


Trees and Wheels and Balloons and Hoops

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ωεβ:=http://www.math.toronto.edu/~drorbn/Talks/Toronto-1303



15 Minutes on Algebra

Let T be a finite set of "tail labels" and H a finite set of "head labels". Set

$$M_{1/2}(T; H) := FL(T)^H,$$

" H -labeled lists of elements of the degree-completed free Lie algebra generated by T ".

$$FL(T) = \left\{ 2t_2 - \frac{1}{2}[t_1, [t_1, t_2]] + \dots \right\} / \left(\begin{array}{c} \text{anti-symmetry} \\ \text{Jacobi} \end{array} \right)$$


... with the obvious bracket.

$$M_{1/2}(u, v; x, y) = \left\{ \left(x \rightarrow \begin{array}{c} u \quad v \\ \diagdown \quad \diagup \\ x \end{array}, y \rightarrow \begin{array}{c} v \quad u \\ \diagdown \quad \diagup \\ y \end{array}, -\frac{2}{3} \begin{array}{c} u \quad u \quad v \\ \diagdown \quad \diagup \quad \diagup \\ y \end{array} \right) \dots \right\}$$

15 Minutes on Topology

$\mathcal{K}^{bh}(T; H)$.


T



balloons / tails

ribbon embeddings $\rightarrow \mathbb{R}^4 / S^4$

H



hoops / heads

Operations $M_{1/2} \rightarrow M_{1/2}$ *is aim!* newspeak!

Tail Multiply $tm_w^{uv} : M_{1/2} \rightarrow M_{1/2} \lambda \mapsto \lambda \parallel (u, v \rightarrow w)$
satisfies "meta-associativity", $tm_w^{uv} \parallel tm_w^{uv} = tm_w^{uv} \parallel tm_w^{uv}$.

Head Multiply $hm_z^{xy} : M_{1/2} \rightarrow M_{1/2} \lambda \mapsto (\lambda \{x, y\}) \cup (z \rightarrow \dots)$
bch(λ_x, λ_y), where

$$\text{bch}(\alpha, \beta) := \log(e^\alpha e^\beta) = \alpha + \beta + \frac{[\alpha, \beta]}{2} + \frac{[\alpha, [\alpha, \beta]] + [[\alpha, \beta], \beta]}{12} + \dots$$

satisfies $\text{bch}(\text{bch}(\alpha, \beta), \gamma) = \log(e^{\text{bch}(\alpha, \beta)} e^\gamma) = \text{bch}(\alpha, \text{bch}(\beta, \gamma))$
and hence meta-associativity, $hm_z^{xy} \parallel hm_z^{xy} = hm_z^{xy} \parallel hm_z^{xy}$.

Tail by Head Action $tha^{ux} : M_{1/2} \rightarrow M_{1/2} \lambda \mapsto \lambda \parallel RC_u^{\lambda_x}$
where $C_u^{-\gamma} : FL \rightarrow FL$ is the substitution $u \rightarrow e^{-\gamma} u e^\gamma$, or more precisely,

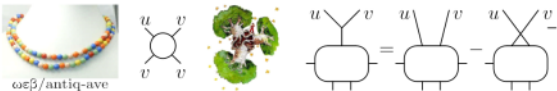
$$C_u^{-\gamma} : u \rightarrow e^{-\text{ad } \gamma}(u) = u - [\gamma, u] + \frac{1}{2}[\gamma, [\gamma, u]] - \dots,$$

and RC_u^γ is the inverse of that. Note that $C_u^{\text{bch}(\alpha, \beta)} = C_u^{\alpha \parallel RC_u^\beta} \parallel C_u^\alpha$ and hence "meta $u^{xy} = (u^x)^y$ ",

$$hm_z^{xy} \parallel tha^{uz} = tha^{ux} \parallel tha^{vy} \parallel hm_z^{xy},$$



and $tm_w^{uv} \parallel C_w^\gamma \parallel tm_w^{uv} = C_w^\gamma \parallel RC_w^{-\gamma} \parallel C_w^\gamma \parallel tm_w^{uv}$ and hence "meta $(uv)^x = u^x v^x$ ", $tm_w^{uv} \parallel tha^{wx} = tha^{ux} \parallel tha^{vy} \parallel tm_w^{uv}$.



