

# PolyPoly on 150708

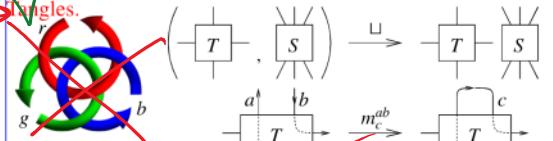
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Dror Bar-Natan: Talks: Qinhuangdao-1507:  
 $\omega\beta := \text{http://www.math.toronto.edu/~drorbn/Talks/Qinhuangdao-1507/}$

**Abstract.** The value of things is inversely correlated with their computational complexity. “Real time” machines, such as our brains, only run linear time algorithms, and there’s still a lot we don’t know. Anything we learn about things doable in linear time is truly valuable. Polynomial time we can in-practice run, even if we have to wait; these things are still valuable. Exponential time we can play with, but just a little, and exponential things must be beautiful or philosophically compelling to deserve attention. Values further diminish and the aesthetic-or-philosophical bar further rises as we go further slower, or un-computable, or ZFC-style intrinsically infinite, or large-cardinalish, or beyond.

I will explain some things I know about polynomial time knot polynomials and explain where there’s more, within reach.



Why Tangles?

- Finitely presented.
- Divide and conquer proofs and computations
- “Algebraic Knot Theory”: If  $K$  is ribbon,

$$z(K) \in \{cl_2(\zeta) : cl_1(\zeta) = 1\}.$$

(Genus and crossing number are also definable properties.)

Faster is better, leaner is meaner!

**Theorem 1.**  $\exists!$  an invariant  $z_0 : \{\text{pure framed } S\text{-component tangles}\} \rightarrow \Gamma_0(S) := R \times M_{S \times S}(R)$ , where  $R = R_S = \mathbb{Z}((T_a)_{a \in S})$  is the ring of rational functions in  $S$  variables, intertwining

$$\begin{aligned} 1. \left( \begin{array}{c|cc} \omega_1 & S_1 & \omega_2 & S_2 \\ \hline S_1 & A_1 & S_2 & A_2 \end{array} \right) &\xrightarrow{\sqcup} \frac{\omega_1 \omega_2}{S_1} \begin{array}{c|cc} S_1 & S_2 \\ A_1 & 0 \\ \hline S_2 & 0 & A_2 \end{array}, \\ 2. \left( \begin{array}{c|ccc} \omega & a & b & S \\ \hline a & \alpha & \beta & \theta \\ b & \gamma & \delta & \epsilon \\ S & \phi & \psi & \Xi \end{array} \right) &\xrightarrow{\mu c^{ab}} \left( \begin{array}{c|cc} \mu \omega & c & S \\ \hline c & \gamma + \alpha \delta / \mu & \epsilon + \delta \theta / \mu \\ S & \phi + \alpha \psi / \mu & \Xi + \psi \theta / \mu \end{array} \right)_{T_a T_b \rightarrow T_c} \end{aligned}$$

and satisfying  $(\mathbf{1}_a, a \nearrow_b, b \nearrow_a) \xrightarrow{z_0} \left( \begin{array}{c|cc} 1 & a & b \\ a & 1 & 1 - T_a^{\pm 1} \\ b & 0 & T_a^{\pm 1} \end{array} \right)$ .

In Addition • The matrix part is just a stitching formula for Burau/Gassner [LD, KLW, CT].



•  $K \mapsto \omega$  is Alexander, mod units.

•  $L \mapsto (\omega, A) \mapsto \omega \det'(A - I)/(1 - T')$  is the MVA, mod units.

• The “fastest” Alexander algorithm.

• There are also formulas for strand deletion, reversal, and doubling.

• Every step along the computation is the invariant of something.

• Extends to and more naturally defined on v/w-tangles.

• Fits in one column, including propaganda & implementation.

**Implementation key idea:**

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 $(\omega, A = (\alpha_{ab})) \leftrightarrow \omega\beta/\text{Demo}$ 
 $(\omega, \lambda = \sum \alpha_{ab} t_a h_b) \quad \quad \quad$ 
 $\text{Module}[\{\alpha, \beta, \gamma, \delta, \theta, \epsilon, \phi, \psi, \Xi, \mu\},$ 
 $\text{Collect}[\{\omega, \lambda\}] := \text{Module}[\{\omega, \lambda\},$ 
 $\text{Format}[\{\omega, \lambda\}, \text{Module}[\{\omega, \lambda\}, \text{Factor}[4]]];$ 
 $\text{S} = \text{UnionCases}[\{\omega, \lambda\}, (b | t_a)_a = a, \#];$ 
 $\text{M} = \text{Outer}[\text{Factor}[\alpha_{ab}], \mathbf{S}, \mathbf{S}];$ 
 $\text{M} = \text{Prepend}[\text{M}, \mathbf{t}_a \& \theta / \mathbf{S} // \text{Transpose}];$ 
 $\text{M} = \text{Prepend}[\text{M}, \text{Prepend}[\text{M}, \mathbf{t}_a \& \theta / \mathbf{S}, \omega]];$ 
 $\text{M} // \text{MatrixForm};$ 

