

Switch to the KBH labeling convention:

x, y, z	For heads	Blue!	✓
u, v, w	For tails	Red!	✓
a, b, c	For ungendered	Purple!	not done

Find a place for the stepping stones image. ✓

Say "reluctant algebraist"? ✗

Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, I

Dror Bar-Natan at the Newton Institute, January 2013.

<http://www.math.toronto.edu/~drorbn/Talks/Newton-1301>



Abstract. I will define “meta-groups” and explain how one specific meta-group, which in itself is a “meta-bicrossed-product”, gives rise to an “ultimate Alexander invariant” of tangles, that contains the Alexander polynomial (multivariable, if you wish), has extremely good composition properties, is evaluated in a topologically meaningful way, and is least-wasteful in a computational sense. If you believe in categorification, that’s a wonderful playground. This will be a repeat of a talk I gave in Regina in August 2012 and in a number of other places, and I plan to repeat it a good further number of places. Though here at the Newton Institute I plan to make the talk a bit longer, giving me more time to give some further fun examples of meta-structures, and perhaps I will learn from the audience that these meta-structures should really be called something else.

Work is closely related to work by Le Dimet (Comment. Math. Helv. **67** (1992) 306-315), Kirk, Livingston and Wang (arXiv:math/9806035) and Cimasoni and Turaev (arXiv:math.GT/0406269).

Alexander Issues.

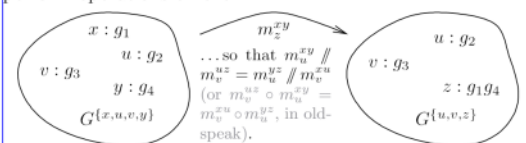
- Quick to compute, but computation departs from topology.
- Extends to tangles, but at an exponential cost.
- Hard to categorify.

Idea. Given a group G and two “YB” pairs $R^\pm = (g_\pm^+, g_\pm^-) \in G^2$, map them to crossings and “multiply along”, so that

$$\begin{array}{c} \text{Crossing} \xrightarrow{Z} \text{Crossing} \\ \left(\begin{array}{c} g_o^+ g_o^+ g_o^- g_o^- g_o^+ g_o^+ \\ g_u g_o \end{array} \right) \end{array}$$

This Fails! R2 implies that $g_o^\pm g_o^\mp = e = g_u^\pm g_u^\mp$ and then R3 implies that g_o^\pm and g_u^\pm commute, so the result is a simple counting invariant.

A Group Computer. Given G , can store group elements and perform operations on them:



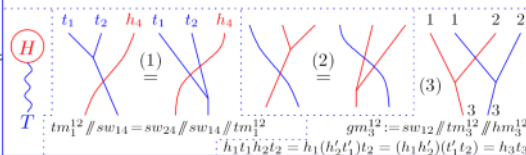
Also has S_x for inversion, e_x for unit insertion, d_x for register deletion, Δ_{xy}^z for element cloning, ρ_y^x for renamings, and $(D_1, D_2) \mapsto D_1 \cup D_2$ for merging, and many obvious composition axioms relating those.

A Meta-Group. Is a similar “computer”, only its internal structure is unknown to us. Namely it is a collection of sets $\{G_\gamma\}$ indexed by all finite sets γ , and a collection of operations m_γ^{xy} , S_x , e_x , d_x , Δ_{xy}^z (sometimes), ρ_γ^x and \cup , satisfying the exact same linear properties.

Example 1. The non-meta example, $G_\gamma := G^\gamma$.
Example 2. $G_\gamma := M_{\gamma \times \gamma}(\mathbb{Z})$, with simultaneous row and column operations, and “block diagonal” merges. Here if $P = \begin{pmatrix} x & a & b \\ y & c & d \end{pmatrix}$ then $d_y P = (x : a)$ and $d_x P = (y : d)$ so $\{d_y P\} \cup \{d_x P\} = \begin{pmatrix} x & a & 0 \\ y & 0 & d \end{pmatrix} \neq P$. So this G is truly meta.

