end{array} \mapsto \begin{array}{c} \nearrow \nearrow \\ \nearrow \end{array}$$

Crossing the Crossings



My moment of reckoning.

~~Boden and Manturov~~ Boden's $VG\mathbb{X}(K)$

$$\begin{array}{ccc} z \nearrow w \rightarrow z = xyx^{-1} & w \nearrow z \rightarrow z = x^{-1}yx & z \nearrow w \rightarrow z = q^{-1}yq \\ x \nearrow y \quad w = x & y \nearrow x \quad w = x & x \nearrow y \quad w = qxq^{-1} \end{array}$$

Manturov's $vB_n \rightarrow \text{Aut}(F(x_1, \dots, x_n, q))$:

$$\begin{array}{ccc} \nearrow \times \times \rightarrow \left\{ \begin{array}{l} x_i \mapsto x_i x_{i+1} x_i^{-1} \\ x_{i+1} \mapsto x_i \end{array} \right. & \times \times \rightarrow \left\{ \begin{array}{l} x_i \mapsto q x_{i+1} q^{-1} \\ x_{i+1} \mapsto q^{-1} x_i q \end{array} \right. \end{array}$$

Easy resolution: setting $y_i = q^i x_i$, we find that this is equivalent to

But where does this come from?

$$wT = vT / OC \text{ (but not } UC) \quad A$$

$$(v+1w)T = vT_{*+1q} / OC \text{ on } q$$

References.

[BND] D. Bar-Natan and Z. Dancso, *Finite Type Invariants of W-Knotted Objects I: Braids, Knots and the Alexander Polynomial*, weβ/WKO1, arXiv:1405.1956; and II: Tangles and the Kashiwara-Vergne Problem, weβ/WKO2, arXiv:1405.1955.

And Boden, Manturov, Keßler, Brandenbur

$$\mathcal{K}_w : vT \rightarrow wT$$

$$\mathcal{K} : vT \rightarrow (v+1w)T$$

$$\mathcal{K} : vB_n \rightarrow wB_{n+1} \text{ (really, } B_nv + 1w)$$

Claims 1. \mathcal{K} is well-defined.

2. \mathcal{K} vanishes on uT
3. \mathcal{K} does not satisfy OC .

$$4. VG(K) = \pi_1(\mathcal{K}(K))$$

$$5. \mathcal{N}(B) = \emptyset(\mathcal{K}(B))$$

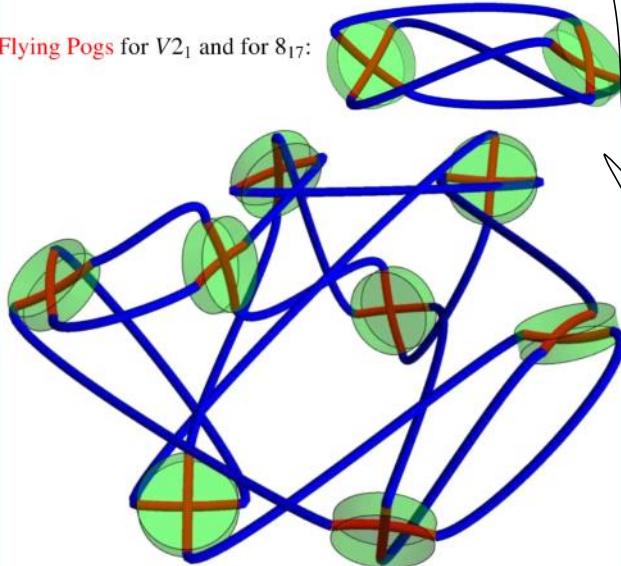
Conjecture: Brandenbur

Even better, can pull back any int from

Prejudices should always be re-evaluated

a rainbow box

Flying Pogs for $V2_1$ and for 8_{17} :



A: $\overline{x} \overline{y} = \overline{\overline{x} \overline{y}}$

T_I is defined on WT , Artins ϕ is
defined on WB_n