Wheels. Let $M(T; H) := M_{1/2}(T; H) \times CW(T)$, where $CW(T)$ is the (completed graded) vector space of cyclic words on T , or equally well, on $FL(T)$:



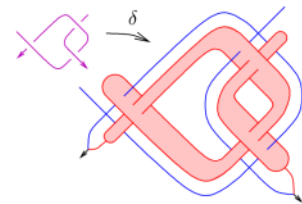
$\omega \in \beta / \text{antiq-ave}$

Examples.


$\epsilon_x : x \rightarrow$  $\epsilon_u : u \rightarrow$ 

$\rho_{ux}^+ : u \rightarrow x$  $\rho_{ux}^- : u \rightarrow x$ 







"the generators"



More on δ



Satisfies R123, VR123, D, and

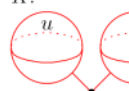
OC:  \leftrightarrow  as  \leftrightarrow  yet not UC:  \leftrightarrow 

- δ injects u-Knots into \mathcal{K}^{bh} (likely u-tangles too).
- δ maps v-tangles to \mathcal{K}^{bh} ; the kernel is as above, and **conjecturally**, that's all. Allowing punctures and cuts, δ is onto.

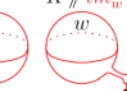
Operations Punctures & Cuts Connected Sums.

If X is a space, $\pi_1(X)$ is a group, $\pi_2(X)$ is an Abelian group, and π_1 acts on π_2 .

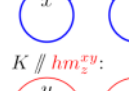
K :



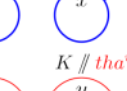
$K \parallel tm_w^{uv}$:




$K \parallel hm_z^{xy}$:



$K \parallel tha^{ux}$:

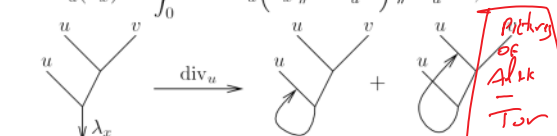




Operations. On $M(T; H)$, define tm_w^{uv} and hm_z^{xy} as before, and tha^{ux} by adding some J-spice:

$$(\lambda; \omega) \mapsto (\lambda, \omega + J_u(\lambda_x)) \parallel RC_u^{\lambda_x},$$

where $J_u(\lambda_x) := \int_0^1 ds \text{div}_u(\lambda_x \parallel RC_u^{s\lambda_x}) \parallel C_u^{-s\lambda_x}$, and

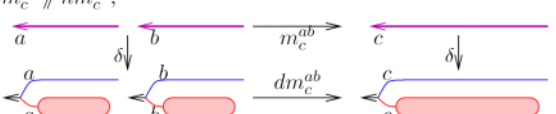


Anything of all k - Tor

Theorem Green. All green identities still hold.

Merge Operation. $(\lambda_1; \omega_1) * (\lambda_2; \omega_2) := (\lambda_1 \cup \lambda_2; \omega_1 + \omega_2)$.

Tangle concatenations $\rightarrow \pi_1 \times \pi_2$. With $dm_c^{ab} := tha^{ab} \parallel tm_c^{ab} \parallel hm_c^{ab}$,



Moral. To construct an M -valued invariant ζ of (v-)tangles, and nearly an invariant on \mathcal{K}^{bh} , it is enough to declare ζ on the generators, and verify the relations that δ satisfies.

→

→

→

✓

→

It is the universal solution to a topological problem and it has many siblings (who talk to each other). It is explicitly computable. Its target space is in itself a space of "universal formulas in Lie algebras" (that's "the miracle"). It seems to be a complete(?) evaluation a certain gauge theory. It is related to a deep conjecture in Lie theory proven by Alekseev and Meinrenken. It has even-better-computable

✓

specializations, including one which is an "ultimate Alexander invariant". And plenty of work remains to be done.

Trees and Wheels and Balloons and Hoops: Why I Care

The Invariant ζ . Set $\zeta(\epsilon_x) = (x \rightarrow 0; 0)$, $\zeta(\epsilon_u) = ((); 0)$, and