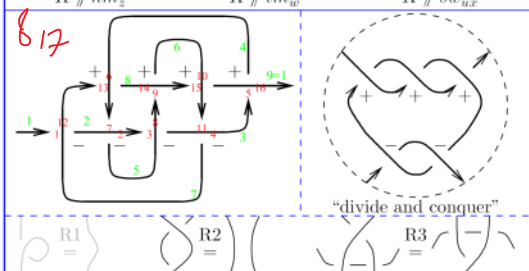
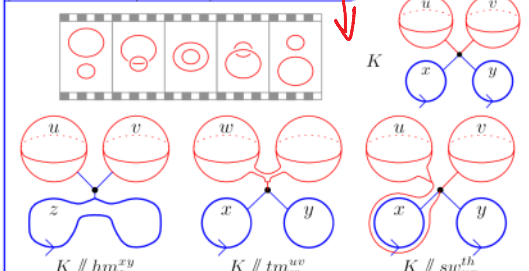
Claim. From a meta-group G and YB elements $R^\pm \in G_2$ we can construct an Alexander invariant.

Bicrossed Products. If $G = HT$ is a group presented as a product of two of its subgroups, with $H \cap T = \{e\}$, then also $G = TH$ and G is determined by H , T , and the “swap” map $sw^{th} : (t, h) \mapsto (h', t')$ defined by $th = h't'$. The map sw satisfies (1) and (2) below; conversely, if $sw : T \times H \rightarrow H \times T$ satisfies (1) and (2) (+ lesser conditions), then (3) defines a group structure on $H \times T$, the “bicrossed product”.



Asd "Comment" Summary

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A Standard Alexander Formula. Label the arcs 1 through $(n+1) = 1$, make an $n \times n$ matrix as below, delete one row and one column, and compute the determinant:

$$\begin{array}{c} \begin{array}{ccc} c & b & \\ \swarrow & \nearrow & \\ a & & \end{array} \rightarrow \begin{array}{ccc} a & b & c \\ c & -1 & 1-X & X \end{array} \\ \begin{array}{ccc} b & c & \\ \swarrow & \nearrow & \\ a & & \end{array} \rightarrow \begin{array}{ccc} a & b & c \\ c & -X & X-1 & 1 \end{array} \end{array}$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & X-1 & 0 & -X \\ -1 & X & 0 & 0 & 0 & 0 & 1-X & 0 \\ 0 & -1 & X & 0 & 1-X & 0 & 0 & 0 \\ X-1 & 0 & -X & 1 & 0 & 0 & 0 & 0 \\ 0 & 1-X & 0 & -1 & X & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -X & 1 & 0 & X-1 \\ 0 & 0 & 1-X & 0 & 0 & -1 & X & 0 \\ 0 & 0 & 0 & X-1 & 0 & 0 & -X & 1 \end{pmatrix} \quad [[1 : ; 7, 1 : ; 7]] // \text{Det}$$

$$-1 + 4X - 8X^2 + 11X^3 - 8X^4 + 4X^5 - X^6$$

Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 2

A **Meta-Bicrossed-Product** is a collection of sets $\beta(\eta, \tau)$ and operations tm_z^{xy} , hm_z^{xy} and sw_z^{th} (and lesser ones), such that tm and hm are "associative" and (1) and (2) hold (+ lesser conditions). A meta-bicrossed-product defines a meta-group with $G_\gamma := \beta(\gamma, \gamma)$ and gm as in (3).