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## Polynomial Time Knot Polynomials, A

**Meta-Associativity**  $\zeta = \Gamma[\omega, \{t_1, t_2, t_3, t_8\}, \begin{pmatrix} \alpha_{11} & \alpha_{12} & \alpha_{13} & \theta_1 \\ \alpha_{21} & \alpha_{22} & \alpha_{23} & \theta_2 \\ \alpha_{31} & \alpha_{32} & \alpha_{33} & \theta_3 \\ \phi_1 & \phi_2 & \phi_3 & \Xi \end{pmatrix}, \{\mathbf{h}_1, \mathbf{h}_2, \mathbf{h}_3, \mathbf{h}_8\}]$ , **Runs**

$$(\zeta // m_{12+1} // m_{13+1}) = (\zeta // m_{23+2} // m_{12+1})$$

True R3 ... divide and conquer!

$$(Rm_{51} Rm_{62} Rp_{34} // m_{14+1} // m_{25+2} // m_{36+3}, Rp_{61} Rp_{24} Rm_{35} // m_{14+1} // m_{25+2} // m_{36+3})$$

$$\left[ \begin{array}{c|ccc} 1 & h_1 & h_2 & h_3 \\ \hline t_1 & \frac{T_3}{T_2} & 0 & 0 \\ t_2 & \frac{-1+T_2}{T_2} & \frac{1}{T_3} & 0 \\ t_3 & \frac{-1+T_3}{T_2} & \frac{-1+T_3}{T_3} & 1 \end{array} \right], \left[ \begin{array}{c|ccc} 1 & h_1 & h_2 & h_3 \\ \hline t_1 & \frac{T_3}{T_2} & 0 & 0 \\ t_2 & \frac{-1+T_2}{T_2} & \frac{1}{T_3} & 0 \\ t_3 & \frac{-1+T_3}{T_2} & \frac{-1+T_3}{T_3} & 1 \end{array} \right]$$

$$\mathbf{z} = Rm_{12,1} Rm_{27} Rm_{83} Rm_{4,11} Rp_{16,5} Rp_{6,13} Rp_{14,9} Rp_{10,15};$$

Do [ $\mathbf{z} = \mathbf{z} // m_{1k+1}$ , {k, 2, 16}]; 817

$$\mathbf{z} = \left( 11 - \frac{1}{T_1^3} + \frac{4}{T_1^2} - \frac{8}{T_1} - 8 T_1 + 4 T_1^2 - T_1^3 \right) h_1$$

$$= \left( \begin{array}{c|ccc} + & + & + & + \\ \hline + & - & - & - \\ - & + & - & - \\ - & - & + & - \\ - & - & - & + \end{array} \right) h_1$$

**Closed Components.** The Halacheva trace satisfies  $m_c^{ab} // \text{tr}_c = m_{ca}^{ba} // \text{tr}_c$  and computes the MVA for all links in the atlas, but its domain is not understood: in range

$$\text{tr}_c[\Gamma[w_\alpha, \lambda_\alpha]] := \text{Module}[\{\alpha, \theta, \psi, \Xi\},$$

$$\left( \begin{array}{c|c} \alpha & \theta \\ \psi & \Xi \end{array} \right) = \left( \begin{array}{c|c} \partial_{t_\alpha} h_0 & \lambda \\ \partial_{t_\alpha} \lambda & \lambda \end{array} \right) / . (t | h)_\alpha = 0;$$

$$\Gamma[\omega (1 - \alpha), \Xi + \psi \theta + (1 - \alpha)] // \text{RCollect};$$

$$(\zeta // m_{12+1} // \text{tr}_c) = (\zeta // m_{23+2} // \text{tr}_c)$$

True

**Weaknesses.** •  $m_c^{ab}$  and  $\text{tr}_c$  are non-linear. • The product  $\omega A$  is always Laurent, but my current proof takes induction with exponentially many conditions. • I still don’t understand  $\text{tr}_c$ . • I still don’t understand “unitarity”. breeze through

v-Tangles. sort it out

$$vT := \text{PA} \langle \cancel{\times}, \cancel{\times} \rangle / \cancel{\times} = \text{VR1} = \text{VR2} = \text{VR3}$$

$$= \text{CA} \langle \cancel{\times} \rangle / \cancel{\times} = \text{VR1} = \text{VR2} = \text{VR3}$$

Flying pogs

Let  $\mathcal{I} := \langle \cancel{\times} - \cancel{\times} \rangle$ . Then  $\mathcal{A}^\vee := \prod \mathcal{I}^n / \mathcal{I}^{n+1} = \text{“universal } \mathcal{U}(Dg)^{\otimes S} \text{”}$

$$\langle \cancel{\times} - \cancel{\times} \rangle / \cancel{\times} = \text{STU}_1 = \text{STU}_2 = \text{STU}_3 = \text{HX}$$

for sources no sinks, AS vertices, grabbed by half # of vertices internally acyclic.