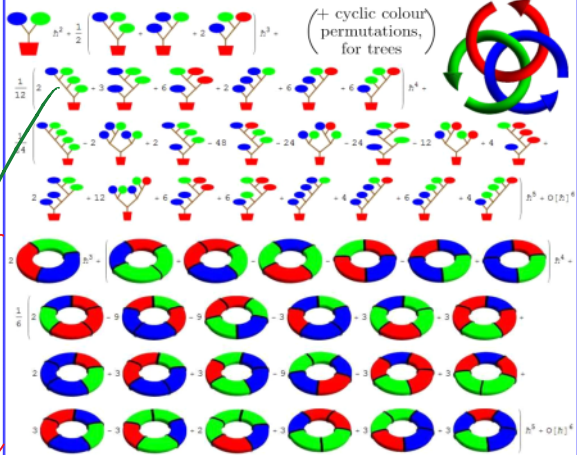
$$\zeta: \begin{array}{c} \text{tree} \\ \downarrow \\ \text{matrix} \end{array} \quad \begin{array}{c} \text{tree} \\ \downarrow \\ \text{matrix} \end{array}$$

Theorem. ζ is (the log of) the unique homomorphic universal finite type invariant. *on knot*

... and it is the tip of an iceberg ...

put some of handouts here?

ζ is computable! ζ of the Borromean tangle, to degree 5:



smaller wheels, same size

Tensorial Interpretation. Let \mathfrak{g} be a finite dimensional Lie algebra (any!). Then there's $\tau: FL(T) \rightarrow \text{Fun}(\oplus_T \mathfrak{g} \rightarrow \mathfrak{g})$ and $\tau: CW(T) \rightarrow \text{Fun}(\oplus_T \mathfrak{g})$. Together, $\tau: M(T; H) \rightarrow \text{Fun}(\oplus_T \mathfrak{g} \rightarrow \oplus_H \mathfrak{g})$, and hence

$$e^\tau: M(T; H) \rightarrow \text{Fun}(\oplus_T \mathfrak{g} \rightarrow \mathcal{U}^{\otimes H}(\mathfrak{g})).$$

ζ and BF Theory. (See Cattaneo-Rossi, arXiv:math-ph/0210037) Let A denote a \mathfrak{g} -connection on S^4 with curvature F_A , and B a \mathfrak{g}^* -valued 2-form on S^4 . For a hoop γ_x , let $\text{hol}_{\gamma_x}(A) \in \mathcal{U}(\mathfrak{g})$ be the holonomy of A along γ_x . For a ball γ_u , let $\mathcal{O}_{\gamma_u}(B) \in \mathfrak{g}^*$ be the integral of B (transported via A to ∞) on γ_u .



Cattaneo

Loose Conjecture. For $\gamma \in \mathcal{K}(T; H)$,

$$\int \mathcal{D}A \mathcal{D}B e^{\int B \wedge F_A} \prod_u e^{\mathcal{O}_{\gamma_u}(B)} \otimes_x \text{hol}_{\gamma_x}(A) = e^\tau(\zeta(\gamma)).$$

That is, ζ is a complete evaluation of the BF TQFT.

Issues. How exactly is B transported via A to ∞ ? How does the ribbon condition arise? Or if it doesn't, could it be that ζ can be generalized??

The β quotient, 1. • Arises when \mathfrak{g} is the 2D non-Abelian Lie algebra. *move to next column.*

• Arises when reducing by relations satisfied by the weight system of the Alexander polynomial.



"God created the knots, all else in topology is the work of mortals." Leopold Kronecker (modified)



www.katlas.org The Knot Atlas

Paper in progress: $\omega\epsilon\beta/kbh$

Class next year: $\omega\epsilon\beta/1350$

The β quotient. ~~X~~ Let $R = \mathbb{Q}[\{c_u\}_{u \in T}]$ and $L_\beta := R \otimes T$ with central R and with $[u, v] = c_u v - c_v u$ for $u, v \in T$. Then $FL \rightarrow L_\beta$ and $CW \rightarrow R$. Under this,

$$\mu \rightarrow (\bar{\lambda}; \omega) \text{ with } \omega = \sum_{u \in T} \lambda_{ux} u x, \quad \lambda_{ux}, \omega \in R,$$

$$\text{bch}(u, v) \rightarrow \frac{c_u + c_v}{e^{c_u + c_v} - 1} \left(\frac{e^{c_u} - 1}{c_u} u + e^{c_v} \frac{e^{c_v} - 1}{c_v} v \right),$$

if $\lambda = \sum \lambda_v v$ then with $c_\lambda := \sum \lambda_v c_v$,

$$u // RC^\lambda = \left(1 + c_u \lambda_u \frac{e^{c_\lambda} - 1}{c_\lambda} \right)^{-1} \left(e^{c_\lambda} u - c_u \frac{e^{c_\lambda} - 1}{c_\lambda} \sum_{v \neq u} \lambda_v v \right)$$

$\text{div}_u \lambda = c_u \lambda_u$, and the ODE for J integrates to

$$J_u(\lambda) = \log \left(1 + \frac{e^{c_\lambda} - 1}{c_\lambda} c_u \lambda_u \right),$$

so ζ is formula-computable to all orders! Can we simplify?