Example. Take $\beta(\eta, \tau) = M_{\tau \times \eta}(\mathbb{Z})$ with row operations for the tails, column operations for the heads, and a trivial swap

β Calculus. Let $\beta(\eta, \tau)$ be

$$\left\{ \begin{array}{c|ccc} \omega & h_1 & h_2 & \dots \\ t_1 & \alpha_{11} & \alpha_{12} & \dots \\ t_2 & \alpha_{21} & \alpha_{22} & \dots \\ \vdots & \dots & \dots & \dots \end{array} \middle| \begin{array}{l} h_j \in \eta, t_i \in \tau, \text{ and } \omega \text{ and} \\ \text{the } \alpha_{ij} \text{ are rational functions} \\ \text{in a variable } X \end{array} \right\}$$

$$tm_z^{xy} : \begin{array}{c|ccc} \omega & \dots & & \\ t_x & \alpha & & \\ t_y & \beta & & \\ \vdots & \gamma & & \end{array} \mapsto \begin{array}{c|ccc} \omega & \dots & & \\ t_z & \alpha + \beta & & \\ & \vdots & & \gamma \end{array}, \quad \begin{array}{c|ccc} \omega_1 & \eta_1 & \cup & \omega_2 & \eta_2 \\ \tau_1 & \alpha_1 & & \tau_2 & \alpha_2 \\ & \omega_1 \omega_2 & & \eta_1 & \eta_2 \\ & \tau_1 & & 0 & \alpha_2 \end{array}$$

$$hm_z^{xy} : \begin{array}{c|ccc} \omega & h_x & h_y & \dots \\ \vdots & \alpha & \beta & \gamma \end{array} \mapsto \begin{array}{c|ccc} \omega & h_z & & \dots \\ \vdots & \alpha + \beta + \langle \alpha \rangle \beta & & \gamma \end{array}$$

$$sw_z^{th} : \begin{array}{c|ccc} \omega & h_y & \dots & \omega \epsilon \\ t_x & \alpha & \beta & \mapsto t_x | \alpha(1 + \langle \gamma \rangle / \epsilon) & \beta(1 + \langle \gamma \rangle / \epsilon) \\ \vdots & \gamma & \delta & \vdots & \gamma / \epsilon & \delta - \gamma \beta / \epsilon \end{array}$$

where $\epsilon := 1 + \alpha$ and $\langle c \rangle := \sum_i c_i$, and let

$$R_{xy}^p := \begin{array}{c|cc} 1 & h_x & h_y \\ t_x & 0 & X-1 \\ t_y & 0 & 0 \end{array} \quad R_{xy}^m := \begin{array}{c|cc} 1 & h_x & h_y \\ t_x & 0 & X^{-1}-1 \\ t_y & 0 & 0 \end{array}$$

Theorem. Z^β is a tangle invariant (and more). Restricted to knots, the ω part is the Alexander polynomial. On braids, it is equivalent to the Burau representation. A variant for links contains the multivariable Alexander polynomial.

- Why Happy?**
- Applications to w-knots.
 - Everything that I know about the Alexander polynomial can be expressed cleanly in this language (even if without proof), except HF, but including genus, ribboness, cabling, v-knots, knotted graphs, etc., and there's potential for vast generalizations.
 - The least wasteful "Alexander for tangles" I'm aware of.
 - Every step along the computation is the invariant of something.
 - Fits on one sheet, including implementation & propaganda.

Further meta-monoids. Π (and variants), \mathcal{A} (and quotients), vT, \dots

Further meta-bicrossed-products. Π (and variants), $\vec{\mathcal{A}}$ (and quotients), $M_0, M, \mathcal{K}^{bh}, \mathcal{K}^{rbh}, \dots$

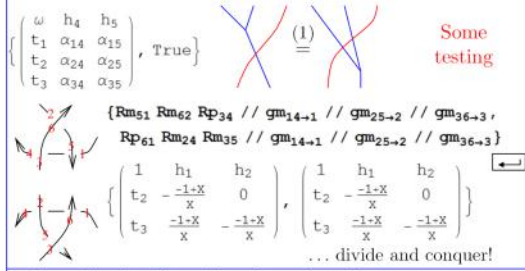
Meta-Lie-algebras. \mathcal{A} (and quotients), \mathcal{S}, \dots