**Theorem.** [EK, En] There exists a homomorphic expansion  $Z : vT \rightarrow \mathcal{A}^\vee$ .

Too hard! Let's look for meta-monoid quotients. sort it out

Pixellization issues.

sort it out  
blue  
brace

purple  
intensity  
edges

red

Add some  $\text{tr}_c$  computations? The Borromean  $\cancel{\times}$

2. The relationship w/  $A^w$  w/ an explicit map  $\cancel{\times}$

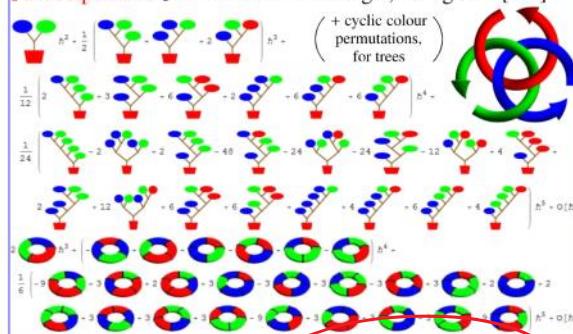
2. The relationship w/  $A^w$  w/ an explicit map! X
3. A picture from Gompf ✓

and the "v-gatient".

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**Theorem 2** [BND].  $\exists!$  a homomorphic expansion, aka a homomorphic universal finite type invariant  $Z^w$  of pure w-tangles.  $Z^w := \log Z^w$  takes values in  $FL(S)^S \times CW(S)$ .

$z$  is computable!  $z$  of the Borromean tangle, to degree 5 [BN]:



**Proposition** [BN]. Modulo all relations that universally hold for the 2D non-Abelian Lie algebra and after some changes-of-variable,  $z^w$  reduces to  $z_0$ .

The "Jones" patient

### Polynomial Time Knot Polynomials, B

**Definition.** (Compare [BNS, BN]) A meta-monoid is a functor  $M$ : (finite sets, injections)  $\rightarrow$  (sets) (think " $M(S)$  is quantum  $G^S$ ", for  $G$  a group) along with natural operations  $*: M(S_1) \times M(S_2) \rightarrow M(S_1 \sqcup S_2)$  whenever  $S_1 \cap S_2 = \emptyset$  and  $m_c^{ab}: M(S \sqcup \{a, b\}) \rightarrow M(S \sqcup \{c\})$  whenever  $a \neq b \notin S$  and  $c \notin S$ , such that

$$\text{meta-associativity: } m_a^{ab} // m_a^{ac} = m_b^{bc} // m_a^{ab}$$

$$\text{meta-locality: } m_c^{ab} // m_f^{de} = m_f^{de} // m_c^{ab}$$

and, with  $\epsilon_b = M(S \hookrightarrow S \sqcup \{b\})$ ,

$$\text{meta-unit: } \epsilon_b // m_a^{ab} = Id = \epsilon_b // m_a^{ba}.$$

**Claim.** Pure virtual tangles  $P\Gamma$  form a meta-monoid.

**Theorem.**  $S \mapsto \Gamma_0(S)$  is a meta-monoid and  $z_0: P\Gamma \rightarrow \Gamma_0$  is a morphism of meta-monoids.

**Strong Conviction.** There exists an extension of  $\Gamma_0$  to a bigger meta-monoid  $\Gamma_{01}(S) = \Gamma_0(S) \times \Gamma_1(S)$ , along with an extension of  $z_0$  to  $z_{01}: P\Gamma \rightarrow \Gamma_{01}$ , with (more or less)

$$\Gamma_1(S) =$$

Furthermore, upon reducing to a single variable everything is polynomial size and polynomial time.

Furthermore,  $\Gamma_{01}$  is given using a "meta-2-cocycle over  $\Gamma_0$ ". In addition to  $m_c^{ab} \rightarrow m_{0c}^{ab}$ , there are  $R_S$ -linear  $m_{1c}^{ab}: \Gamma_1(S \sqcup \{a, b\}) \rightarrow \Gamma_1(S \sqcup \{c\})$  a meta-right-action  $\alpha^{ab}: \Gamma_1(S) \times \Gamma_0(S) \rightarrow \Gamma_1(S)$   $R_S$ -linear in the first variable, and an order 1 differential operator (over  $R_S$ )  $\rho_c^{ab}: \Gamma_0(S \sqcup \{a, b\}) \rightarrow \Gamma_1(S \sqcup \{c\})$  such that

... and then

$$(\zeta_0, \zeta_1) // m_c^{ab} = (\zeta_0 // m_{0c}^{ab}, (\zeta_1, \zeta_0) // \alpha^{ab} // m_{1c}^{ab} + \zeta_0 // \rho_c^{ab})$$

The 1-co-gatient

everything should be polynomial...

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Let's talk about China, America, Taiwan, economy, ecology, religion, democracy, censorship, and all else.



"God created the knots, all else in topology is the work of mortals."

Leopold Kronecker (modified)



[www.katlas.org](http://www.katlas.org)

Really add:

Something about one-co, abc calculus,

"I'm 2D and I'm a friend of Vaughan Jones")