Repackaging. Given $((x: \lambda_{ux}); \omega)$, set $c_x := \sum_v c_v \lambda_{vx}$, replace $\lambda_{ux} \rightarrow \alpha_{ux} := c_u \lambda_{ux} \frac{e^{c_x} - 1}{c_x}$ and $\omega \rightarrow \omega_{\alpha}$, use $t_u = e^{c_u}$, and write α_{ux} as a matrix. Get " β calculus".

β Calculus. Let $\beta(H, T)$ be

$$\left\{ \begin{array}{c|ccc} \omega & x & y & \dots \\ \hline u & \alpha_{ux} & \alpha_{uy} & \dots \\ v & \alpha_{vx} & \alpha_{vy} & \dots \\ \vdots & \cdot & \cdot & \cdot \end{array} \middle| \begin{array}{l} \omega \text{ and the } \alpha_{ux} \text{'s are} \\ \text{rational functions in} \\ \text{variables } t_u, \text{ one for} \\ \text{each } u \in T. \end{array} \right\}$$



Selman Akbulut

$$tm_{uv}^{xy}: \begin{array}{c|ccc} \omega & \dots & \dots & \dots \\ \hline u & \alpha & \alpha + \beta & \dots \\ v & \beta & \gamma & \dots \\ \vdots & \cdot & \cdot & \cdot \end{array} \mapsto \begin{array}{c|ccc} \omega & \dots & \dots & \dots \\ \hline \omega_1 & H_1 & \omega_2 & H_2 \\ \hline T_1 & \alpha_1 & T_2 & \alpha_2 \\ \hline \omega_1 \omega_2 & H_1 & H_2 & \\ \hline T_1 & \alpha_1 & 0 & \\ T_2 & 0 & \alpha_2 & \end{array}$$

$$hm_z^{xy}: \begin{array}{c|ccc} \omega & x & y & \dots \\ \hline \alpha & \beta & \gamma & \dots \end{array} \mapsto \begin{array}{c|ccc} \omega & z & \dots & \dots \\ \hline \alpha & \alpha + \beta + \langle \alpha \rangle \beta & \gamma & \dots \end{array}$$

$$tha^{ux}: \begin{array}{c|ccc} \omega & x & \dots & \omega \epsilon \\ \hline u & \alpha & \beta & u \left(\alpha(1 + \langle \gamma \rangle / \epsilon) \right) \beta(1 + \langle \gamma \rangle / \epsilon) \\ \vdots & \gamma & \delta & \gamma / \epsilon \quad \delta - \gamma \beta / \epsilon \end{array}$$

where $\epsilon := 1 + \alpha$, $\langle \alpha \rangle := \sum_v \alpha_v$, and $\langle \gamma \rangle := \sum_{v \neq u} \gamma_v$, and let

$$R_{ux}^+ := \frac{1}{u} \left| \begin{array}{c} x \\ u \quad t_u - 1 \end{array} \right| \quad R_{ux}^- := \frac{1}{u} \left| \begin{array}{c} x \\ u \quad t_u^{-1} - 1 \end{array} \right|.$$

On long knots, ω is the Alexander polynomial!

Why bother? (1) An ultimate Alexander invariant: Manifestly polynomial (time and size) extension of the (multivariable) Alexander polynomial to tangles. Every step of the computation is the computation of the invariant of some topological thing (no fishy Gaussian elimination!). If there should be an Alexander invariant ~~more~~ algebraic categorification, it is this one. See also $\omega\epsilon\beta/\text{regina}$, $\omega\epsilon\beta/\text{gwu}$.



Why bother? (2) Related to A-T, K-V, and E-K, should have vast generalization beyond w-knots and the Alexander polynomial. See also $\omega\epsilon\beta/\text{wko}$, $\omega\epsilon\beta/\text{caen}$, $\omega\epsilon\beta/\text{swiss}$.

move right.

more near handouts

Handouts to insert:

Tennessee, SwissKnots/Strasbourg, Bonn