

Monads

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From http://en.wikipedia.org/wiki/Monad_%28category_theory%29:

Formal definition

If C is a **category**, a **monad** on C consists of a functor $T: C \rightarrow C$ together with two **natural transformations**: $\eta: 1_C \rightarrow T$ (where 1_C denotes the identity functor on C) and $\mu: T^2 \rightarrow T$ (where T^2 is the functor $T \circ T$ from C to C). These are required to fulfill the following conditions (sometimes called **coherence conditions**):

- $\mu \circ T\mu = \mu \circ \mu T$ (as natural transformations $T^3 \rightarrow T$);
- $\mu \circ T\eta = \mu \circ \eta T = 1_T$ (as natural transformations $T \rightarrow T$; here 1_T denotes the identity transformation from T to T).

We can rewrite these conditions using following **commutative diagrams**:

$$\begin{array}{ccc}
 T^3 & \xrightarrow{T\mu} & T^2 \\
 \mu T \downarrow & & \downarrow \mu \\
 T^2 & \xrightarrow{\mu} & T
 \end{array}
 \qquad
 \begin{array}{ccc}
 T & \xrightarrow{\eta T} & T^2 \\
 T\eta \downarrow & \searrow & \downarrow \mu \\
 T^2 & \xrightarrow{\mu} & T
 \end{array}$$

See the article on **natural transformations** for the explanation of the notations $T\mu$ and μT , or see below the commutative diagrams not using these notions:

$$\begin{array}{ccc}
 T(T(T(X))) & \xrightarrow{T(\mu_X)} & T(T(X)) \\
 \mu_{T(X)} \downarrow & & \downarrow \mu_X \\
 T(T(X)) & \xrightarrow{\mu_X} & T(X)
 \end{array}
 \qquad
 \begin{array}{ccc}
 T(X) & \xrightarrow{\eta_{T(X)}} & T(T(X)) \\
 T(\eta_X) \downarrow & \searrow & \downarrow \mu_X \\
 T(T(X)) & \xrightarrow{\mu_X} & T(X)
 \end{array}$$

The first axiom is akin to the **associativity** in **monoids**, the second axiom to the existence of an **identity element**. Indeed, a monad on C can alternatively be defined as a **monoid** in the category **End_C** whose objects are the endofunctors of C and whose morphisms are the **natural transformations** between them, with the **monoidal structure** induced by the composition of endofunctors.

See also: [F-algebra](#)

Suppose that (T, η, μ) is a given monad on a category C .

A **T -algebra** (x, h) is an object x of C together with an arrow $h: Tx \rightarrow x$ of C called the **structure map** of the algebra such that the diagrams

$$\begin{array}{ccc}
 T^2x & \xrightarrow{Th} & Tx \\
 \mu_x \downarrow & & \downarrow h \\
 Tx & \xrightarrow{h} & x
 \end{array}
 \quad \text{and} \quad
 \begin{array}{ccc}
 x & \xrightarrow{\eta_x} & Tx \\
 \searrow 1_x & & \downarrow h \\
 & & x
 \end{array}$$

commute.

A morphism $f: (x, h) \rightarrow (x', h')$ of T -algebras is an arrow $f: x \rightarrow x'$ of C such that the diagram

$$\begin{array}{ccc}
 Tx & \xrightarrow{Tf} & Tx' \\
 h \downarrow & & \downarrow h' \\
 x & \xrightarrow{f} & x'
 \end{array}$$

commutes.

The category C^T of T -algebras and their morphisms is called the **Eilenberg-Moore category** or **category of (Eilenberg-Moore) algebras** of the monad T . The forgetful functor $C^T \rightarrow C$ has a left adjoint $C \rightarrow C^T$ taking x to the free algebra (Tx, μ_x) .

Given the monad T , there exists another "canonical" category C_T called the **Kleisli category** of the monad T . This category is equivalent to the *category of free algebras* for the monad T , i. e. the **full subcategory** of C^T whose objects are of the form (Tx, μ_x) , for x an object of C .