

ef:quandle | **Definition 6.5.** A quandle is a set Q with a binary operation $\uparrow : Q \times Q \rightarrow Q$ satisfying the following axioms:

- (1) $\forall x \in Q, x \uparrow x = x.$
- (2) For any fixed $y \in Q$, the map $x \mapsto x \uparrow y$ is invertible²⁹.
- (3) Self-distributivity: $\forall x, y, z \in Q, (x \uparrow y) \uparrow z = (x \uparrow z) \uparrow (y \uparrow z).$

We say that a quandle Q has a unit, or is unital, if there is a distinguished element $1 \in Q$ satisfying the further axiom:

- (4) $\forall x \in Q, x \uparrow 1 = x$ and $1 \uparrow x = 1.$

If G is a group, it is also a (unital) quandle by setting $x \uparrow y := y^{-1}xy$, yet there are many quandles that do not arise from groups in this way.

prop:ProjQ | **Proposition 6.6.** If Q is a unital quandle, $\text{proj}_0 Q$ is one-dimensional and $\text{proj}_{>0} Q$ is a graded Lie algebra generated by $\text{proj}_1 Q$.

Proof. For any algebraic structure A with just one kind of objects, $\text{proj}_0 A$ is one-dimensional, generated by the equivalence class $[x]$ of any single object x . In particular, $\text{proj}_0 Q$ is one-dimensional and generated by $[1]$. Let $\mathcal{I} \subset \mathbb{Q}Q$ be the augmentation ideal of Q . For any $x \in Q$ set $\bar{x} := x - 1 \in \mathcal{I}$. Then \mathcal{I} is generated by the \bar{x} 's, and therefore \mathcal{I}^m is generated by expressions involving the operation \uparrow applied to some m elements of $\bar{Q} := \{\bar{x} : x \in Q\}$ and possibly some further elements $y_i \in Q$. When regarded in $\mathcal{I}^m / \mathcal{I}^{m+1}$, any y_i is such a generating expression can be replaced by 1 , for the difference would be the same expression with y_i replaced by \bar{y}_i , and this is now a member of \mathcal{I}^{m+1} . But for any element $z \in \mathcal{I}$ we have $z \uparrow 1 = z$ and $1 \uparrow z = 0$, so all the 1 's can be eliminated from the expressions generating \mathcal{I}^m . Thus $\text{proj}_{>0} Q$ is generated by \bar{Q} and hence by $\text{proj}_1 Q$.

Let $\Delta : \mathbb{Q}Q \rightarrow \mathbb{Q}Q \otimes \mathbb{Q}Q$ be the linear extension of the operation $x \mapsto x \otimes x$ defined on $x \in Q$, and extend \uparrow to a binary operator $\uparrow_2 : (\mathbb{Q}Q \otimes \mathbb{Q}Q) \otimes (\mathbb{Q}Q \otimes \mathbb{Q}Q) \rightarrow \mathbb{Q}Q \otimes \mathbb{Q}Q$ by using \uparrow twice, to pair the first and third tensor factors and then to pair the second and the fourth tensor factors. With this language in place, the self-distributivity axiom becomes the following *linear* statement, which holds for every $x, y, z \in \mathbb{Q}Q$:

$$(x \uparrow y) \uparrow z = \uparrow \circ \uparrow_2(x \otimes y \otimes \Delta z). \tag{25} \quad \text{eq:LinSelf}$$

Clearly, we need to understand Δ better. By direct computation, if $x \in Q$ then $\Delta \bar{x} = \bar{x} \otimes 1 + 1 \otimes \bar{x} + \bar{x} \otimes \bar{x}$. We claim that in general, if z is a generating expression of \mathcal{I}^m (that is, a formula made of m elements of \bar{Q} and $m - 1$ applications of \uparrow), then

$$\Delta z = z \otimes 1 + 1 \otimes z + \sum z'_i \otimes z''_i, \quad \text{with} \quad \sum z'_i \otimes z''_i \in \sum_{\substack{m'+m''=m+1, \\ m',m''>0}} \mathcal{I}^{m'} \otimes \mathcal{I}^{m''}. \tag{26} \quad \text{eq:Quandle}$$

Indeed, for the generators of \mathcal{I}^1 this had just been shown, and if $z = z_1 \uparrow z_2$ is a generator of \mathcal{I}^m , with z_1 and z_2 generators of \mathcal{I}^{m_1} and \mathcal{I}^{m_2} with $1 \leq m_1, m_2 < m$ and $m_1 + m_2 = m$,

foot:upinv | ²⁹This can alternatively be stated as “there exists a second binary operation \uparrow^{-1} so that $\forall x, x = (x \uparrow y) \uparrow^{-1} y = (x \uparrow^{-1} y) \uparrow y$ ”, so this axiom can still be phrased within the language of “algebraic structures”. Yet note that below we do not use this axiom at all.

then (using $w \uparrow 1 = w$ and $1 \uparrow w = 0$ for $w \in \mathcal{I}$),

$$\begin{aligned} \Delta z &= \Delta(z_1 \uparrow z_2) = (\Delta z_1) \uparrow_2 (\Delta z_2) \\ &= (z_1 \otimes 1 + 1 \otimes z_1 + \sum z'_{1j} \otimes z''_{1j}) \uparrow_2 (z_2 \otimes 1 + 1 \otimes z_2 + \sum z'_{2k} \otimes z''_{2k}) \\ &= (z_1 \uparrow z_2) \otimes 1 + 1 \otimes (z_1 \uparrow z_2) \\ &\quad + \sum_j \left((z'_{1j} \uparrow z_2) \otimes z''_{1j} + z'_{1j} \otimes (z''_{1j} \uparrow z_2) + \sum_k (z'_{1j} \uparrow z'_{2k}) \otimes (z''_{1j} \uparrow z''_{2k}) \right), \end{aligned}$$

and it is easy to see that the last line agrees with (26).

We can now combine Equations (25) and (26) to get that for any $x, y, z \in \mathbb{Q}\mathbb{Q}$,

$$(x \uparrow y) \uparrow z = (x \uparrow z) \uparrow y + x \uparrow (y \uparrow z) + \sum (x \uparrow z'_i) \uparrow (y \uparrow z''_i).$$

If $x \in \mathcal{I}^{m_1}$, $y \in \mathcal{I}^{m_2}$, and $z \in \mathcal{I}^{m_3}$, then by (26) the last term above is in $\mathcal{I}^{m_1+m_2+m_3+1}$, and so the above identity becomes the Jacobi identity $(x \uparrow y) \uparrow z = (x \uparrow z) \uparrow y + x \uparrow (y \uparrow z)$ in $\text{proj}_{m_1+m_2+m_3} Q$.

MORE. It remains to show that within $\text{proj}_{>0} Q$, the operation \uparrow is anti-symmetric.

Exercise 6.7. Verify that in the above proof axiom (2) of Definition 6.5 was not used. Verify also that if this axiom is introduced as in footnote 29 using a second operation \uparrow^{-1} (thus enlarging the set of algebraic expressions that we need to consider as in MORE), Proposition 6.6 remains true.

Compare with Leibnitz:

$$[x, [y, z]] = [[x, y], z] - [[x, z], y] \xrightarrow{\text{Flip } []},$$

$$[[z, y], x] = [z, [y, x]] - [y, [z, x]] \xrightarrow{(x, y, z) \mapsto (z, y, x)}$$

$$[[x, y], z] = [x, [y, z]] - [y, [x, z]]$$

\Rightarrow It would be lovely to know that

$$(x \uparrow z) \uparrow y = -y \uparrow (x \uparrow z)$$