Meta-Lie-bialgebras. $\vec{\mathcal{A}}$ (and quotients), \dots



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mean business!
Simp = Factor[SetAttributes[Collect, Listable];
Collect[#, #2] := #[Simp[#]] &];
Collect[#, #2] := #2[Collect[#, #2]] &];
From[#, #2] := Module[{t, h, s},
t = Union[Cases[#, #2, #, Infinity]];
h = Union[Cases[#, #2, #, Infinity]];
s = Outer[Simp[Coefficient[#, #2, #2] &, h, s]];
PrependTo[#, t, #/s];
h = Prepend[Transpose[h], Prepend[h, #/s]];
MatrixForm[h]];
From[#, #2] := #2[From[#, #2]];
Format[#, #2, StandardForm] := #Form[#, #2];
```

$$\{\beta = B[\omega, \text{Sum}[\alpha_{10+i} t_i h_j, \{i, \{1, 2, 3\}\}, \{j, \{4, 5\}\}]], (\beta // tm_{12+1} // sw_{14}) = (\beta // sw_{24} // sw_{14} // tm_{12+1})\}$$



$$\left\{ \begin{array}{c|ccc} 1 & h_1 & h_2 & \\ t_2 & -\frac{-1-X}{X} & 0 & \\ t_3 & -\frac{-1-X}{X} & -\frac{-1-X}{X} & \end{array} \right\}, \left\{ \begin{array}{c|ccc} 1 & h_1 & h_2 & \\ t_2 & -\frac{-1-X}{X} & 0 & \\ t_3 & -\frac{-1-X}{X} & -\frac{-1-X}{X} & \end{array} \right\}$$

... divide and conquer!

$$\beta = Rm_{12,1} Rm_{27} Rm_{63} Rm_{4,11} Rp_{16,5} Rp_{6,13} Rp_{14,9} Rp_{10,15} \quad 8_{17}$$

$$\begin{array}{c|cccccccc} 1 & h_1 & h_3 & h_5 & h_7 & h_9 & h_{11} & h_{13} & h_{15} \\ t_2 & 0 & 0 & 0 & -\frac{-1-X}{X} & 0 & 0 & 0 & 0 \\ t_4 & 0 & 0 & 0 & 0 & 0 & -\frac{-1-X}{X} & 0 & 0 \\ t_6 & 0 & 0 & 0 & 0 & 0 & 0 & -1+X & 0 \\ t_8 & 0 & -\frac{-1-X}{X} & 0 & 0 & 0 & 0 & 0 & 0 \\ t_{10} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1+X \\ t_{12} & -\frac{-1-X}{X} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ t_{14} & 0 & 0 & 0 & 0 & -1+X & 0 & 0 & 0 \\ t_{16} & 0 & 0 & -1+X & 0 & 0 & 0 & 0 & 0 \end{array}$$

Do[$\beta = \beta // gm_{1k-1}, \{k, 2, 10\}$]; β 8₁₇, cont.

$$\begin{array}{c|cccc} \frac{1}{X} & h_1 & h_{11} & h_{13} & h_{15} \\ t_1 & -\frac{(-1-X)(1-X)}{X} & -(-1+X)(1-X+X^2) & (-1+X)(1-X+X^2) & -1+X \\ t_{12} & -\frac{-1-X}{X} & 0 & 0 & 0 \\ t_{14} & -1+X & \frac{(-1-X)^2(1-X-X^2)}{X} & -\frac{(-1-X)^2(1-X-X^2)}{X} & 0 \\ t_{16} & -\frac{-1-X}{X} & (-1+X)^2 & -\frac{(-1-X)^2}{X} & 0 \end{array}$$

Do[$\beta = \beta // gm_{1k-1}, \{k, 11, 16\}$]; β

$$\left(\frac{-1-4X+8X^2-11X^3+8X^4-4X^5-X^6}{X^3} \right)$$

- A Partial To Do List.**
1. Where does it *more simply* come from?
 2. Remove all the denominators.
 3. How do determinants arise in this context?
 4. Understand links. (i.e., meta-conjugacy classes).
 5. Find the "reality condition".
 6. Do some "Algebraic Knot Theory".
 7. Categorify.
 8. Do the same in other natural quotients of the v/w-story.

